

Aspects of duality in 2-categories

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Introduction

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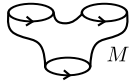


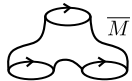
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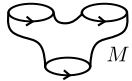
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Stupid is as stupid does. — Forrest Gump's mum.

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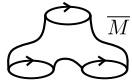
	nCob	Hilb
objects	 $(n-1)$ -dim space	H fin dim Hilbert space
morphisms	 n -dim spacetime	H_1 $\downarrow A$ H_2 linear map
monoidal		$H \otimes H$
duals for objects		\bar{H}
duals for morphisms		H_2 $\downarrow A^*$ H_1 adjoint

Introduction

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objects	 $(n-1)$ -dim space	H fin dim Hilbert space
morphisms	 n -dim spacetime	$ \begin{array}{c} H_1 \\ \downarrow A \\ H_2 \end{array} $ linear map

For instance, quantum entanglement:

$$\begin{array}{ccc}
 \text{Diagram of a pair of pants} & \neq & \text{Diagram of two separate circles} \\
 \downarrow \psi & \neq & \downarrow \psi_1 + \downarrow \psi_2 \\
 H \otimes H & & H + H
 \end{array}$$

duals for morphisms		$ \begin{array}{c} \dots \\ \downarrow A^* \\ H_1 \end{array} $ adjoint
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Baez and Dolan proposed that a unitary extended n -dimensional TQFT is a *unitary representation* of the cobordism n -category on the n -category of n -Hilbert spaces:

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So understanding duality in higher categories will aid our understanding of spacetime and quantum theory!

» Skip extra stuff

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$$Z(\text{pt}) = 2\text{Rep}(G).$$

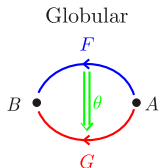
Indeed, at least when G is finite:

Theorem (BB, SW).

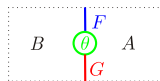
$$\underbrace{\text{Bun}_G(G)}_{Z(S^1)} \overset{\text{braided mon}}{\cong} \underbrace{\text{Dim } 2\text{Rep}(G)}_{\text{category of weak transformations and modifications of identity 2-functor}}.$$

2. String diagram notation for 2-categories

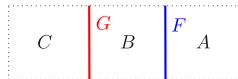
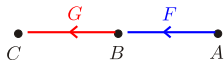
Objects, morphisms
and 2-morphisms



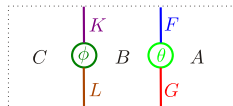
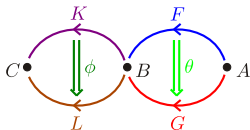
String



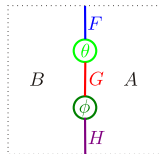
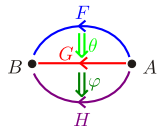
Composition
of 1-morphisms



Horizontal composition
of 2-morphisms

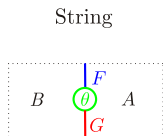
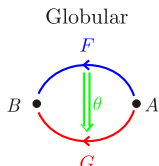


Vertical composition
of 2-morphisms

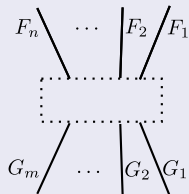


2. String diagram notation for 2-categories

Objects, morphisms
and 2-morphisms



Note — string diagrams work perfectly well for *weak* 2-categories, as long as the the parenthesis scheme of the input and output 1-morphisms are understood in each context, eg.:



unique interpretation
 \mapsto
by coherence

$$[F_n \circ (F_{n-1} \circ F_{n-2})] \circ \cdots \circ [F_2 \circ F_1]$$



$$G_m \circ \cdots \circ [G_3 \circ (G_2 \circ G_1)]$$

\overleftarrow{H}



Ambijunction groupoid I

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An *ambidextrous adjoint* of a morphism $F: A \rightarrow B$ in a 2-category is a morphism $F^*: B \rightarrow A$ equipped with unit and counit 2-morphisms expressing F^* as a right adjoint of F , and unit and counit 2-morphisms expressing F^* as a left adjoint of F :

$$\langle F^* \rangle \equiv \left(\begin{array}{c} \downarrow \\ F^* \\ \uparrow \end{array}, \begin{array}{c} \uparrow \\ F^* \\ \downarrow \\ F \end{array}, \begin{array}{c} \downarrow \\ F \\ \uparrow \\ F^* \end{array}, \begin{array}{c} \downarrow \\ F \\ \uparrow \\ F^* \end{array}, \begin{array}{c} \downarrow \\ F^* \\ \uparrow \\ F \end{array} \right)$$

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$$\langle F^* \rangle \equiv \left(\downarrow^{F^*}, F^* \curvearrowright F, F \curvearrowleft F^*, F \curvearrowright F^*, F^* \curvearrowleft F \right)$$

Taken together these form the *ambijunction groupoid* $\text{Amb}(F)$ of F , where a morphism $\gamma: \langle F^* \rangle \rightarrow \langle (F^*)' \rangle$ is defined to be an invertible 2-morphism

$$\begin{array}{c} \downarrow^{F^*} \\ \circlearrowleft \gamma \\ \downarrow^{(F^*)'} \end{array}$$

such that

$$\left(\downarrow^{F^*}, \circlearrowleft \gamma, \downarrow^{(F^*)'}, \circlearrowright \gamma^{-1}, \downarrow^{F^*}, \circlearrowright \gamma, \downarrow^{(F^*)'}, \circlearrowleft \gamma^{-1} \right) = \left(\downarrow^{F^*}, \curvearrowright F, F \curvearrowleft F^*, F \curvearrowright F^*, F^* \curvearrowleft F \right).$$

Ambijunction groupoid II

We write $[\text{Amb}(F)]$ for the isomorphism classes in $\text{Amb}(F)$.

Properties of the ambijunction groupoid

If $\text{Amb}(F)$ is not empty, then

- There is at most one arrow between any two ambidextrous adjunctions in $\text{Amb}(F)$.
- The group $\text{Aut}(F)$ of automorphisms of F acts freely and transitively on $[\text{Amb}(F)]$.

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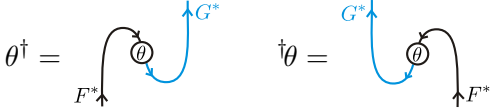
An automorphism $\alpha: F \rightarrow F$ acts on an ambidextrous adjoint by twisting the right unit and counit maps,

$$\alpha \cdot [F^*] = \left[\begin{array}{c} \uparrow \\ F^* \\ \downarrow \end{array}, \begin{array}{c} \text{arc} \\ \uparrow \downarrow \\ F^* \quad F \end{array}, \begin{array}{c} \text{arc} \\ \downarrow \uparrow \\ F \quad F^* \end{array}, \begin{array}{c} \text{arc} \\ \uparrow \downarrow \\ F \quad F^* \\ \circlearrowleft \alpha \end{array}, \begin{array}{c} \text{arc} \\ \downarrow \uparrow \\ F^* \quad F \\ \circlearrowright \alpha^{-1} \end{array} \right].$$

Even-handed structures

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Suppose $\theta: F \Rightarrow G$ is a 2-morphism, and that choices of ambidextrous adjoints $\langle F^* \rangle$ and $\langle G^* \rangle$ have been made. The *right and left daggers* of θ are:



Whether $\theta^\dagger = {}^\dagger\theta$ only depends on $[F^*]$ and $[G^*]$.

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$$\theta^\dagger = \begin{array}{c} \curvearrowright \\ \theta \\ \curvearrowleft \\ F^* \end{array} \quad \begin{array}{c} \curvearrowright \\ G^* \\ \curvearrowleft \\ \theta \\ \curvearrowright \\ F^* \end{array}$$

Whether $\theta^\dagger = {}^\dagger\theta$ only depends on $[F^*]$ and $[G^*]$.

We say that a 2-category \mathcal{C} has *ambidextrous adjoints* if every morphism has an (unspecified) ambidextrous adjoint.

Even-handed structures

Suppose $\theta: F \Rightarrow G$ is a 2-morphism, and that choices of ambidextrous adjoints $\langle F^* \rangle$ and $\langle G^* \rangle$ have been made. The *right and left daggers* of θ are:

The diagram shows two equations for the right and left daggers of a 2-morphism θ . On the left, $\theta^\dagger =$ is followed by a diagram where a blue arrow labeled F^* enters from the bottom left, curves around a circle containing θ , and exits at the top left. A second blue arrow labeled G^* enters from the bottom right, curves around the same circle, and exits at the top right. On the right, $\dagger\theta =$ is followed by a similar diagram where the blue arrow labeled F^* enters from the bottom right and the blue arrow labeled G^* enters from the bottom left.

Whether $\theta^\dagger = \dagger\theta$ only depends on $[F^*]$ and $[G^*]$.

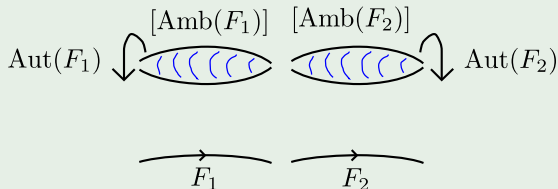
We say that a 2-category \mathcal{C} has *ambidextrous adjoints* if every morphism has an (unspecified) ambidextrous adjoint.

Definition

An *even-handed structure* on a 2-category with ambidextrous adjoints is a choice $F^{[*]} \in [\text{Amb}(F)]$ for every morphism F , such that:

- 1 $\text{id}^{[*]} =$ class of trivial ambidextrous adjunction for all identity 1-cells,
- 2 $(G \circ F)^{[*]} = F^{[*]} \circ G^{[*]}$ for all composable pairs of morphisms, and
- 3 $\theta^\dagger = \dagger\theta$ for every 2-morphism $\theta: F \Rightarrow G$, provided they are computed using ambidextrous adjoints from the classes $F^{[*]}$ and $G^{[*]}$.

So... an even-handed structure on a 2-category is an 'even-handed trivialization of the *ambijunction gerbe*':



Analagous to Murray and Singer's reformulation of a spin structure on a manifold as a trivialization of the *spin gerbe*.

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Lemma

For each choice of system of right duals \star on C , there is a canonical bijection

$$\left\{ \begin{array}{l} \text{Even-handed structures on } C, \\ \text{considered as a one object 2-category} \end{array} \right\} \cong \left\{ \begin{array}{l} \text{Pivotal structures on} \\ C \text{ with respect to } \star \end{array} \right\}.$$

Moreover, this bijection is natural with respect to changing the system of right duals \star .

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- In good cases, an even-handed structure $[*]$ gives rise to a ring homomorphism

$$\dim_{[*]}: K(C) \rightarrow \text{End}(1)$$

$$[V] \mapsto V^* \begin{array}{c} \circlearrowleft \\ \circlearrowright \end{array} V.$$

If C is a semisimple linear category, $[*]$ is characterized by $\dim_{[*]}$.

Examples: Monoidal categories (one-object 2-categories)

- ① For the monoidal category (G, ω) coming from a 3-cocycle $\omega \in Z^3(G, U(1))$,

$$\left\{ \begin{array}{c} \text{Even-handed structures} \\ \text{on } (G, \omega) \end{array} \right\} \cong \left\{ \begin{array}{c} \text{Group homomorphisms} \\ f: G \rightarrow U(1) \end{array} \right\}.$$

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$$\left\{ \begin{array}{c} \text{Even-handed structures} \\ \text{on } C \end{array} \right\} \cong \left\{ \begin{array}{c} \text{Pretwists on } C \\ \text{with respect to } \sigma \end{array} \right\}.$$

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$\omega \in \text{Aut}(\text{id})$ is a *pretwist* with respect to σ if

$$\begin{array}{c} V \\ | \\ | \\ | \\ \bigcirc \theta \\ | \\ | \\ W \end{array} = \begin{array}{c} V \\ \diagdown \\ \bigcirc \theta \\ \diagup \\ W \end{array} \begin{array}{c} W \\ \diagup \\ \bigcirc \theta \\ \diagdown \\ V \end{array} .$$

- Such a

ns }.

right

Gives even-handed structure

$$V[*] = \left[\begin{array}{c} \uparrow V^*, V^* \uparrow \downarrow V, V \downarrow \uparrow V^* \\ \uparrow V^* \uparrow \downarrow V, V \downarrow \uparrow V^* \\ \bigcirc \theta \\ \uparrow V^* \uparrow \downarrow V, V \downarrow \uparrow V^* \\ \bigcirc \theta^{-1} \end{array} \right] .$$

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- ① For the monoidal category (G, ω) coming from a 3-cocycle $\omega \in Z^3(G, U(1))$,

$$\left\{ \begin{array}{c} \text{Even-handed structures} \\ \text{on } (G, \omega) \end{array} \right\} \cong \left\{ \begin{array}{c} \text{Group homomorphisms} \\ f: G \rightarrow U(1) \end{array} \right\}.$$

- ② Suppose \mathcal{C} is a monoidal category where every object has a right dual. If \mathcal{C} can be equipped with a braiding σ , then

$$\left\{ \begin{array}{c} \text{Even-handed structures} \\ \text{on } \mathcal{C} \end{array} \right\} \cong \left\{ \begin{array}{c} \text{Pretwists on } \mathcal{C} \\ \text{with respect to } \sigma \end{array} \right\}.$$

- ③ For *fusion categories* (semisimple rigid monoidal k -linear categories), some mysteries remain:

- What are the equations which have the symmetry group

$$\text{Aut}_{\otimes}(\text{id}) = \{\theta_i \in k^\times : \theta_i = \theta_j \theta_k \text{ whenever } X_i \text{ appears in } X_j \otimes X_k\}?$$

- Will they ensure that an even-handed structure always exists?

Examples: Monoidal categories (one-object 2-categories)

The manifest invariants of a fusion category are the fusion ring $K(C)$ and Müger's *squared norms* d_i of the simple objects:

$$d_i = X_i^* \circlearrowleft X_i \quad X_i \circlearrowright X_i^* \quad \dots \text{this is independent of choice of } \langle X_i^* \rangle.$$

If an even-handed structure exists, then $\dim_{[*]}: K(C) \rightarrow k$ satisfies:

- It is a ring homomorphism taking nonzero values on simple objects,
- $\dim[X_i] \dim[X_i^*] = d_i$ for all simple objects.

Do *these* equations have $\text{Aut}_{\otimes}(\text{id})$ as their symmetry group? To make an even-handed structure, crucially need:

$$X_k^* \circlearrowleft X_j^* \circlearrowright X_i \quad \begin{matrix} \circlearrowright a^p \\ \circlearrowleft a_q \end{matrix} = \delta_q^p X_i \circlearrowright X_i^*$$

Even-handed structures on sub-2-categories of $\mathcal{C}at$

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If $F: A \rightarrow B$ is a functor between categories, express right and left adjunctions via

$$\phi: \text{Hom}(Fx, y) \xrightarrow{\cong} \text{Hom}(x, F^*y)$$

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For $\theta: F \Rightarrow G$, can express $\theta^\dagger, \dagger\theta: G^* \Rightarrow F^*$ as follows:

$$\text{post}(\theta_y^\dagger) = \phi_F \circ \text{pre}(\theta_x) \circ \phi_G^{-1}$$

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In this way, an even-handed structure on a sub-2-category $\mathcal{C} \subseteq \mathcal{C}at$ with ambidextrous adjoints translates into a system of bijective maps

$$\Psi_{F, F^*}: \text{Adj}(F \dashv F^*) \rightarrow \text{Adj}(F^* \dashv F)$$

which transform correctly under isomorphism 2-cells, respect composition, and satisfy the *even-handed equation*:

$$\text{post}^{-1}(\phi_F \circ \text{pre}(\theta) \circ \phi_G^{-1}) = \text{pre}^{-1}(\Psi(\phi_G)^{-1} \circ \text{post}(\theta) \circ \Psi(\phi_F))$$

Even-handed structures from traces on linear categories

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Definition

A *trace* on a k -linear category is a linear map $\text{Tr}_x: \text{End}(x) \rightarrow k$ for each object satisfying:

- $\text{Tr}_x(gf) = \text{Tr}_y(fg)$.
- Nondegeneracy: $s: \text{Hom}(x, y) \xrightarrow{\cong} \text{Hom}(y, x)^\vee$.

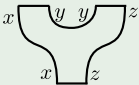

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Remark

A trace on C is the same thing as a symmetric monoidal functor

$\text{PlanarCob}_{\text{Ob}(C)}$	\rightarrow	Vect_k
$x \text{---} y$	\mapsto	$\text{Hom}(x, y)$
	\mapsto	$\text{Hom}(x, y) \otimes \text{Hom}(y, z)$
		$\downarrow \circ$
		$\text{Hom}(x, z)$
$x \text{---} x$	\mapsto	$\text{Hom}(x, x)$
		$\downarrow \text{Tr}_x$
		k
	\mapsto	$\frac{1}{\sum_i f_i \otimes f^i}$ etc.

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Turn right adjoints into left adjoints:

$$\begin{array}{ccc} \text{Hom}(F^*y, x) & \xrightarrow{\phi^T} & \text{Hom}(y, Fx) \\ s_A \downarrow & & \downarrow s_B \\ \text{Hom}(x, F^*y)^\vee & \xrightarrow{\phi^\vee} & \text{Hom}(Fx, y)^\vee \end{array} \quad \phi^T := s_B^{-1} \circ \phi^\vee \circ s_A$$

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Theorem

If each category in $\mathcal{C} \subseteq \text{LinearCat}_k$ comes equipped with a trace, then sending $\phi \mapsto \phi^T$ gives an even handed structure on \mathcal{C} .

Example: The 2-category of 2-Hilbert spaces

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A 2-Hilbert space is an abelian Hilb-category H equipped with antilinear maps $*$: $\text{Hom}(x, y) \rightarrow \text{Hom}(y, x)$ compatible with composition and inner products. They form a 2-category $2\mathcal{H}ilb$.

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The resulting even-handed structure on $2\mathcal{Hilb}$ can be expressed as

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where ϕ^ is the adjoint linear map of ϕ .*

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where ϕ^* is the adjoint linear map of ϕ .

Indeed, can prove that in the semisimple context,

$$\left\{ \begin{array}{l} \text{Even-handed structures on} \\ \text{a full sub-2-category} \\ \mathcal{S} \subset \mathcal{SLCat}_k \end{array} \right\} \cong \left\{ \begin{array}{l} \text{Weightings on the} \\ \text{simple objects in } \mathcal{S} \\ \text{up to a global scale factor} \end{array} \right\}.$$

This is similar to

$\dim\{\text{harmonic spinors on a spin manifold}\}$ is a conformal invariant.

- Every 2-Hilbert space comes equipped with a natural even-handed structure $f \mapsto (\text{id}_x, f)$.

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X an n -dimensional Calabi-Yau manifold. Have the *graded derived category* $\mathbf{D}(X)$:

- An object is a bounded complex \mathcal{E} of coherent sheaves on X .
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Corollary

The 2-category \mathcal{CYau} comes equipped with a canonical even-handed structure arising from Serre duality on each $\mathbf{D}(X)$.

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 - goes hand in hand with a powerful string-diagram calculus
- The concept of an even-handed structure, and the ‘moduli space’ of even-handed structures on a given 2-category, have geometric overtones.
- Personal motivation: needed an even-handed structure on $2\mathcal{Hilb}$ to make the 2-character into a *functor*

$$\chi: \underbrace{2\mathrm{Rep}(G)}_{\text{homotopy category}} \rightarrow \mathrm{Bun}_G(G)$$

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