Game Semantics

Andrzej S. Murawski and Nikos Tzevelekos

Lecture 1b: Introduction to Game Semantics

Sequentiality

The problem with **parallel-or** is that programs are not really functions.

For example, a program with two inputs:

- either does not use one of the inputs,
- or, if it does, it picks one to use first.

► Put otherwise, programs are **sequential computations**.

This is what makes behaviours like parallel-or non-programmable.

This mismatch between

- functions in domain-based models
- and programs in PCF

makes us look beyond functions for a model of programs.

Dynamic behaviours

Example. Consider a program

count : int \rightarrow int

such that:

- the first time we call it, it returns 1;
- the second time we call it, it returns 2;
- • •
- \blacksquare the *i*-th time we call it, it returns *i*.

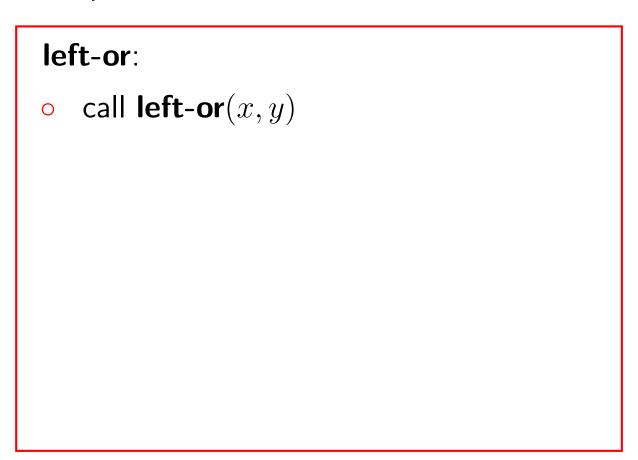
While we cannot write such a program in PCF, it is easy to do it in any language like Java, Python, OCaML, etc.

Programs can change their behaviour dynamically.

This is another source of mismatch between programs and functions.

Being sequential and dynamic, programs are best described as sequences of computation steps:

Being sequential and dynamic, programs are best described as sequences of computation steps:



Being sequential and dynamic, programs are best described as sequences of computation steps:

- o call **left-or**(x,y)
- \bullet evaluate x

Being sequential and dynamic, programs are best described as sequences of computation steps:

- call **left-or**(x, y)
- \bullet evaluate x
 - \circ x is 1
 - return 1

Being sequential and dynamic, programs are best described as sequences of computation steps:

left-or:

- \circ call **left-or**(x,y)
- evaluate x
 - \circ x is 1

 \circ x is 0

• return 1

- \bullet evaluate y
 - \circ y is v
 - \bullet return v

Being sequential and dynamic, programs are best described as games (or dialogues) between the program and its calling context:

left-or:

• what is the result of **left-or**(x, y)?

Being sequential and dynamic, programs are best described as games (or dialogues) between the program and its calling context:

- what is the result of **left-or**(x, y)?
- what is the value of x?

Being sequential and dynamic, programs are best described as games (or dialogues) between the program and its calling context:

- what is the result of **left-or**(x, y)?
- what is the value of x?
 - \circ x is 1
 - the result is 1

Being sequential and dynamic, programs are best described as games (or dialogues) between the program and its calling context:

- what is the result of **left-or**(x, y)?
- what is the value of x?
 - \circ x is 1
 - the result is 1

- \circ x is 0
- what is the result of *y*?
 - \circ y is v
 - the result is v

Programs \mapsto games

Programs \mapsto games

```
\vdash count : int \rightarrow int :
  what is the result of count?
```

Programs \mapsto games

```
\vdash count : int \rightarrow int :
• what is the result of count?
it is a function
```

```
\vdash count : int \rightarrow int :
```

- what is the result of count?
- it is a function
 - what is the result of the function on 42?

```
\vdash count : int \rightarrow int :
```

- what is the result of count?
- it is a function
 - what is the result of the function on 42?
 - it is 1

```
\vdash count : int \rightarrow int :
```

- what is the result of count?
- it is a function
 - what is the result of the function on 42?
 - it is 1
 - what is the result of the function on 23?
 - it is 2

Being sequential and dynamic, programs are best described as games (or dialogues) between the program and its calling context:

```
\vdash count : int \rightarrow int :
```

- what is the result of count?
- it is a function
 - what is the result of the function on 42?
 - it is 1
 - what is the result of the function on 23?
 - it is 2
 - what is the result of the function on 25?
 - it is 3

. . .

Higher-order programs

Example. Consider a program

$$f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$$

given by: **plusOne** $\equiv \lambda x^{\text{int}}.fx + 1$

- \blacksquare if we call it with a function f
- \blacksquare it returns a function that, on input x, returns f(x) + 1

Higher-order programs

Example. Consider a program

$$f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$$

given by: **plusOne** $\equiv \lambda x^{\text{int}}.fx + 1$

- \blacksquare if we call it with a function f
- lacktriangle it returns a function that, on input x, returns f(x)+1

E.g. taking
$$2x \equiv \lambda y^{\text{int}} \cdot 2 * y : \text{int} \rightarrow \text{int} :$$

$$(\lambda f.\mathsf{plusOne}) \ \mathbf{2x} \ 0 \ \longrightarrow \ (\lambda x^{\mathsf{int}}.\mathbf{2x} \ x+1) \ 0 \ \longrightarrow \ \mathbf{2x} \ 0+1 \ \longrightarrow^* \ 1$$

$$(\lambda f.\mathsf{plusOne}) \ \mathbf{2x} \ 42 \ \longrightarrow \ (\lambda x^{\mathsf{int}}.\mathbf{2x} \ x+1) \ 42 \ \longrightarrow \ \mathbf{2x} \ 42+1 \ \longrightarrow^* \ 85$$

```
Example. Consider a program f: \text{int} \to \text{int} \vdash \textbf{plusOne}: \text{int} \to \text{int} given by: \textbf{plusOne} \equiv \lambda x^{\text{int}}.fx + 1.
```

- given function f, what is the result of **plusOne**?
- it is a function f'

```
Example. Consider a program f: \text{int} \to \text{int} \vdash \textbf{plusOne}: \text{int} \to \text{int} given by: \textbf{plusOne} \equiv \lambda x^{\text{int}}.fx + 1.
```

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on 42?

```
Example. Consider a program f: \text{int} \to \text{int} \vdash \textbf{plusOne}: \text{int} \to \text{int} given by: \textbf{plusOne} \equiv \lambda x^{\text{int}}.fx + 1.
```

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne} : \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne} \equiv \lambda x^{\operatorname{int}} \cdot fx + 1.
```

- given function f, what is the result of **plusOne**?
- ullet it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 84

```
Example. Consider a program f: \text{int} \to \text{int} \vdash \text{plusOne}: \text{int} \to \text{int}
given by: plusOne \equiv \lambda x^{\text{int}}.fx + 1.
Game:
• given function f, what is the result of plusOne?
• it is a function f'
   • what is the result of f' on 42?
       • what is the result of f on 42?
       o it is 84
   • it is 85
```

```
Example. Consider a program f: \text{int} \to \text{int} \vdash \textbf{plusOne}: \text{int} \to \text{int} given by: \textbf{plusOne} \equiv \lambda x^{\text{int}}.fx + 1.
```

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne} : \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne} \equiv \lambda x^{\operatorname{int}} \cdot fx + 1.
```

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 35

```
Example. Consider a program f: \text{int} \to \text{int} \vdash \text{plusOne}: \text{int} \to \text{int}
given by: plusOne \equiv \lambda x^{\text{int}}.fx + 1.
Game:
• given function f, what is the result of plusOne?
• it is a function f'
   • what is the result of f' on 42?
       • what is the result of f on 42?
       o it is 35
   • it is 71
```

More higher-order programs

Example. Consider a program

$$f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne}^2: \mathsf{int} \to \mathsf{int}$$

given by: $\mathbf{plusOne}^2 \equiv \lambda x^{\mathsf{int}}. f(fx+1) + 1$

- lacksquare if we call it with a function f
- \blacksquare it returns a function that, on input x, returns f(f(x) + 1) + 1

More higher-order programs

Example. Consider a program

$$f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne}^2: \mathsf{int} \to \mathsf{int}$$

given by:
$$\mathbf{plusOne}^2 \equiv \lambda x^{\mathsf{int}} \cdot f(fx+1) + 1$$

- \blacksquare if we call it with a function f
- \blacksquare it returns a function that, on input x, returns f(f(x)+1)+1

E.g. taking
$$2x \equiv \lambda y^{\text{int}} \cdot 2 * y : \text{int} \rightarrow \text{int} :$$

$$(\lambda f.\mathsf{plusOne}^2) \, \mathbf{2x} \, 0 \longrightarrow (\lambda x^\mathsf{int}. \, \mathbf{2x} (\mathbf{2x} \, x + 1) + 1) \, 0 \longrightarrow \mathbf{2x} (\mathbf{2x} \, 0 + 1) + 1 \longrightarrow^* 1$$

$$(\lambda f.\mathsf{plusOne}^2) \, \mathbf{2x} \, 42 \longrightarrow (\lambda x^\mathsf{int}. \, \mathbf{2x}(\mathbf{2x} \, x+1)+1) \, 42 \longrightarrow \mathbf{2x}(\mathbf{2x} \, 42+1)+1 \longrightarrow^* 171$$

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of *f* on 42?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of *f* on 42?
 - o it is 84

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 84
 - what is the result of *f* on 85?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 84
 - what is the result of f on 85?
 - o it is 170

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

Game:

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 84
 - what is the result of f on 85?
 - o it is 170
 - it is 171

. . .

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of *f* on 42?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of *f* on 42?
 - o it is 35

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 35
 - what is the result of *f* on 36?

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 35
 - what is the result of f on 36?
 - o it is 15

```
Example. Consider a program f: \operatorname{int} \to \operatorname{int} \vdash \operatorname{plusOne}^2: \operatorname{int} \to \operatorname{int} given by: \operatorname{plusOne}^2 \equiv \lambda x^{\operatorname{int}}. f(fx+1)+1.
```

Game:

- given function f, what is the result of **plusOne**²?
- it is a function f'
 - what is the result of f' on 42?
 - what is the result of f on 42?
 - o it is 35
 - what is the result of *f* on 36?
 - o it is 15
 - it is 16

. . .

- games are sequences of **moves**, called **plays**
- following some **formal conditions**:

- games are sequences of moves, called plays
- following some formal conditions:
 - games are played between two players:
 - Opponent (O), that represents the program's context
 - Proponent (P), that represents the program

Revisit count

Being sequential and dynamic, programs are best described as games between a Proponent (P) and an Opponent (O):

Revisit count

Being sequential and dynamic, programs are best described as games between a Proponent (P) and an Opponent (O):

```
\vdash count : int \rightarrow int :
   what is the result of count?
  it is a function
  O what is the result of the function on 42?
  P it is 1
     what is the result of the function on 23?
  P it is 2
     what is the result of the function on 25?
     it is 3
```

- games are sequences of moves, called plays
- following some formal conditions:
 - games are played between two players:
 - Opponent (O), that represents the program's context
 - Proponent (P), that represents the program

- games are sequences of **moves**, called **plays**
- following some formal conditions:
 - games are played between two players:
 - Opponent (O), that represents the program's context
 - Proponent (P), that represents the program
 - ◆ P and O alternate; O starts first

- games are sequences of moves, called plays
- following some formal conditions:
 - games are played between two players:
 - Opponent (O), that represents the program's context
 - Proponent (P), that represents the program
 - ◆ *P* and *O* alternate; *O* starts first
- moves come in two forms:
 - lacktriangle moves that call functions are **questions** (Q)
 - lacktriangle moves that return function calls are **answers** (A)

Revisit count

Being sequential and dynamic, programs are best described as games between a Proponent (P) and an Opponent (O):

Revisit count

Being sequential and dynamic, programs are best described as games between a Proponent (P) and an Opponent (O):

```
\vdash count : int \rightarrow int :
OQ what is the result of count?
PA it is a function
   OQ what is the result of the function on 42?
   PA it is 1
   OQ what is the result of the function on 23?
   PA it is 2
   OQ what is the result of the function on 25?
   PA it is 3
```

- games are sequences of moves, called plays
- following some formal conditions:
 - games are played between two players:
 - Opponent (O), that represents the program's context
 - Proponent (P), that represents the program
 - lacktriangle P and O alternate; O starts first
 - moves come in two forms:
 - lacktriangle moves that call functions are **questions** (Q)
 - \blacksquare moves that return function calls are **answers** (A)

- games are sequences of moves, called plays
- following some formal conditions:
 - games are played between two players:
 - Opponent (O), that represents the program's context
 - Proponent (P), that represents the program
 - lacktriangle P and O alternate; O starts first
 - moves come in two forms:
 - \blacksquare moves that call functions are **questions** (Q)
 - \blacksquare moves that return function calls are **answers** (A)
 - lacktriangle each A refers to a Q by the opposite player
 - **•** ...

 $\llbracket \vdash \mathbf{count} : \mathsf{int} \to \mathsf{int} \rrbracket$ is a set of plays of the form:

$$\star$$
 † 42 1 23 2 25 3 \cdots

where:

the first move is played by O and asks the result of **count**, given an empty context (so, \star is a move representing the empty context)

 $\llbracket \vdash \mathbf{count} : \mathsf{int} \to \mathsf{int} \rrbracket$ is a set of plays of the form:

$$\star$$
 † 42 1 23 2 25 3 \cdots

where:

- the first move is played by O and asks the result of **count**, given an empty context (so, \star is a move representing the empty context)
- the second move is played by P and answers the initial question saying the result is a function (so, \dagger is a move representing a function)

 $\llbracket \vdash \mathbf{count} : \mathsf{int} \to \mathsf{int} \rrbracket$ is a set of plays of the form:

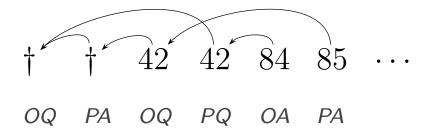
$$\star$$
 † 42 1 23 2 25 3 \cdots

where:

- the first move is played by O and asks the result of **count**, given an empty context (so, \star is a move representing the empty context)
- \blacksquare the second move is played by P and answers the initial question saying the result is a function (so, \dagger is a move representing a function)
- from there on, we engage in a OQ-PA pattern:
 - O asks the result of the function on some input number
 - ◆ P answers by simply playing the number of times it has been called

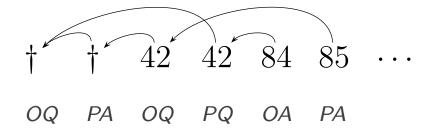
Recall **plusOne** $\equiv \lambda x^{\text{int}}.fx + 1.$

 $\llbracket f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne}: \mathsf{int} \to \mathsf{int} \rrbracket$ is a set of plays of the form:



Recall **plusOne** $\equiv \lambda x^{\text{int}}. fx + 1.$

 $[\![f:\mathsf{int}\to\mathsf{int}\vdash\mathsf{plusOne}:\mathsf{int}\to\mathsf{int}]\!]$ is a set of plays of the form:

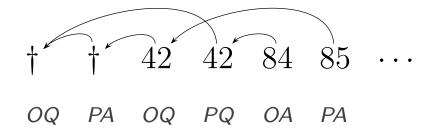


where, additionally to what we saw in [count],

- we also have **pointers** between moves
- each move has a pointer to an earlier move that justifies it

Recall **plusOne** $\equiv \lambda x^{\text{int}}.fx + 1.$

 $[\![f:\mathsf{int}\to\mathsf{int}\vdash\mathsf{plusOne}:\mathsf{int}\to\mathsf{int}]\!]$ is a set of plays of the form:

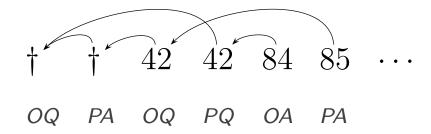


where, additionally to what we saw in [count],

- we also have pointers between moves
- each move has a pointer to an earlier move that justifies it
 - e.g. an answer points to its corresponding question

Recall **plusOne** $\equiv \lambda x^{\text{int}}.fx + 1.$

 $[\![f:\mathsf{int}\to\mathsf{int}\vdash\mathsf{plusOne}:\mathsf{int}\to\mathsf{int}]\!]$ is a set of plays of the form:

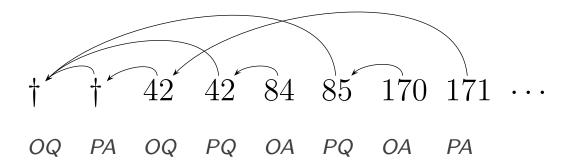


where, additionally to what we saw in [count],

- we also have **pointers** between moves
- each move has a pointer to an earlier move that justifies it
 - e.g. an answer points to its corresponding question
 - ◆ a question points to the † move that the question refers to
 - e.g. the 42 played by O is a question to second \dagger (i.e. f')
 - whereas the 42 played by P is a question to first \dagger (i.e. f)

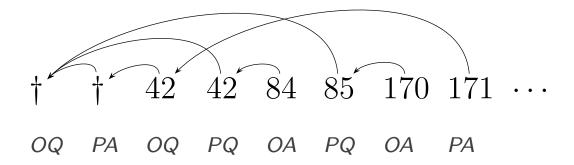
Recall **plusOne**² $\equiv \lambda x^{\text{int}}. f(fx+1) + 1.$

 $[f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne}^2: \mathsf{int} \to \mathsf{int}]$ is a set of plays of the form:



Recall **plusOne**² $\equiv \lambda x^{\text{int}} \cdot f(fx+1) + 1$.

 $[f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne}^2 : \mathsf{int} \to \mathsf{int}]$ is a set of plays of the form:



So, plays combine:

- the idea of executing a program and exchanging moves with its context
- with those moves potentially representing higher-order functions
 - the exchanged moves themselves enable more moves to be played
 - pointers are used to keep an order on what-is-played-where

We know that P in a game is the modelled program. But, who is O?

■ It is another program. But then, that program also has a game semantics!

We know that P in a game is the modelled program. But, who is O?

■ It is another program. But then, that program also has a game semantics!

```
\vdash \lambda y^{\mathsf{int}}.2 * y : \mathsf{int} \to \mathsf{int}
```

O what is the result?

We know that P in a game is the modelled program. But, who is O?

■ It is another program. But then, that program also has a game semantics!

```
\vdash \lambda y^{\mathsf{int}}.2 * y : \mathsf{int} \to \mathsf{int}
```

- O what is the result?
- P it is a function

We know that P in a game is the modelled program. But, who is O?

■ It is another program. But then, that program also has a game semantics!

```
\vdash \lambda y^{\mathsf{int}}.\ 2*y:\mathsf{int}\to\mathsf{int}
```

- O what is the result?
- P it is a function
- *O* what is the result for 42?

We know that P in a game is the modelled program. But, who is O?

■ It is another program. But then, that program also has a game semantics!

```
\vdash \lambda y^{\mathsf{int}}.2 * y : \mathsf{int} \to \mathsf{int}
```

- O what is the result?
- P it is a function
- *O* what is the result for 42?
- *P* it is 84

We know that P in a game is the modelled program. But, who is O?

It is another program. But then, that program also has a game semantics!

- $\vdash \lambda y^{\mathsf{int}}.2 * y : \mathsf{int} \to \mathsf{int}$
- O what is the result?
- P it is a function
- O what is the result for 42?
- *P* it is 84

- $f: \mathsf{int} \to \mathsf{int} \vdash \lambda x^{\mathsf{int}}. fx + 1: \mathsf{int} \to \mathsf{int}$
- $oldsymbol{O}$ given function f, what is the result?
- P it is a function
- O what is the result for 42?
- P what is the result of f for 42?
- *O* it is 84
- *P* it is 85

We know that P in a game is the modelled program. But, who is O?

■ It is another program. But then, that program also has a game semantics!

$$\vdash \lambda y^{\mathsf{int}}.2 * y : \mathsf{int} \to \mathsf{int}$$

- O what is the result?
- P it is a function

- *O* what is the result for 42?
- *P* it is 84

$$f: \mathsf{int} \to \mathsf{int} \vdash \lambda x^{\mathsf{int}}. fx + 1: \mathsf{int} \to \mathsf{int}$$

- $oldsymbol{O}$ given function f, what is the result?
- P it is a function
- O what is the result for 42?
- P what is the result of f for 42?
- O it is 84
- *P* it is 85

If we write moves formally and re-arrange moves in space so that we see where they come from

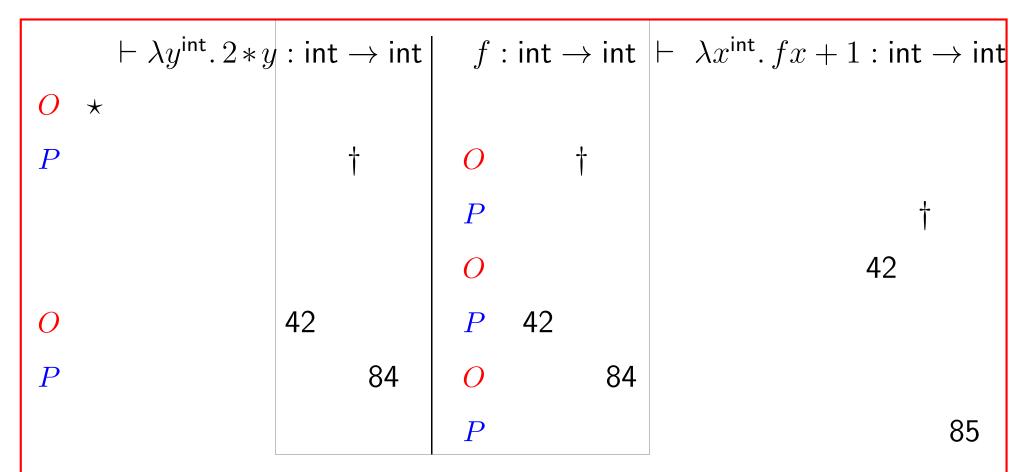
 \blacksquare we see a **duality** between P and O in the middle component

ŀ	$-\lambda y^{int}.2*y$	$:int\toint$	f	: int —	→ int	 - λa	$z^{int}. f$	x +	1 : in	t o i	int
0 *											
P		†	O	-	-						
			O P							†	
			0						42		
0		42	P	42							
P		84	0		84						
			P							85)
						_					

Game duality

If we write moves formally and re-arrange moves in space so that we see where they come from

 \blacksquare we see a duality between P and O in the middle component



What is P on the LHS of the grey box, is O on the RHS, and viceversa.

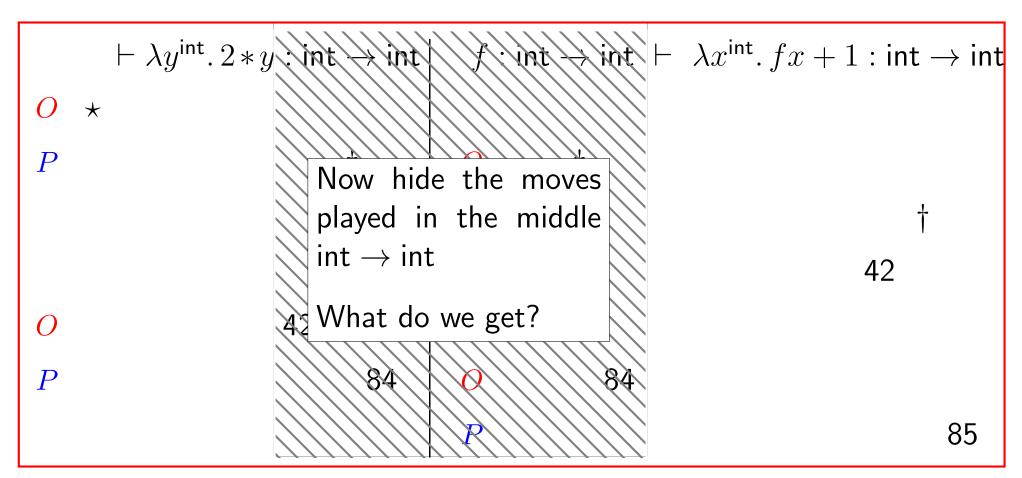
If we write moves formally and re-arrange moves in space so that we see where they come from

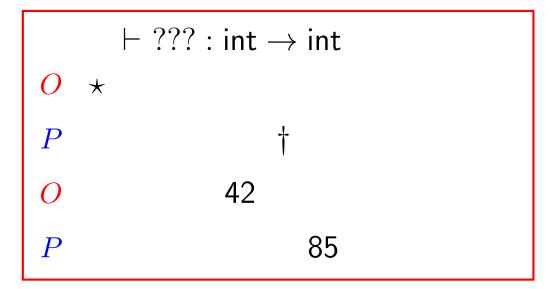
 \blacksquare we see a **duality** between P and O in the middle component

$\vdash \lambda$	$y^{int}.2\!*\!y$: in^{r}	t o int	f	: int o int	nt	$\lambda x^{\mathrm{int}}.fx+1:\mathrm{int} o \mathrm{int}$
O *						
P		†	0	†		
			P			†
			0			42
0	42		P	42		
P		84	O	8	34	
			P			85

If we write moves formally and re-arrange moves in space so that we see where they come from

 \blacksquare we see a duality between P and O in the middle component





$$\vdash ??? : \text{int} \rightarrow \text{int}$$
 $O \star$
 $P \uparrow$
 $O \downarrow$
 O

We get the game corresponding to the **composition** of the two programs:

$$\vdash \operatorname{let}(f = \lambda y^{\operatorname{int}}. 2 * y) \operatorname{in} \lambda x^{\operatorname{int}}. fx + 1$$

i.e.
$$\vdash \lambda x^{\mathsf{int}} \cdot 2 * x + 1$$
.

So, we can compose games with a common right/left component:

- synchronising moves in common component (using duality)
- and hiding those moves.

This is analogous to how functions compose.

Example. Consider a program $f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$, where $\mathsf{plusOne} \equiv \lambda x^{\mathsf{int}} \cdot fx + 1$.

- given function f, what is the result of **plusOne**?
- it is a function f'

Example. Consider a program $f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$, where $\mathsf{plusOne} \equiv \lambda x^{\mathsf{int}} \cdot fx + 1$.

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on integer i?

Example. Consider a program $f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$, where $\mathsf{plusOne} \equiv \lambda x^{\mathsf{int}} \cdot fx + 1$.

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on integer i?
 - what is the result of f on i?

Example. Consider a program $f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$, where $\mathsf{plusOne} \equiv \lambda x^{\mathsf{int}} \cdot fx + 1$.

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on integer i?
 - what is the result of f on i?
 - \circ it is j

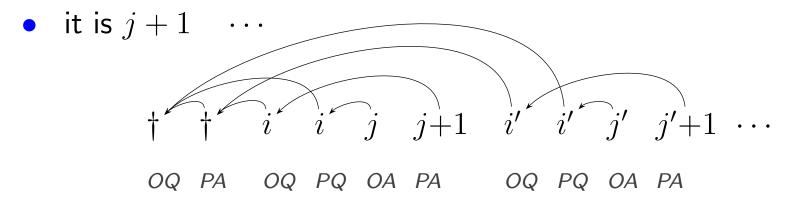
Example. Consider a program $f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$, where $\mathsf{plusOne} \equiv \lambda x^{\mathsf{int}} \cdot fx + 1$.

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on integer i?
 - what is the result of f on i?
 - \circ it is j
 - it is j+1 ····

Example. Consider a program $f: \mathsf{int} \to \mathsf{int} \vdash \mathsf{plusOne} : \mathsf{int} \to \mathsf{int}$, where $\mathsf{plusOne} \equiv \lambda x^{\mathsf{int}} \cdot fx + 1$.

The corresponding game has plays for every possible behaviour of f:

- given function f, what is the result of **plusOne**?
- it is a function f'
 - what is the result of f' on integer i?
 - what is the result of f on i?
 - \circ it is j



[plusOne] contains all plays of that form (i.e. for all $i, j, i', j', \dots \in \mathbb{Z}$)

Results

Games model programs under any possible context:

- they contain plays for every O move allowed
- they include conditions that disallow spurious plays.

And this is the key to full-abstraction results:

$$M \cong N \iff \llbracket M \rrbracket = \llbracket N \rrbracket$$

Results

Games model programs under any possible context:

- they contain plays for every O move allowed
- they include conditions that disallow spurious plays.

And this is the key to full-abstraction results:

$$M \cong N \iff \llbracket M \rrbracket = \llbracket N \rrbracket$$

After the original papers on PCF, there was a series of works covering extensions of PCF with effects: local state, local higher-order state, non-determinism, probabilities, control operators, etc.

More recently, games have been extended to languages with data-generating effects, like references, objects, channels, etc. Nowadays, games capture a wide-range of higher-order languages, typically fragments of OCaML and Java.

See the Lecture Notes for references.

Exercises

1. Consider the following alternative notion of equivalence for PCF terms:

Given $\Gamma \vdash M_1 : \theta$ and $\Gamma \vdash M_2 : \theta$, we let $M_1 \cong' M_2$ if, for every context C such that $\vdash C[M_1] : \text{int and } \vdash C[M_2] : \text{int, and for all } i \in \mathbb{Z}, \ C[M_1] \longrightarrow^* i \iff C[M_2] \longrightarrow^* i.$

Prove that \cong' coincides with \cong .

- 2. We use the following shorthand notation: let x=M in $N\equiv (\lambda x.N)M$. Using the operational semantics of PCF:
 - Verify that $(\operatorname{let}(f=\lambda y^{\operatorname{int}}.\ 2*y)\operatorname{in}\lambda x^{\operatorname{int}}.\ fx+1)z \longrightarrow^* 2*z+1.$
 - Compute $(\operatorname{let}(f = \lambda y^{\operatorname{int}}. \ 2 * y) \operatorname{in} \lambda x^{\operatorname{int}}. \ f(fx+1) + 1) z.$
- 3. Working as in Slides 24-25, compose game-semantically $\vdash \lambda y^{\text{int}}.\ 2*y: \text{int} \to \text{int}$ and $f: \text{int} \to \text{int} \vdash \lambda x^{\text{int}}.\ f(fx+1)+1: \text{int} \to \text{int}.$
- 4. Consider the program

$$\vdash \lambda f^{\mathsf{int} \to \mathsf{int}} . \lambda x^{\mathsf{int}} . fx + 1 : (\mathsf{int} \to \mathsf{int}) \to (\mathsf{int} \to \mathsf{int})$$

obtained by λ -abstracting **plusOne**. Working first informally (using dialogues) and then formally (using plays), express the plays in $[\![\lambda f^{\mathsf{int} \to \mathsf{int}}.\lambda x^{\mathsf{int}}.fx+1]\!]$:

- First, assuming that O plays a question \dagger (for f) only once.
- Consider how the plays would look like if O plays the † more than once.