

A Layered Aggregate Engine for Analytics Workloads

`fdbresearch.github.io`

`relational.ai`



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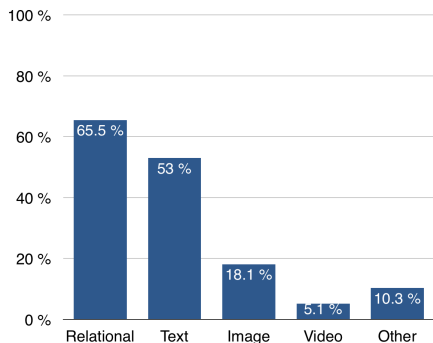
relationalAI
AI for the enterprise

ACM SIGMOD

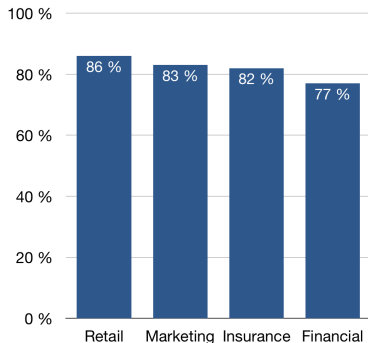
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Relational Data is Ubiquitous

Kaggle Survey: Most Data Scientists use Relational Data at Work!



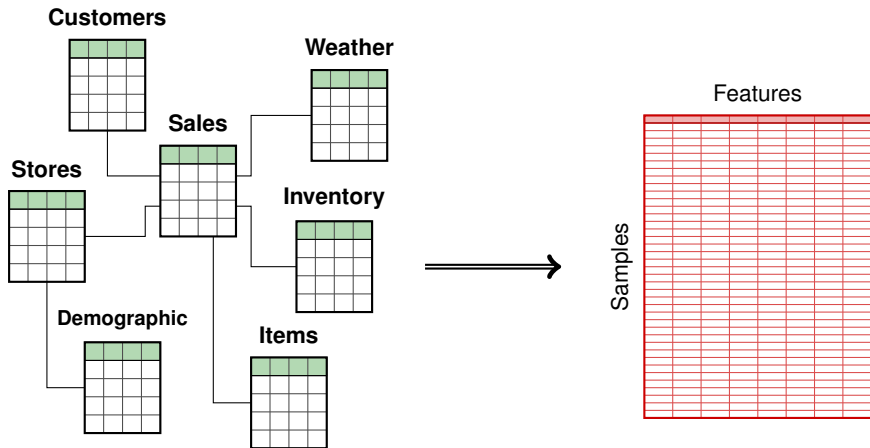
Overall



By Industry

Source: The State of Data Science & Machine Learning 2017, Kaggle, October 2017
(based on 2017 Kaggle survey of 16,000 ML practitioners)

Current State of Affairs in Analytics Workloads



- Carefully crafted by domain experts
- Comes with relational structure
- Throws away relational structure
- Can be order-of-magnitude larger

Turn Analytics Workload into Database Workload!

Many analytics workloads require computation of

batches of aggregate queries.

Advantages:

1. Use DB tools for optimization
2. Decompose Aggregates into views over join tree
 - ▶ Using different roots and directional views
 - ▶ Pushing aggregate computation past joins
3. Avoid materialization of data matrix

Challenge:

1. Workloads require **many** aggregate queries

In contrast:

1. Many ML systems rely on Linear Algebra packages for optimizations

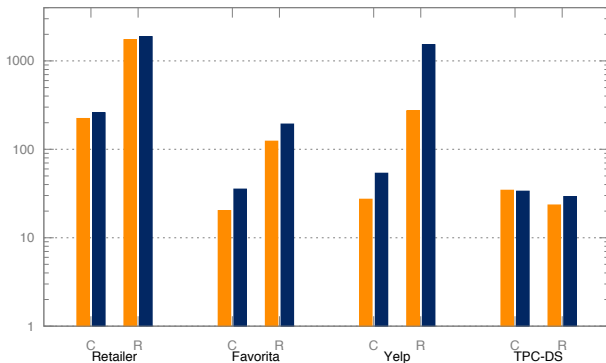
Aggregates are at the Core of Analytics Workloads

| Workload | Query Batch | # Queries |
|---------------------------------------|---|-----------|
| Linear Regression | $SUM(X_i * X_j)$ | 814 |
| Covariance Matrix | $SUM(X_i) \text{ GROUP BY } X_j$ $COUNT(*) \text{ GROUP BY } X_i, X_j$ | |
| Decision Tree (Regression, 1 Node) | $VARIANCE(Y) \text{ WHERE } X_j = c_j$ | 3,141 |
| Mutual Information | $COUNT(*) \text{ GROUP BY } X_i$ | 56 |
| Chow-Liu Trees | $COUNT(*) \text{ GROUP BY } X_i, X_j$ | |
| Data Cubes | $SUM(M) \text{ GROUP BY } X_1, \dots, X_d$ | 40 |

(# Queries shown for Retailer dataset)

Existing DBMSs are **NOT** Designed for Query Batches

Relative Speedup for **Our Approach** over **DBX** and **MonetDB**



C = Covariance Matrix; R = Regression Tree Node; AWS d2.xlarge (4 vCPUs, 32GB)

Tools of a Database Researcher

1. Exploit structure in the data

- ▶ Algebraic structure: Factorized aggregate computation
- ▶ Combinatorial structure: Query complexity measures

2. Sharing computation and data access

- ▶ Aggregates decomposed into views over join tree
- ▶ Share data access across views

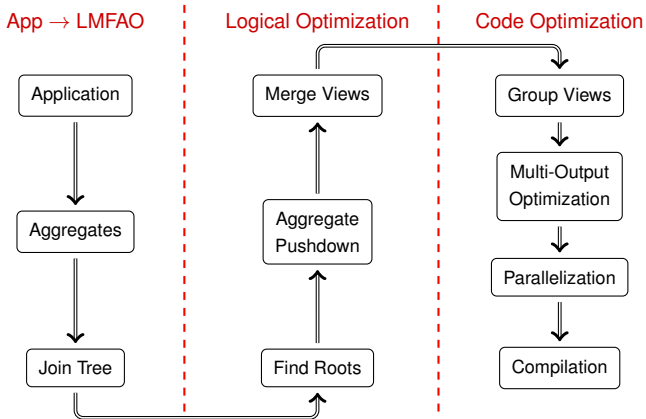
3. Specialization for workload and data

- ▶ Generate code specific to the query batch and dataset
- ▶ Improve cache locality for hot data

4. Parallelization

- ▶ Task and domain parallelism

LMFAO: Layered Multi Functional Aggregate Optimization

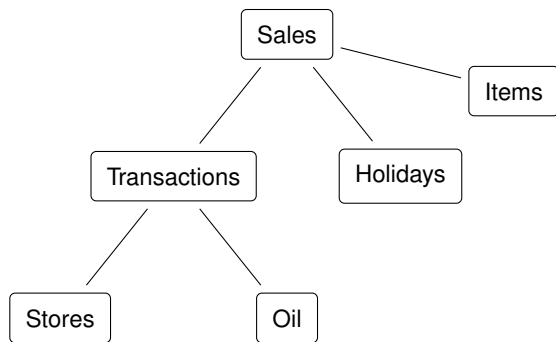


The Layers of LMFAO: Logical Optimization

Q_1 : $\text{SUM}(f(\text{units}))$

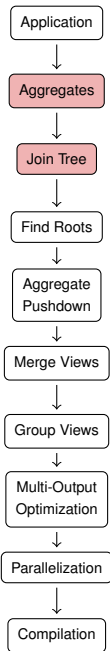
Q_2 : $\text{SUM}(g(\text{item}) \cdot h(\text{date}, \text{color}))$ GROUP BY store

Q_3 : $\text{SUM}(f(\text{units}) \cdot g(\text{item}))$ GROUP BY color



Favorita Kaggle Dataset:

Units Sales for different store, date, item.

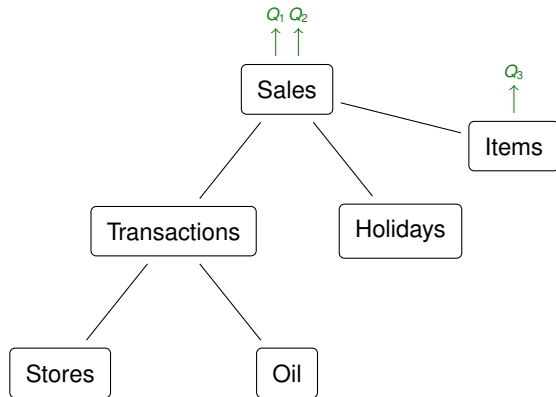


The Layers of LMFAO: Logical Optimization

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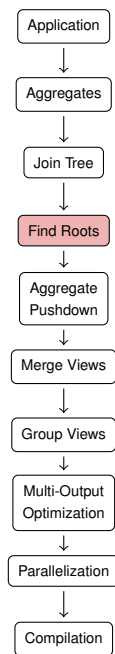
Q_3 : $\text{SUM}(f(\text{units}) \cdot g(\text{item}))$ GROUP BY color



Find Roots Layer:

For each query, decide its output (root) node.

Choose root which minimizes sizes of views.

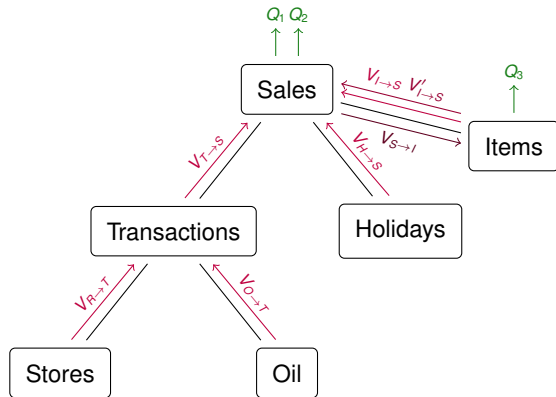


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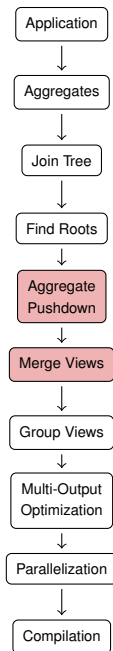
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Aggregate Pushdown Layer:

Break down each query into *directional views* over the join tree.

Reuse Partial Aggregates & **Merge Views** with same group-by attributes.

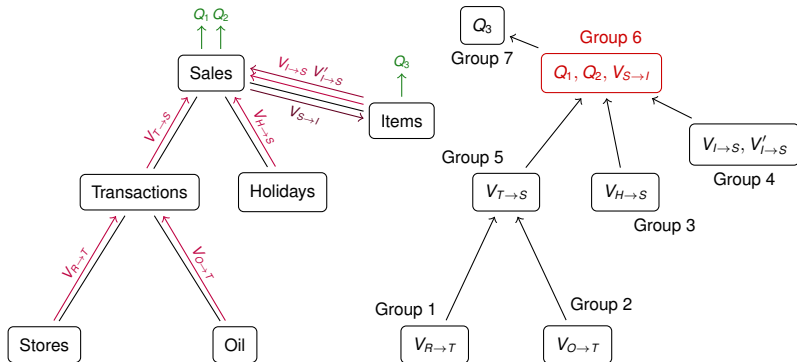


The Layers of LMFAO: Code Optimization

Q_1 : $\text{SUM}(f(\text{units}))$

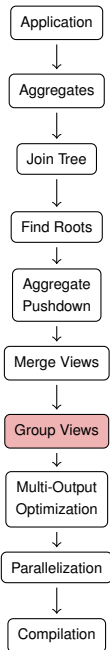
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Group Views Layer:

1. Construct Dependency Graph,
2. Group Views that are computed over same relation.

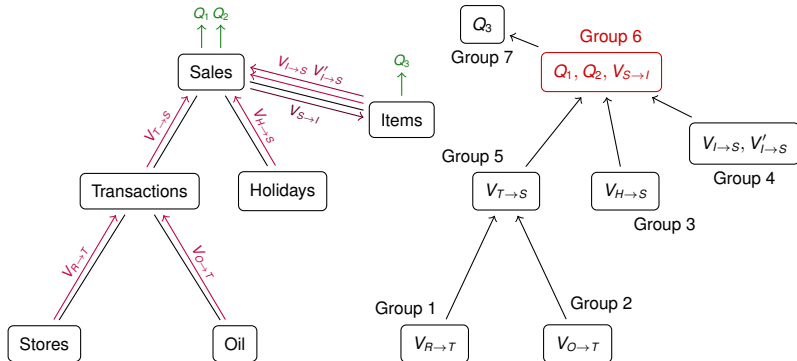


The Layers of LMFAO: Code Optimization

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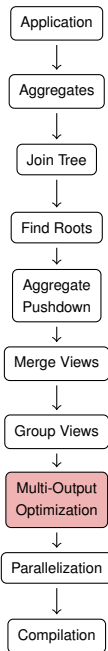
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Multi-Output Optimization Layer:

View Group is a **computational unit** in LMFAO.

All views in one group are computed in one scan over the relation.

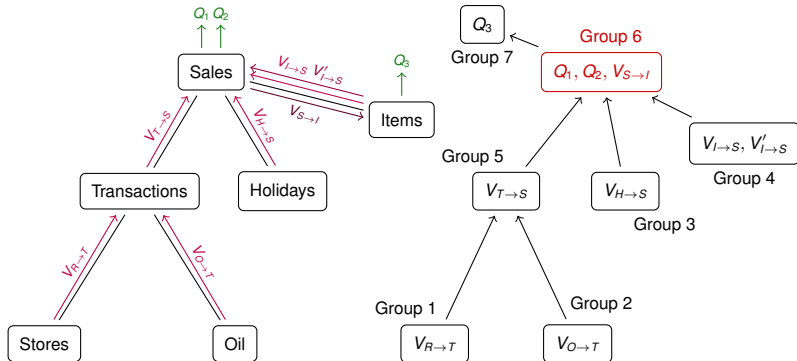


The Layers of LMFAO: Code Optimization

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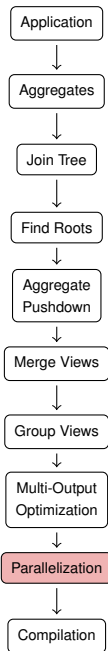
Q_3 : $\text{SUM}(f(\text{units}) \cdot g(\text{item}))$ GROUP BY color



Parallelization Layer:

Task parallelism: Evaluate independent groups in parallel

Domain parallelism: Partition the large relation used by each group

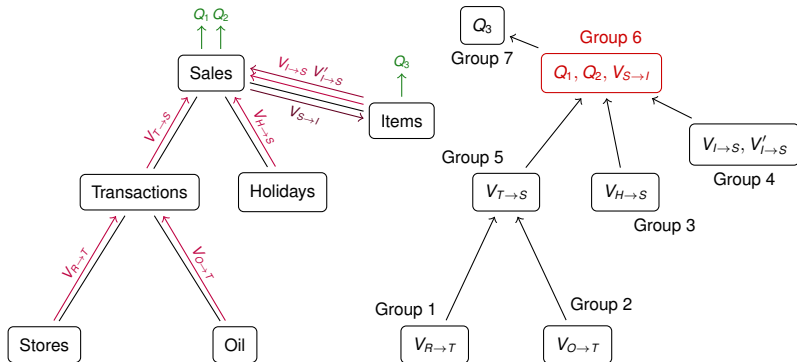


The Layers of LMFAO: Code Optimization

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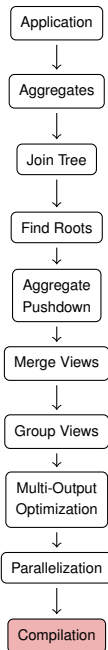
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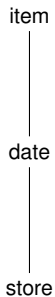


Compilation Layer:

Generate C++ code to compute each View Group.



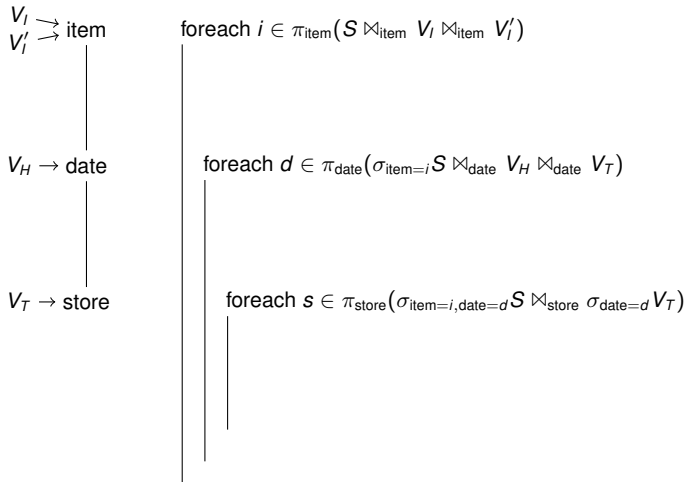
Code Generation for Executing View Group 6 over Sales



Q_1 : SUM ($f(\text{units})$)

Traverse Sales as a trie following an order of its join attributes

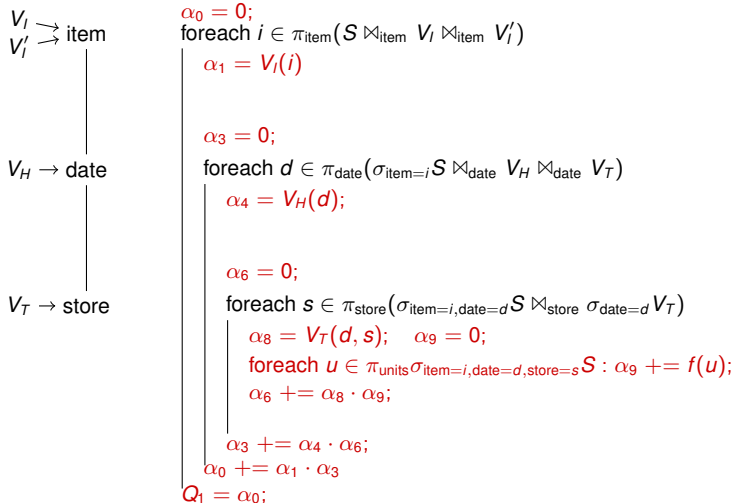
Code Generation for Executing View Group 6 over Sales



Q_1 : SUM ($f(\text{units})$)

Lookup into incoming views, e.g., V_H , as early as possible

Code Generation for Executing View Group 6 over Sales



Q_1 : SUM ($f(units)$)

Insert code for partial aggregates as early as possible

Reduces number of executed instructions

Code Generation for Executing View Group 6 over Sales

$V_I \rightarrow$
 $V'_I \rightarrow$ item

$V_H \rightarrow$ date

$V_T \rightarrow$ store

```
 $\alpha_0 = 0;$ 
foreach  $i \in \pi_{\text{item}}(\mathcal{S} \bowtie_{\text{item}} V_I \bowtie_{\text{item}} V'_I)$ 
   $\alpha_1 = V_I(i)$ 
   $\alpha_2 = g(i);$ 
   $\alpha_3 = 0;$ 
  foreach  $d \in \pi_{\text{date}}(\sigma_{\text{item}=i} \mathcal{S} \bowtie_{\text{date}} V_H \bowtie_{\text{date}} V_T)$ 
     $\alpha_4 = V_H(d);$ 
     $\alpha_6 = 0;$ 
    foreach  $s \in \pi_{\text{store}}(\sigma_{\text{item}=i, \text{date}=d} \mathcal{S} \bowtie_{\text{store}} \sigma_{\text{date}=d} V_T)$ 
       $\alpha_8 = V_T(d, s); \quad \alpha_9 = 0;$ 
      foreach  $u \in \pi_{\text{units}} \sigma_{\text{item}=i, \text{date}=d, \text{store}=s} \mathcal{S} : \alpha_9 += f(u);$ 
       $\alpha_6 += \alpha_8 \cdot \alpha_9;$ 
     $\alpha_3 += \alpha_4 \cdot \alpha_6;$ 
   $\alpha_0 += \alpha_1 \cdot \alpha_3 \quad V_{S \rightarrow I}(i) = \alpha_3 \cdot \alpha_2;$ 
 $Q_1 = \alpha_0;$ 
```

$V_{S \rightarrow I}$: SUM ($f(\text{units}) \cdot g(\text{item})$) GROUP BY item

Different outputs share partial aggregates

Code Generation for Executing View Group 6 over Sales

$V_I \rightarrow$
 $V'_I \rightarrow$ item

$V_H \rightarrow$ date

$V_T \rightarrow$ store

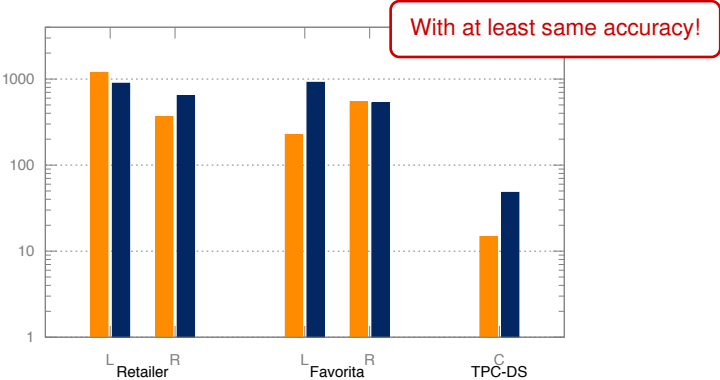
```
 $\alpha_0 = 0;$ 
foreach  $i \in \pi_{\text{item}}(\mathcal{S} \bowtie_{\text{item}} V_I \bowtie_{\text{item}} V'_I)$ 
   $\alpha_1 = V_I(i)$ 
   $\alpha_2 = g(i);$ 
   $\alpha_3 = 0;$ 
  foreach  $d \in \pi_{\text{date}}(\sigma_{\text{item}=i} \mathcal{S} \bowtie_{\text{date}} V_H \bowtie_{\text{date}} V_T)$ 
     $\alpha_4 = V_H(d); \quad \alpha_5 = 0;$ 
    foreach  $c \in \pi_{\text{color}} \sigma_{\text{item}=i} V'_I : \alpha_5 += h(d, c) \cdot V'_I(i, c);$ 
     $\alpha_6 = 0; \quad \alpha_7 = \alpha_5 \cdot \alpha_2 \cdot \alpha_4;$ 
    foreach  $s \in \pi_{\text{store}}(\sigma_{\text{item}=i, \text{date}=d} \mathcal{S} \bowtie_{\text{store}} \sigma_{\text{date}=d} V_T)$ 
       $\alpha_8 = V_T(d, s); \quad \alpha_9 = 0; \quad \alpha_{10} = |\sigma_{\text{item}=i, \text{date}=d, \text{store}=s} \mathcal{S}|;$ 
      foreach  $u \in \pi_{\text{units}} \sigma_{\text{item}=i, \text{date}=d, \text{store}=s} \mathcal{S} : \alpha_9 += f(u);$ 
       $\alpha_6 += \alpha_8 \cdot \alpha_9; \quad \alpha_{11} = \alpha_7 \cdot \alpha_8 \cdot \alpha_{10};$ 
      if  $Q_2(s)$  then  $Q_2(s) += \alpha_{11}$  else  $Q_2(s) = \alpha_{11};$ 
       $\alpha_3 += \alpha_4 \cdot \alpha_6;$ 
     $\alpha_0 += \alpha_1 \cdot \alpha_3 \quad V_{S \rightarrow I}(i) = \alpha_3 \cdot \alpha_2;$ 
   $Q_1 = \alpha_0;$ 
```

Q_2 : SUM ($g(\text{item}) \cdot h(\text{date}, \text{color})$) GROUP BY store

Different outputs share partial aggregates

Experimental Evaluation

Relative Speedup for **LMFAO** over **TensorFlow** and **MADlib**



L = Linear Regression; R = Regression Tree; C = Classification Tree;
TensorFlow learns only 1 Decision Tree Node. Intel i7-4770 (8 CPUs, 32GB)