Copy-Cat Strategies and Information Flow in Physics, Geometry, Logic and Computation

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- How to Beat a Grand-Master
- Does Copy-Cat still work here?
- Some themes

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The Copy-Cat Strategy
Does Copy-Cat still work here?

Kasparov

W
B

Short

W
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Short

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- Correspondence between interactive and geometric views of the same phenomena
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- Correspondence between interactive and geometric views of the same phenomena
- Emergence
- Copy-cat vs. Cloning:
  - Linear vs. Classical Logic
  - Quantum vs. Classical Physics
  - Linear-time vs. Computationally universal
Some themes

- Common fundamental structures of interaction: logical, computational, physical, geometric

- Correspondence between *interactive* and *geometric* views of the same phenomena

- Emergence

- Copy-cat vs. Cloning:
  - Linear vs. Classical Logic
  - Quantum vs. Classical Physics
  - Linear-time vs. Computationally universal

- Many pictures — serious mathematical structure underneath! (Joyal, Street, Kelly, Penrose, ...)

- Game aspect not emphasized — more at primitive interaction level (‘GoI’).
The essential information in a (cut-free) proof in MLL is the axiom links. Accordingly, we define a proof structure on a sequent $\Gamma$ to be a fixpoint-free involution $f$ (so $f^2 = 1$ and $f(a) \neq a$) on its occurrences of literals such that if $f(a) = b$, $l(a) = l(b)^\bot$. 
From Proof Nets to Semantics

Note that proof structures as we have defined them are simply certain permutations acting on finite sets (of literals). This leads to the following compositional interpretation of formulas by finite sets, and of proofs by permutations on these sets.

- A literal is interpreted by a one-point set; Tensor and Par by disjoint union. A sequent is treated like the Par of its formulas. Thus the set $|\Gamma|$ associated to a sequent is in bijection with its set of occurrences of literals.
Assignment of Permutations to Sequent Proofs

Axiom

\[ \vdash a, a^\perp \text{ Id} \]

Multiplicatives

\[ \vdash \Gamma, A \vdash \Delta, B \quad \vdash \Gamma, \Delta, A \otimes B \]  
\[ \vdash \Gamma, A \otimes B \]

- Axiom: assign the transposition \( a \leftrightarrow a^\perp \)
- Tensor: assign the disjoint union of the two permutations
- Par: assign the same permutation!
MLL Proof Nets

Which proof structures really come from proofs in MLL?

Switching Graphs: A switching $S$ of $\Gamma$ assigns $L$ or $R$ to each occurrence of $\otimes$. Given a sequent $\Gamma$, a proof structure $f$, and a switching $S$, the switching graph $G(\Gamma, f, S)$ has:

- subformula occurrences in $\Gamma$ as vertices;
- an edge connecting $A$ to $A \otimes B$ and an edge connecting $B$ to $A \otimes B$ for each occurrence of $A \otimes B$;
- an edge connecting $A$ to $A \otimes B$ if $S$ assigns $L$ to $A \otimes B$, and an edge connecting $B$ to $A \otimes B$ if $S$ assigns $R$ to $A \otimes B$;
- an edge connecting literal occurrences $a$ and $b$ if $f(a) = b$.

The Danos-Regnier criterion: A proof-structure $f$ for $\Gamma$ is an MLL proof-net if for every switching $\;S$, $G(\Gamma, f, S)$ is acyclic and connected.
Results on Proof Nets

Every sequent proof in MLL canonically maps to a proof structure.

Proposition 1 (Soundness)  The proof structures arising from sequent proofs are proof nets.

Theorem 2 (Sequentialization Theorem)  Every proof net arises from a sequent proof.

This is the Geometric Criterion.
We can think of a proof structure (set of axiom links) as a **copy-cat strategy**, and a switching as a counter-strategy. A proof structure will be a proof-net if its interaction with every counter-strategy yields a correct result. Hence we define (Girard 1988):

\[ f \perp g \equiv fg \text{ is cyclic} \]

\[ i.e. \ (fg)^k = 1 \text{ where } k \text{ is the cardinality of the underlying set (of literal occurrences), and this is true for no smaller value of } k. \]

This condition is directly inspired by the **long trip condition**, the earlier version of the proof net correctness condition used by (Girard 1987).

We can then define

\[ S^\perp = \{ g \mid \forall f \in S. f \perp g \}. \]
Semantics of MLL Proofs

We now give a semantics of MLL proofs by specifying, for each formula $A$, a set $S$ of permutations on the set of literal occurrences $|A|$, such that $S = S^{\perp \perp}$.

For a literal, we specify the unique permutation (the identity).

$$S(A \otimes B) = \{ f + g \mid f \in S(A) \land g \in S(B) \}^{\perp \perp}$$

$$S(A \land B) = S(A^{\perp} \otimes B^{\perp})^{\perp}.$$

Note that, for every formula $A$: $S(A^{\perp}) = S(A)^{\perp}$.

We extend this assignment to sequents $\Gamma$ by treating $\Gamma$ as the Par of its formulas.
Semantics: Soundness and Completeness

Proposition 3 (Semantic Soundness) If $f$ is the permutation assigned to a sequent proof of $\Gamma$, then $f \in S(\Gamma)$. 
Semantics: Soundness and Completeness

**Proposition 5 (Semantic Soundness)**  If $f$ is the permutation assigned to a sequent proof of $\Gamma$, then $f \in S(\Gamma)$.

**Theorem 6 (Full Completeness)**  If $f \in S(\Gamma)$ is a literal-respecting involution, then $f$ is a proof-net, and hence is the denotation of a sequent proof.

This shows the equivalence of the geometric and interactive criteria for proofs.

**Proof Outline:** Given $\sigma \in S(\Gamma)$, we assume that for some switching $S$, $G(\Gamma, \sigma, S')$ is not a tree. Then we construct a counter-strategy $\tau \in S(\Gamma)^\perp$ such that $\neg(\sigma \perp \tau)$. Contradiction.

**Discussion:** Non-emptiness of $S(A)$, “paraproofs”, and uniformity.
We consider performing Cut-elimination on $\text{twist} \circ \text{twist}$: The proof net for $\text{twist} \circ \text{twist}$ before cut elimination is:

![Proof net before cut elimination](image)

The proof net for $\text{twist} \circ \text{twist}$ after one step of Cut elimination is:

![Proof net after cut elimination](image)
Generally, in this fragment, we can apply this ‘decomposition rule’ repeatedly for tensors cut against par until all cuts are between axiom links. We can say that the whole purpose of these transformations is to match up the corresponding axiom links correctly; the ‘real’ information flow is then accomplished by the axiom reductions:

$$\alpha \perp \alpha \alpha \perp \alpha \alpha$$

or more generally,

$$\alpha \perp \alpha \alpha \perp \alpha \alpha \ldots \alpha \perp \alpha \alpha$$

Two views: geometric and interactive.
From Proof Nets to Diagram Algebras

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Copy-Cat Strategies and Information Flow
We now make a transition from an apparently very specialized corner of Proof Theory to a broad topic arising in Representation Theory, Knot Theory, and with connections to Mathematical Physics.

We shall on the one hand **lose** some structure, and on the other **gain** some.

- We shall obliterate the distinction between $\otimes$ and $\otimes$; this corresponds to moving from $\ast$-autonomous to **compact closed** categories. This means that we can forget about the formula tree structure altogether; we are simply connecting up literal occurrences, which we shall draw as “joining up the dots”.

Motivation: compact closed categories show up in many contexts of interest!
Diagrams for Arrows

- On the other hand, rather than one-sided sequents, we shall represent general arrows or two-sided sequents diagrammatically. This means we represent arrows

\[
A_1 \otimes \cdots \otimes A_n \rightarrow B_1 \otimes \cdots \otimes B_m
\]

where each \(A_i\) and \(B_j\) is a literal. We represent such arrows by literal-preserving involutions on \(\{1, \ldots, n\} + \{1, \ldots, m\}\), where literal-preserving now means:

- We connect **opposite** literals in the domain or codomain, or
- We connect occurrences of the **same** literal in the domain and the codomain.

An advantage of this representation is that we express composition very transparently, by “stacking” arrows.
Example

The composition of

\[
\begin{align*}
\sigma & \sigma^* \sigma \\
\sigma & \sigma^* \sigma
\end{align*}
\]

and

\[
\begin{align*}
\tau & \tau^* \tau \\
\tau & \tau^* \tau
\end{align*}
\]

is given by

\[
\begin{align*}
\sigma & \sigma^* \sigma^* \sigma \\
\sigma & \sigma^* \sigma^* \sigma
\end{align*}
\]

\[=\]

\[
\begin{align*}
\tau & \tau^* \tau \\
\tau & \tau^* \tau
\end{align*}
\]
Temperley-Lieb Algebra

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Temperley-Lieb Algebra

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Temperley-Lieb Algebra

The **Temperley-Lieb algebra** played a central role in the **Jones polynomial invariant of knots** and ensuing developments.

The TL algebra was originally presented, rather forbiddingly, in terms of abstract generators and relations. It was recast in beautifully elementary and conceptual terms by Louis Kauffman as a **planar diagram algebra**.
TL algebra: generators and relations

We fix a ring $R$. Given a choice of parameter $\tau \in R$ and a dimension $n \in \mathbb{N}$, we define the Temperley-Lieb algebra $A_n(\tau)$ to be the unital, associative $R$-linear algebra with generators

$$U_1, \ldots, U_{n-1}$$

and relations

$$U_i U_j U_i = U_i \quad |i - j| = 1$$

$$U_i^2 = \tau \cdot U_i$$

$$U_i U_j = U_j U_i \quad |i - j| > 1$$

Note that the only relations used in defining the algebra are multiplicative ones. This suggests that we can present the multiplicative monoid $\mathcal{M}_n$, and then obtaining $A_n(\tau)$ as the **monoid algebra** of formal $R$-linear combinations $\sum_i r_i \cdot a_i$ over $\mathcal{M}_n$, with multiplication defined by bilinear extension:

$$\left(\sum_i r_i \cdot a_i\right) \left(\sum_j s_j \cdot b_j\right) = \sum_{i,j} (r_i s_j) \cdot (a_i b_j).$$
We define $\mathcal{M}_n$ as the monoid with generators
\[
\delta, U_1, \ldots, U_{n-1}
\]
and relations
\[
\begin{align*}
U_i U_j U_i &= U_i & |i - j| &= 1 \\
U_i^2 &= \delta U_i \\
U_i U_j &= U_j U_i & |i - j| &> 1 \\
\delta U_i &= U_i \delta
\end{align*}
\]
We can then obtain $\mathcal{A}_n(\tau)$ as the monoid algebra over $\mathcal{M}_n$, subject to the identification
\[
\delta = \tau \cdot 1.
\]
Diagram Monoids: Generators

We start with two parallel rows of \( n \) dots (geometrically, points in the plane). An element of the monoid is obtained by “joining up the dots” pairwise in a smooth, planar fashion, where the arc connecting each pair of dots must lie within the rectangle framing the two parallel rows of dots. Such diagrams are identified up to planar isotopy, i.e. continuous deformation within the portion of the plane bounded by the framing rectangle.

The generators \( U_1, \ldots, U_{n-1} \) can be drawn as follows:

\[
\begin{align*}
1 & \quad 2 \quad 3 \quad \cdots \quad n \\
\bigcirc & \quad \bigcirc \\
1 & \quad 2 \quad 3 \quad \cdots \quad n
\end{align*}
\]

\[
\begin{align*}
1 & \quad 2 \quad 3 \quad \cdots \quad n \\
\bigcirc & \quad \bigcirc \\
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\]

\[
\begin{align*}
1 & \quad 2 \quad 3 \quad \cdots \quad n \\
\bigcirc & \quad \bigcirc \\
1 & \quad 2 \quad 3 \quad \cdots \quad n
\end{align*}
\]

The generator \( \delta \) corresponds to a loop \( \bigcirc \) — all such loops are identified up to isotopy.
Diagram Monoids: Relations

We refer to arcs connecting dots in the top row as **cups**, those connecting dots in the bottom row as **caps**, and those connecting a dot in the top row to a dot in the bottom row as **through lines**.

Multiplication $xy$ is defined by identifying the bottom row of $x$ with the top row of $y$, and composing paths. In general loops may be formed — these are “scalars”, which can float freely across these figures. The relations can be illustrated as follows:

\[
\begin{align*}
U_1 U_2 U_1 &= U_1 \\
U_1^2 &= \delta U_1 \\
U_1 U_3 &= U_3 U_1
\end{align*}
\]
Expressiveness of the Generators

The fact that all planar diagrams can be expressed as products of generators is not entirely obvious. As an illustrative example, consider the planar diagrams in $\mathcal{M}_3$. Apart from the generators $U_1, U_2$, and ignoring loops, there are three:

The first is the identity for the monoid; we refer to the other two as the left wave and right wave respectively. The left wave can be expressed as the product $U_2 U_1$:

The right wave has a similar expression.
Nested Cups and Caps

Once we are in dimension four or higher, we can have nested cups and caps. These can be built using waves, as illustrated by the following:
### The Trace

There is a natural **trace function** on the Temperley-Lieb algebra, which can be defined diagrammatically on $\mathcal{M}_n$ by connecting each dot in the top row to the corresponding dot in the bottom row, using auxiliary cups and cups. This always yields a diagram isotopic to a number of loops — hence to a **scalar**, as expected. This trace can then be extended linearly to $\mathcal{A}_n(\tau)$.

We illustrate this firstly by taking the trace of a wave—which is equal to a single loop:

```
\begin{array}{c}
\text{The Ear is a Circle}
\end{array}
```
Our second example illustrates the important general point that the trace of the identity in $M_n$ is $\delta^n$:
The Connection to Knots

How does this connect to knots? Again, a key conceptual insight is due to Kauffman, who saw how to recast the Jones polynomial in elementary combinatorial form in terms of his bracket polynomial. The basic idea of the bracket polynomial is expressed by the following equation:

\[
\begin{align*}
\langle \begin{array}{c}
\text{\textbullet} \\
\text{\textbullet}
\end{array} \rangle &= A \langle \begin{array}{c}
\text{\\\textbullet} \\
\text{\\\textbullet}
\end{array} \rangle + B \langle \begin{array}{c}
\text{\emptyset}
\end{array} \rangle
\end{align*}
\]

Each over-crossing in a knot or link is evaluated to a weighted sum of the two possible planar smoothings. With suitable choices for the coefficients \( A \) and \( B \) (as Laurent polynomials), this is invariant under the second and third Reidemeister moves. With an ingenious choice of normalizing factor, it becomes invariant under the first Reidemeister move — and yields the Jones polynomial!
An Algebraic View

What this means algebraically is that the braid group $B_n$ has a representation in the Temperley-Lieb algebra $A_n(\tau)$ — the above bracket evaluation shows how the generators $\beta_i$ of the braid group are mapped into the Temperley-Lieb algebra:

$$\beta_i \mapsto A \cdot U_i + B \cdot 1.$$ 

Every knot arises as the closure (i.e. the diagrammatic trace) of a braid; the invariant arises by mapping the open braid into the Temperley-Lieb algebra, and taking the trace there.

This is just the beginning of a huge swathe of further developments, including: Topological Quantum Field Theories, Quantum Groups, Quantum Statistical mechanics, Diagram Algebras and Representation Theory, and more.
Characterizing Planarity

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- First condition
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Characterizing Planarity

A map \( f \in \text{Inv}(N(n, m)) \) will be called **planar** if it satisfies the following two conditions, for all \( i, j \in N(n, m) \):

1. \( i < j < f(i) \implies f(j) < f(i) \) (PL1)
2. \( f(i) \neq i < j \neq f(j) \implies f(i) < f(j) \) (PL2)

It is instructive to see which possibilities are **excluded** by these conditions.
First condition

\[(PL1) \quad i < j < f(i) \implies f(j) < f(i)\]

\((PL1)\) rules out

\[\ldots \quad i \quad j \quad f(i)\]

where \(f(j) \neq f(i)\), and also

\[\ldots \quad f(j)\]

where \(f(i) < f(j)\).
Second condition

\[(PL2) \quad f(i) \not= i < j \not= f(j) \implies f(i) < f(j).\]

Similarly, (PL2) rules out

We write \(\mathcal{P}(n, m)\) for the set of planar maps in \(\text{Inv}(\mathbb{N}(n, m))\).

**Proposition 7**

1. Every planar diagram satisfies the two conditions.
2. Every involution satisfying the two conditions can be drawn as a planar diagram.
Planarity for Points

Rather than proving this directly, it is simpler, and also instructive, to reduce it to a special case. We consider arrows in $\mathcal{D}$ of the special form $I \rightarrow n$. Such arrows consist only of caps. They correspond to points, or states in the terminology of Categorical QM. Since the top row of dots is empty, in this case we have a linear order, and the premise of condition (PL2) can never arise. Hence planarity for such arrows is just the simple condition (PL1) — which can be seen to be equivalent to saying that, if we write a left parenthesis for each left end of a cap, and a right parenthesis for each right end, we get a well-formed string of parentheses. Thus

![Diagram](diagram.png)

corresponds to

$()(()).$
Names and Conames

Now we recall that quite generally, in any pivotal category we have the Hom-Tensor adjunction

$$A \otimes B^* \overset{[f]}{\longrightarrow} I \quad \overset{\sim}{\leftrightarrow} \quad A \overset{f}{\longrightarrow} B \quad \overset{\sim}{\leftrightarrow} \quad I \overset{[f]}{\longrightarrow} A^* \otimes B$$

$$\llbracket f \rrbracket = (1_{A^*} \otimes f) \circ \eta_A : I \to A^* \otimes B \quad \llcorner f \lrcorner = \epsilon_B \circ (f \otimes 1_{B^*}) : A \otimes B^* \to I.$$  

We call $\llbracket f \rrbracket$ the name of $f$, and $\llcorner f \lrcorner$ the coname. The inverse to the map $f \mapsto \llbracket f \rrbracket$ is defined by

$$g : I \to A^* \otimes B \quad \mapsto \quad (\epsilon_A \otimes 1_B) \circ (1_A \otimes g) : A \to B.$$
Example

We compute the name of the left wave:

\[ \begin{align*}
\text{Applying the inverse transformation:} & \\
\end{align*} \]

Note also that the unit is the name of the identity: \( \eta_n = \lceil 1_n \rceil \), and similarly \( \epsilon_n = \lfloor 1_n \rfloor \).
Composition

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The Temperley-Lieb Category

Our aim is now to define a category $\mathcal{T}$, which will yield the desired description of the Temperley-Lieb monoids. The objects of $\mathcal{T}$ are the natural numbers. The homset $\mathcal{T}(n, m)$ is defined to be the cartesian product $\mathbb{N} \times \mathcal{P}(n, m)$. Thus a morphism $n \to m$ in $\mathcal{T}$ consists of a pair $(s, f)$, where $s$ is a natural number, and $f \in \mathcal{P}(n, m)$ is a planar map in $\text{Inv}(\mathcal{N}(n, m))$.

It remains to define the composition and identities in this category. Clearly (even leaving aside the natural number components of morphisms) composition cannot be defined as ordinary function composition. This does not even make sense — the codomain of a morphism $f : n \to m$ does not match the domain of a morphism $g : m \to p$ — let alone yield a function with the necessary properties to be a morphism in the category.
Consider a map \( f : [n] + [m] \rightarrow [n] + [m] \). Each input lies in either \([n]\) or \([m]\) (exclusive or), and similarly for the corresponding output. This leads to a decomposition of \( f \) into four disjoint partial maps:

\[
\begin{align*}
    f_{n,n} : [n] &\rightarrow [n] \\
    f_{n,m} : [n] &\rightarrow [m] \\
    f_{m,n} : [m] &\rightarrow [n] \\
    f_{m,m} : [m] &\rightarrow [m]
\end{align*}
\]

so that \( f \) can be recovered as the disjoint union of these four maps. If \( f \) is an involution, then these maps will be partial involutions.

Now suppose we have maps \( f : [n] + [m] \rightarrow [n] + [m] \) and \( g : [m] + [p] \rightarrow [m] + [p] \). We write the decompositions of \( f \) and \( g \) as above in matrix form:

\[
\begin{pmatrix}
    f_{n,n} & f_{n,m} \\
    f_{m,n} & f_{m,m}
\end{pmatrix}
\quad
\begin{pmatrix}
    g_{m,m} & g_{m,p} \\
    g_{p,m} & g_{p,p}
\end{pmatrix}
\]
The ‘Execution Formula’

We can view these maps as **binary relations** on \([n] + [m]\) and \([m] + [p]\) respectively, and use relational algebra (union \(R \cup S\), relational composition \(R; S\) and reflexive transitive closure \(R^*\)) to define a **new relation** \(\theta\) on \([n] + [p]\). If we write

\[
\theta = \begin{pmatrix}
\theta_{n,n} & \theta_{n,p} \\
\theta_{p,n} & \theta_{p,p}
\end{pmatrix}
\]

so that \(\theta\) is the disjoint union of these four components, then we can define it component-wise as follows:

\[
\begin{align*}
\theta_{n,n} &= f_{n,n} \cup f_{n,m}; g_{m,m}; (f_{m,m}; g_{m,m})^*; f_{m,n} \\
\theta_{n,p} &= f_{n,m}; (g_{m,m}; f_{m,m})^*; g_{m,p} \\
\theta_{p,n} &= g_{p,m}; (f_{m,m}; g_{m,m})^*; f_{m,n} \\
\theta_{p,p} &= g_{p,p} \cup g_{p,m}; f_{m,m}; (g_{m,m}; f_{m,m})^*; g_{m,p}.
\end{align*}
\]
Reading the Execution Formula

We can give clear intuitive readings for how these formulas express composition of paths in diagrams in terms of relational algebra:

- The component \( \theta_{n,n} \) describes the **cups** of the diagram resulting from the composition. These are the union of the cups of \( f \left( f_{n,n} \right) \), together with paths that start from the top row with a through line of \( f \), given by \( f_{n,m} \), then go through an alternating odd-length sequence of cups of \( g \left( g_{m,m} \right) \) and caps of \( f \left( f_{m,m} \right) \), and finally return to the top row by a through line of \( f \left( f_{m,n} \right) \).

* \( f_{n,m}; g_{m,m}; (f_{m,m}; g_{m,m})^*; f_{m,n} \)
Reading the Execution Formula

- Similarly, $\theta_{p,p}$ describes the caps of the composition.
- $\theta_{n,p} = \theta^c_{p,n}$ describe the through lines. Thus $\theta_{n,p}$ describes paths which start with a through line of $f$ from $n$ to $m$, continue with an alternating even-length (and possibly empty) sequence of cups of $g$ and caps of $f$, and finish with a through line of $g$ from $m$ to $p$.

\[
\begin{align*}
&f_{n,m}; (g_{m,m}; f_{m,m})^*; g_{m,p} \\
&\text{All through lines from } n \text{ to } p \text{ must have this form.}
\end{align*}
\]

**Proposition 8** If $f$ and $g$ are planar, so is $\theta$. 
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Cloning
An Example

We shall consider the **bracketing combinator**

\[ B \equiv \lambda x.\lambda y.\lambda z. x(yz) : (B \rightarrow C) \rightarrow (A \rightarrow B) \rightarrow (A \rightarrow C). \]

This is characterized by the equation \( Babc = a(bc) \).
An Example

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This is characterized by the equation \( Babc = a(bc) \).

We take \( A = B = C = 1 \) in TL. The interpretation of the open term

\[ x : B \to C, y : A \to B, z : A \vdash x(yz) : C \]

is as follows:

Here \( x^+ \) is the output of \( x \), and \( x^- \) the input, and similarly for \( y \). The output of the whole expression is \( o \).
Example ctd

When we abstract the variables, we obtain the following caps-only diagram:
Example ctd

When we abstract the variables, we obtain the following caps-only diagram:

Now we consider an application $B_{abc}$:
An Example

We shall consider the **commuting combinator**

\[ C \equiv \lambda x. \lambda y. \lambda z. xzy : (A \rightarrow B \rightarrow C) \rightarrow B \rightarrow A \rightarrow C. \]

This is characterized by the equation \( C_{abc} = acb \).
An Example

We shall consider the **commuting combinator**

\[ C \equiv \lambda x.\lambda y.\lambda z. xzy : (A \to B \to C) \to B \to A \to C. \]

This is characterized by the equation \( Cabc = acb \).

The interpretation of the open term

\[ x : A \to B \to C, y : B, z : A \vdash xzy : C \]

is as follows:

Here \( x^+ \) is the output of \( x \), \( x^1 \) the first input, and \( x^2 \) the second input. The output of the whole expression is \( o \).
Example ctd

When we abstract the variables, we obtain the following caps-only diagram:
Example ctd

When we abstract the variables, we obtain the following caps-only diagram:

Now we consider an application $C_{abc}$:
Wider perspective
The **Brauer algebra** (1931) arises if we drop the planarity condition on the TL algebra. This plays an important role in the representation theory of the Orthogonal group (‘Schur-Weyl duality’). A whole genre of ‘diagram algebras’ in Representation Theory.
Wider perspective

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- With **BCI** combinators one can interpret **Linear λ-calculus**.
Wider perspective

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- With **BCI** combinators one can interpret **Linear λ-calculus**.

- One can retrieve the Kelly-Laplaza construction of the free compact closed category by a straightforward generalization.
Logic of Quantum Information Flow

- Quantum Entanglement
- From ‘paradox’ to ‘feature’: Teleportation
- What is the output?
- Follow the line!
- Graphical Calculus for Information Flow
- Projectors
- Decomposed
- Compositionality
- Compositionality ctd
- Teleportation
Quantum Entanglement

Bell state: $|00\rangle + |11\rangle$

EPR state: $|01\rangle + |10\rangle$
Quantum Entanglement

Bell state:

\[ |00\rangle + |11\rangle \]

EPR state:

\[ |01\rangle + |10\rangle \]

Compound systems are represented by tensor product: \( \mathcal{H}_1 \otimes \mathcal{H}_2 \). Typical element:

\[ \sum_i \lambda_i \cdot \phi_i \otimes \psi_i \]

Superposition encodes correlation. Einstein’s ‘spooky action at a distance’. Even if the particles are spatially separated, measuring one has an effect on the state of the other. Bell’s theorem: QM is essentially non-local.
From ‘paradox’ to ‘feature’: Teleportation

\[ |\phi\rangle \]

\[ x \in \mathbb{B}^2 \]

\[ U_x \]

\[ M_{\text{Bell}} \]

\[ |00\rangle + |11\rangle \]
What is the output?

\[ (\mathcal{P}_f \otimes 1) \circ (1 \otimes \mathcal{P}_f) \circ \mathcal{P}_f : \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \mathcal{H}_3 \rightarrow \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \mathcal{H}_3 \]
What is the output?

\[(P_{f_4} \otimes 1) \circ (1 \otimes P_{f_3}) \circ (P_{f_2} \otimes 1) \circ (1 \otimes P_{f_1}) : \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \mathcal{H}_3 \longrightarrow \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \mathcal{H}_3\]

\[\phi_{\text{out}} = f_3 \circ f_4 \circ f_2^\dagger \circ f_3^\dagger \circ f_1 \circ f_2(\phi_{\text{in}})\]
Follow the line!

\[ f_3 \circ f_4 \circ f_2^\dagger \circ f_3^\dagger \circ f_1 \circ f_2 \]
Compact Closure: The basic algebraic laws for units and counits.

\[(\epsilon_A \otimes 1_A) \circ (1_A \otimes \eta_A) = 1_A\]

\[(1_{A^*} \otimes \epsilon_A) \circ (\eta_A \otimes 1_{A^*}) = 1_{A^*}\]
Projectors Decomposed

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- Compositionality ctd
- Teleportation
**Compositionality**

The key algebraic fact from which teleportation (and many other protocols) can be derived.
Compositionality ctd

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Compositionality

Compositionality ctd

Teleportation
Teleportation diagrammatically

\[ \beta_i \beta_i^{-1} = \beta_i \]
It’s Logic!

The graphical calculus can be seen as a calculus of proofs for a certain logic — which is highly non-classical, (in particular resource-sensitive, so e.g. it builds in ‘No Cloning’), but also very different from the Birkhoff-von Neumann quantum logic.

Simplification of diagrams — ‘straightening out the lines’ — corresponds to normalization or cut-elimination of proofs.
Cloning

- Cloning vs. Copy-cat
- Some References
Cloning vs. Copy-cat

Copy-cat is **linear copying**: swapping $A \leftrightarrow A$, rather than **cloning** $A \rightarrow A, A$. 
Cloning vs. Copy-cat

Copy-cat is **linear copying**: swapping $A \leftrightarrow A$, rather than **cloning** $A \rightarrow A, A$.

- In **Logic**, Cloning corresponds to the Contraction rule

\[
\frac{\Gamma, A, A \vdash B}{\Gamma, A \vdash B}
\]

and takes us from Linear to Classical Logic.
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- In **Computation**, Cloning allows us to define combinators such as

\[
W xy = xyy
\]

and takes us from linear-time to universal computational power.
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Copy-cat is **linear copying**: swapping \( A \leftrightarrow A \), rather than **cloning** \( A \rightarrow A, A \).

- **In Logic**, Cloning corresponds to the Contraction rule

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  \Gamma, A \vdash B
  \]

  and takes us from Linear to Classical Logic.
- **In Computation**, Cloning allows us to define combinators such as

  \[
  W_{xy} = xyy
  \]

  and takes us from linear-time to universal computational power.
- **In Physics**, Cloning can be used to express the passage from quantum to classical: *given the choice of a basis*, we can define a linear map

  \[
  \mathcal{H} \rightarrow \mathcal{H} \otimes \mathcal{H}.
  \]
Some References

Papers available from my webpages
http://web.comlab.ox.ac.uk/oucl/work/samson.abramsky/