On the Hardness of Robust Classification

P. Gourdeau, V. Kanade, M. Kwiatkowska and J. Worrell

University of Oxford

Overview

A computational and information-theoretic study of the hardness of robust learning.

A computational and information-theoretic study of the hardness of robust learning.

Setting: Binary classification tasks on input space $\mathcal{X} = \{0, 1\}^n$ in the presence of an adversary.

A computational and information-theoretic study of the hardness of robust learning.

Setting: Binary classification tasks on input space $\mathcal{X} = \{0, 1\}^n$ in the presence of an adversary.

E.g.: distinguishing between handwritten 0's and 1's:

111111111111111

A computational and information-theoretic study of the hardness of robust learning.

Setting: Binary classification tasks on input space $\mathcal{X} = \{0, 1\}^n$ in the presence of an adversary.

E.g.: distinguishing between handwritten 0's and 1's:

 $\{((0,1,\ldots,1),0),((1,1,\ldots,1),1),\ldots,((0,1,\ldots,0),0)\}$

A computational and information-theoretic study of the hardness of robust learning.

Setting: Binary classification tasks on input space $\mathcal{X} = \{0, 1\}^n$ in the presence of an adversary.

E.g.: distinguishing between handwritten 0's and 1's:

 $\{((0,1,\ldots,1),0),((1,1,\ldots,1),1),\ldots,((0,1,\ldots,0),0)\}$



A computational and information-theoretic study of the hardness of robust learning.

Setting: Binary classification tasks on input space $\mathcal{X} = \{0, 1\}^n$ in the presence of an adversary.

E.g.: distinguishing between handwritten 0's and 1's:

 $\{((0,1,\ldots,1),0),((1,1,\ldots,1),1),\ldots,((0,1,\ldots,0),0)\}$



• A comparison of different notions of robust risk,

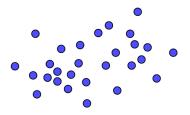
- A comparison of different notions of robust risk,
- A result on the impossibility of sample-efficient *distribution-free* robust learning,

- A comparison of different notions of robust risk,
- A result on the impossibility of sample-efficient *distribution-free* robust learning,
- *Robustness thresholds* to robustly learn monotone conjunctions under log-Lipschitz distributions,

- A comparison of different notions of robust risk,
- A result on the impossibility of sample-efficient *distribution-free* robust learning,
- *Robustness thresholds* to robustly learn monotone conjunctions under log-Lipschitz distributions,
- A simple proof of the computational hardness of robust learning.

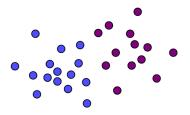
Big picture:

Big picture:



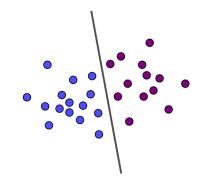
Data i.i.d. from unknown distribution

Big picture:



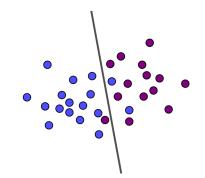
Data i.i.d. from unknown distribution labelled from some concept.

Big picture:



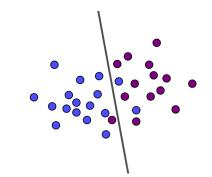
Data i.i.d. from unknown distribution labelled from some concept. We focus on the *realizable setting*,

Big picture:



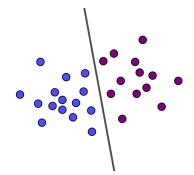
Data i.i.d. from unknown distribution labelled from some concept. We focus on the *realizable setting*, as opposed to the *agnostic setting*.

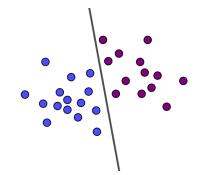
Big picture:

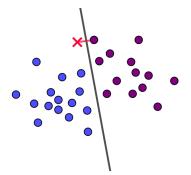


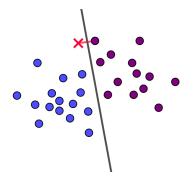
Data i.i.d. from unknown distribution labelled from some concept. We focus on the *realizable setting*, as opposed to the *agnostic setting*.

Learning algorithm A with sample complexity m: when given a sample S of size $\geq m$, A outputs a hypothesis that has low error w.h.p. over S.

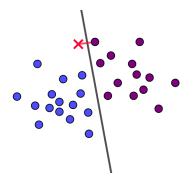




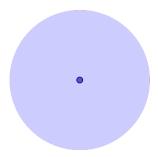


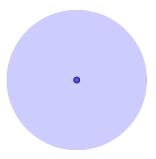


Goal: learn a function that will be robust (with high probability) against an adversary who can perturb the test data.

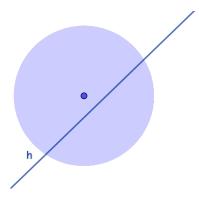


Goal: learn a function that will be robust (with high probability) against an adversary who can perturb the test data. **Question:** How do we define a misclassification?

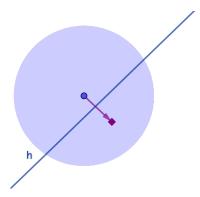




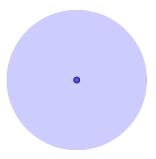
- c: target concept
- h: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)



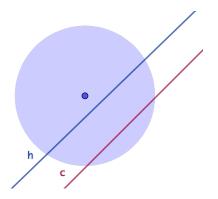
- c: target concept
- h: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)



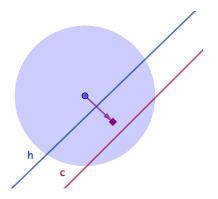
- c: target concept
- h: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)



- c: target concept
- h: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)



- c: target concept
- h: hypothesis
- ρ : robustness parameter (adversary's perturbation budget)



- c: target concept
- h: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)

Robust Risk Definitions

- c: target concept
- *h*: hypothesis
- ρ: robustness parameter (adversary's perturbation budget)

Robust Risk Definitions

- c: target concept
- h: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)

Robust risks:

Constant-in-the-ball: probability that an adversary can perturb a point x drawn from D to z with budget ρ , so that c on x and h on z differ:

$$\mathsf{R}^{\mathsf{C}}_{\rho}\left(h, \ c\right) = \mathop{\mathbb{P}}_{x \sim D}\left(\exists z \in B_{\rho}\left(x\right) \, . \ c(x) \neq h(z)\right) \ .$$

Robust Risk Definitions

- c: target concept
- *h*: hypothesis
- *ρ*: robustness parameter (adversary's perturbation budget)

Robust risks:

Constant-in-the-ball: probability that an adversary can perturb a point x drawn from D to z with budget ρ , so that c on x and h on z differ:

$$\mathsf{R}^{C}_{\rho}\left(h, \ c\right) = \mathop{\mathbb{P}}_{x \sim D}\left(\exists z \in B_{\rho}\left(x\right) \, . \ c(x) \neq h(z)\right)$$

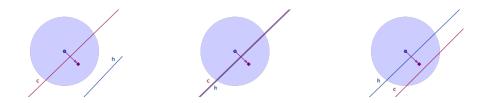
Exact-in-the-ball: probability that an adversary can perturb a point x drawn from D to z with budget ρ , so that c and h disagree on z:

$$\mathsf{R}^{E}_{
ho}\left(oldsymbol{h}, oldsymbol{c}
ight) = \mathop{\mathbb{P}}_{x\sim D}\left(\exists z\in B_{
ho}\left(x
ight) \,.\, oldsymbol{c}(z)
eq oldsymbol{h}(z)
ight) \;.$$

In general, the constant-in-the-ball and the exact-in-the-ball risk functions are not comparable:

Comparing Robust Risk Functions

In general, the constant-in-the-ball and the exact-in-the-ball risk functions are not comparable:



(a) $\mathsf{R}^{E}_{\rho} > 0$, $\mathsf{R}^{C}_{\rho} = 0$, (b) $\mathsf{R}^{E}_{\rho} = 0$, $\mathsf{R}^{C}_{\rho} > 0$, (c) $\mathsf{R}^{E}_{\rho} > 0$, $\mathsf{R}^{C}_{\rho} > 0$.

Choosing a Robust Risk Function

- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,

- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,

- R^{C}_{ρ} pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,



- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,
 - some concept classes are *inherently* not robust w.r.t. to this definition,

- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,
 - some concept classes are *inherently* not robust w.r.t. to this definition,
 - as $\rho \rightarrow n$, we require the function to be constant.

- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,
 - some concept classes are *inherently* not robust w.r.t. to this definition,
 - as $\rho \rightarrow n$, we require the function to be constant.

R^{E}_{ρ} pros and cons:

• requires knowledge of *c* outside of sampled points, e.g. through membership queries,

- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,
 - some concept classes are *inherently* not robust w.r.t. to this definition,
 - as $\rho \rightarrow n$, we require the function to be constant.

R^{E}_{ρ} pros and cons:

• requires knowledge of *c* outside of sampled points, e.g. through membership queries,

•
$$\mathsf{R}^{R}_{\rho}(c,c) = 0.$$

- $\mathsf{R}^{\mathsf{C}}_{\rho}$ pros and cons:
 - simple: only need to know x's correct label to evaluate its loss,
 - can have positive risk when c = h,
 - some concept classes are *inherently* not robust w.r.t. to this definition,
 - as $\rho \rightarrow n$, we require the function to be constant.

 R^{E}_{ρ} pros and cons:

• requires knowledge of *c* outside of sampled points, e.g. through membership queries,

•
$$\mathsf{R}^{R}_{\rho}(c,c) = 0.$$

In our view: adversary's power = creating perturbations that cause $c \neq h$, so we choose R_{ρ}^{E} , despite its drawbacks.

Efficient Robust Learning

 $\mathcal{A} \text{ efficiently } \rho \text{-robustly learns a concept class } \mathcal{C} \text{ with respect to} \\ \text{ distribution class } \mathcal{D}:$

There exists a polynomial sample complexity function poly such that

 for any input dimension n, any target concept c, any distribution D, and any accuracy and confidence parameters ε, δ > 0,

There exists a polynomial sample complexity function poly such that

- for any input dimension n, any target concept c, any distribution D, and any accuracy and confidence parameters ε, δ > 0,
- when \mathcal{A} is given access to a sample $S \sim D^m$, where $m \ge \text{poly}(1/\epsilon, 1/\delta, n)$, \mathcal{A} outputs $h : \{0, 1\}^n \to \{0, 1\}$ such that

There exists a polynomial sample complexity function poly such that

- for any input dimension n, any target concept c, any distribution D, and any accuracy and confidence parameters ε, δ > 0,
- when \mathcal{A} is given access to a sample $S \sim D^m$, where $m \ge \text{poly}(1/\epsilon, 1/\delta, n)$, \mathcal{A} outputs $h : \{0, 1\}^n \to \{0, 1\}$ such that

There exists a polynomial sample complexity function poly such that

- for any input dimension n, any target concept c, any distribution D, and any accuracy and confidence parameters ε, δ > 0,
- when A is given access to a sample $S \sim D^m$, where $m \ge poly(1/\epsilon, 1/\delta, n)$, A outputs $h : \{0, 1\}^n \rightarrow \{0, 1\}$ such that

$$\mathop{\mathbb{P}}_{S\sim D^m} \left(\mathsf{R}^{\mathsf{E}}_{
ho(n)}(h,c) < \epsilon
ight) > 1-\delta$$
 .

There exists a polynomial sample complexity function poly such that

- for any input dimension n, any target concept c, any distribution D, and any accuracy and confidence parameters ε, δ > 0,
- when A is given access to a sample $S \sim D^m$, where $m \ge poly(1/\epsilon, 1/\delta, n)$, A outputs $h : \{0, 1\}^n \rightarrow \{0, 1\}$ such that

$$\mathop{\mathbb{P}}_{S\sim D^m} \left(\mathsf{R}^{\mathsf{E}}_{
ho(n)}(h,c) < \epsilon
ight) > 1-\delta$$
 .

Note:

• We require *polynomial* sample complexity,

There exists a polynomial sample complexity function poly such that

- for any input dimension n, any target concept c, any distribution D, and any accuracy and confidence parameters ε, δ > 0,
- when A is given access to a sample $S \sim D^m$, where $m \ge poly(1/\epsilon, 1/\delta, n)$, A outputs $h : \{0, 1\}^n \rightarrow \{0, 1\}$ such that

$$\mathop{\mathbb{P}}_{S\sim D^m} \left(\mathsf{R}^{\mathcal{E}}_{
ho(n)}(h,c) < \epsilon
ight) > 1-\delta$$
 .

Note:

- We require *polynomial* sample complexity,
- It might make more sense to require *finite* sample complexity in other contexts, e.g. ℝⁿ.

Theorem

 $\ensuremath{\mathcal{C}}$ is efficiently distribution-free robustly learnable iff it is trivial.

Theorem

 ${\mathcal C}$ is efficiently distribution-free robustly learnable iff it is trivial.

Proof idea:

• If C is non-trivial, we can find c_1 and c_2 and x such that

$$(0, 0, \dots, 1, \dots, 0, 0)$$

 $c_1(x) = c_2(x)$.

Theorem

 ${\mathcal C}$ is efficiently distribution-free robustly learnable iff it is trivial.

Proof idea:

• If C is non-trivial, we can find c_1 and c_2 and x such that

 $(0, 0, \dots, 0, \dots, 0, 0)$ $c_1(x) \neq c_2(x)$.

Theorem

 ${\mathcal C}$ is efficiently distribution-free robustly learnable iff it is trivial.

Proof idea:

• If C is non-trivial, we can find c_1 and c_2 and x such that

$$(0, 0, \dots, 0, \dots, 0, 0)$$

 $c_1(x) \neq c_2(x)$.

• Construct a distribution such that c_1 and c_2 will likely agree on a sample of size polynomial in *n* but have $R_{\rho}^{E}(c_1, c_2) = \Omega(1)$.

Theorem

 ${\mathcal C}$ is efficiently distribution-free robustly learnable iff it is trivial.

Proof idea:

• If C is non-trivial, we can find c_1 and c_2 and x such that

$$(0, 0, \dots, 0, \dots, 0, 0)$$

 $c_1(x) \neq c_2(x)$.

- Construct a distribution such that c_1 and c_2 will likely agree on a sample of size polynomial in *n* but have $R_{\rho}^{E}(c_1, c_2) = \Omega(1)$.
- Let $c \sim \text{Unif}(c_1, c_2)$ before labelling the sample. Then any function we learn won't be robust against c with positive probability.

"Nice" Distributions

Idea: We need distributional assumptions to have efficient robust learning.

Idea: We need distributional assumptions to have efficient robust learning.

Log-Lipschitz distributions: D is α -log-Lipschitz if the logarithm of the density function is $log(\alpha)$ -Lipschitz w.r.t. the Hamming distance.

Idea: We need distributional assumptions to have efficient robust learning.

Log-Lipschitz distributions: D is α -log-Lipschitz if the logarithm of the density function is $log(\alpha)$ -Lipschitz w.r.t. the Hamming distance.

$$\begin{array}{l} x_1 = (0, \dots, 1, 1, 1, 1, \dots, 0) \\ x_2 = (0, \dots, 1, 0, 1, \dots, 0) \end{array} \implies \frac{D(x_1)}{D(x_2)} \le \alpha \end{array}$$

Idea: We need distributional assumptions to have efficient robust learning.

Log-Lipschitz distributions: D is α -log-Lipschitz if the logarithm of the density function is $log(\alpha)$ -Lipschitz w.r.t. the Hamming distance.

$$\begin{array}{l} x_1 = (0, \dots, 1, 1, 1, 1, \dots, 0) \\ x_2 = (0, \dots, 1, 0, 1, \dots, 0) \end{array} \implies \frac{D(x_1)}{D(x_2)} \le \alpha \end{array}$$

Intuition: input points that are close to each other cannot have vastly different probability masses.

Examples: uniform distribution, product distribution where the mean of each variable is bounded, etc.

Efficient distribution-free robust learning is not possible in general, but what happens when we restrict the class of distributions?

Efficient distribution-free robust learning is not possible in general, but what happens when we restrict the class of distributions? We look at MON-CONJ : monotone conjunctions

• E.g.:
$$h(x) = x_1 \wedge x_3 \wedge x_5$$
.

Efficient distribution-free robust learning is not possible in general, but what happens when we restrict the class of distributions? We look at MON-CONJ : monotone conjunctions

• E.g.:
$$h(x) = x_1 \wedge x_3 \wedge x_5$$

Theorem

The threshold to robustly learn MON-CONJ under log-Lipschitz distributions is $\rho(n) = O(\log n)$.

The threshold to robustly learn MON-CONJ under log-Lipschitz distributions is $\rho(n) = O(\log n)$.

The threshold to robustly learn MON-CONJ under log-Lipschitz distributions is $\rho(n) = O(\log n)$.

To show that MON-CONJ is not efficiently robustly learnable for $\rho(n) = \omega(\log n)$, we can show that, under the *uniform distribution*

• Choose long enough monotone conjunctions c₁ and c₂

The threshold to robustly learn MON-CONJ under log-Lipschitz distributions is $\rho(n) = O(\log n)$.

To show that MON-CONJ is not efficiently robustly learnable for $\rho(n) = \omega(\log n)$, we can show that, under the *uniform distribution*

- Choose long enough monotone conjunctions c₁ and c₂
- Choose input dimension *n* large enough,

The threshold to robustly learn MON-CONJ under log-Lipschitz distributions is $\rho(n) = O(\log n)$.

To show that MON-CONJ is not efficiently robustly learnable for $\rho(n) = \omega(\log n)$, we can show that, under the *uniform distribution*

- Choose long enough monotone conjunctions c₁ and c₂
- Choose input dimension *n* large enough,
- A sample of size polynomial in *n* will likely look *constant* with fixed probability.

The threshold to robustly learn MON-CONJ under log-Lipschitz distributions is $\rho(n) = O(\log n)$.

To show that MON-CONJ is not efficiently robustly learnable for $\rho(n) = \omega(\log n)$, we can show that, under the *uniform distribution*

- Choose long enough monotone conjunctions c₁ and c₂
- Choose input dimension *n* large enough,
- A sample of size polynomial in *n* will likely look *constant* with fixed probability.
- Again, choose target at random before labelling.

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Example:

Input space: $\mathcal{X} = \{0, 1\}^5$ Target: $x_1 \wedge x_3 \wedge x_5$ Hypothesis: $x_1 \wedge x_2 \wedge x_3 \wedge x_4 \wedge x_5$

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Example:

```
Input space: \mathcal{X} = \{0, 1\}^5

Target: x_1 \wedge x_3 \wedge x_5

Hypothesis: x_1 \wedge x_2 \wedge x_3 \wedge x_4 \wedge x_5

Sample:

(1, 1, 1, 0, 1), 1
```

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Example:

```
Input space: \mathcal{X} = \{0, 1\}^5

Target: x_1 \land x_3 \land x_5

Hypothesis: x_1 \land x_2 \land x_3 \land x_5

Sample:

(1, 1, 1, 0, 1), 1
```

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Example:

```
Input space: \mathcal{X} = \{0, 1\}^5

Target: x_1 \land x_3 \land x_5

Hypothesis: x_1 \land x_2 \land x_3 \land x_5

Sample:

(1, 1, 1, 0, 1), 1

(0, 0, 1, 1, 1), 0
```

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Example:

```
Input space: \mathcal{X} = \{0, 1\}^5

Target: x_1 \land x_3 \land x_5

Hypothesis: x_1 \land x_2 \land x_3 \land x_5

Sample:

(1, 1, 1, 0, 1), 1

(0, 0, 1, 1, 1), 0

(1, 0, 1, 1, 1), 1
```

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Algorithm: Start with $h(x) = \bigwedge_{i \in [n]} x_i$. For each positive example x, if $x_i = 0$, remove *i* from the index set.

Example:

```
Input space: \mathcal{X} = \{0, 1\}^5

Target: x_1 \land x_3 \land x_5

Hypothesis: x_1 \land x_3 \land x_5

Sample:

(1, 1, 1, 0, 1), 1

(0, 0, 1, 1, 1), 0

(1, 0, 1, 1, 1), 1
```

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Theorem

The algorithm to PAC-learn MON-CONJ is an efficient ρ -robust learning algorithm for log-Lipschitz distributions when $\rho = O(\log n)$.

Proof idea: Two cases:

- If the target conjunction is short enough, we have learned exactly, and hence robustly.
- If the target conjunction is large enough, we can use concentration bounds to show that the adversary is unlikely to cause a label change.

- Another learning model (statistical query) [Bubeck et al., 2018],
- Cryptographic assumptions [Degwekar and Vaikuntanathan, 2019].

- Another learning model (statistical query) [Bubeck et al., 2018],
- Cryptographic assumptions [Degwekar and Vaikuntanathan, 2019].

Our proof is quite simple, and only relies on the existence of a hard problem on the boolean hypercube in the PAC-learning framework.

- Another learning model (statistical query) [Bubeck et al., 2018],
- Cryptographic assumptions [Degwekar and Vaikuntanathan, 2019].

Our proof is quite simple, and only relies on the existence of a hard problem on the boolean hypercube in the PAC-learning framework.

 $(\mathcal{C},\mathcal{D},\mathcal{X})$

PAC learning

- Another learning model (statistical query) [Bubeck et al., 2018],
- Cryptographic assumptions [Degwekar and Vaikuntanathan, 2019].

Our proof is quite simple, and only relies on the existence of a hard problem on the boolean hypercube in the PAC-learning framework.

$(\mathcal{C},\mathcal{D},\mathcal{X})$	$(\mathcal{C}',\mathcal{D}',\mathcal{X}')$
PAC learning	Robust learning

Pascale Gourdeau (University of Oxford) On the Hardness of Robust Classification

- Another learning model (statistical query) [Bubeck et al., 2018],
- Cryptographic assumptions [Degwekar and Vaikuntanathan, 2019].

Our proof is quite simple, and only relies on the existence of a hard problem on the boolean hypercube in the PAC-learning framework.

$$(\mathcal{C}, \mathcal{D}, \mathcal{X}) \longrightarrow (\mathcal{C}', \mathcal{D}', \mathcal{X}')$$

PAC learning

Robust learning

• The definitions and models come from previous work in adversarial machine learning theory.

- The definitions and models come from previous work in adversarial machine learning theory.
- At first glance, they seem in many ways *natural* and *reasonable*.

- The definitions and models come from previous work in adversarial machine learning theory.
- At first glance, they seem in many ways *natural* and *reasonable*.
 - Their *inadequacies* surface when viewed under the lens of computational learning theory.

- The definitions and models come from previous work in adversarial machine learning theory.
- At first glance, they seem in many ways natural and reasonable.
 - Their *inadequacies* surface when viewed under the lens of computational learning theory.
- It may be possible to only solve "easy" robust learning problems with strong *distributional assumptions*.

- The definitions and models come from previous work in adversarial machine learning theory.
- At first glance, they seem in many ways natural and reasonable.
 - Their *inadequacies* surface when viewed under the lens of computational learning theory.
- It may be possible to only solve "easy" robust learning problems with strong *distributional assumptions*.
- Other learning models, e.g. when one has access to *membership queries*.

• Generalize robustness threshold for other concept classes:

- Generalize robustness threshold for other concept classes:
 - Majority functions,

- Generalize robustness threshold for other concept classes:
 - Majority functions,
 - Linear threshold functions,

- Generalize robustness threshold for other concept classes:
 - Majority functions,
 - Linear threshold functions,
 - etc.

- Generalize robustness threshold for other concept classes:
 - Majority functions,
 - Linear threshold functions,
 - etc.
- More powerful learning model (e.g., membership queries).

Sébastien Bubeck, Eric Price, and Ilya Razenshteyn. *Adversarial examples from computational constraints.* arXiv preprint. arXiv:1805.10204, 2018.

Akshay Degwekar and Vinod Vaikuntanathan. *Computational limitations in robust classification and win-win results.* arXiv preprint. arXiv:1902.01086, 2019.