

Codensity Games for Bisimilarity

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Abstract—Bisimilarity as an equivalence notion of systems has been central to process theory. Due to the recent rise of interest in quantitative systems (probabilistic, weighted, hybrid, etc.), bisimilarity has been extended in various ways: notably, bisimulation metric between probabilistic systems. An important feature of bisimilarity is its game-theoretic characterization, where Spoiler and Duplicator play against each other; extension of bisimilarity games to quantitative settings has been actively pursued too.

In this paper, we present a general framework that uniformly describes game characterizations of bisimilarity-like notions. Our framework is formalized categorically using fibrations and coalgebras. In particular, our characterization of bisimilarity in terms of fibrational predicate transformers allows us to derive codensity bisimilarity games: a general categorical game characterization of bisimilarity. Our framework covers known bisimilarity-like notions (such as bisimulation metric) as well as new ones (including what we call bisimulation topology).

I. INTRODUCTION

A. Bisimilarity Notions and Games

Since the seminal works by Park and Milner [1], [2], *bisimilarity* has played a central role in theoretical computer science. It is an equivalence notion between branching systems; it abstracts away internal states and stresses the black-box observation-oriented view on process semantics. Bisimilarity is usually defined as the largest *bisimulation*, which is a binary relation that satisfies a suitable mimicking condition. In fact, a bisimulation R can be characterized as a post-fixed point $R \subseteq \Phi(R)$ using a suitable relation transformer Φ ; from this we obtain that bisimilarity is the greatest *fixed* point of Φ by the Knaster–Tarski theorem. This order-theoretic foundation is the basis of a variety of advanced techniques for reasoning about (or using) bisimilarity, such as bisimulation up-to—see, e.g., [3].

Bisimilarity is conventionally defined for state-based systems with nondeterministic branching. However, as the applications of computer systems become increasingly pervasive and diverse (such as cyber-physical systems), extension of

bisimilarity to systems with other branching types has been energetically sought in the literature. One notable example is the bisimulation notion for probabilistic systems in [4]: it is a relation that witnesses that two states are indistinguishable in their behaviors henceforth. This qualitative notion has also been made quantitative, as the notion of *bisimulation metric* [5]. It replaces a relation with a metric that is induced by the probabilistic transition structure.

There is a body of literature (including [6]–[12]) that aims to identify the mathematical essences that are shared by this variety of bisimilarity, and express the identified essences in a rigorous manner using *category theory*. Our particular interest is in the correspondence between bisimilarity notions and (*safety*) *games*; three examples of which are given below. This interest in bisimilarity games is shared by the recent work [10], and the comparison is discussed in §I-D.

1) *Bisimilarity Games*: It is well-known that the following game characterizes the conventional notion of bisimilarity between Kripke frames. Let (X, \rightarrow) be a Kripke frame where $\rightarrow \subseteq X^2$; the game is played between Duplicator (D) and Spoiler (S). In a position (x_1, x_2) , Spoiler challenges Duplicator’s claim that x_1 and x_2 are bisimilar, by choosing one of the states (say x_1) and further choosing a transition $x_1 \rightarrow x'_1$. Duplicator responds by choosing a transition $x_2 \rightarrow x'_2$ from the other state, and the game is continued from (x'_1, x'_2) . Duplicator wins if Spoiler gets stuck, or the game continues infinitely long, and this witnesses that x_1 and x_2 are bisimilar.

2) *Games for Probabilistic Bisimilarity*: A recent step forward in the topic of bisimilarity and games is the characterization of probabilistic bisimulation introduced in [13]. For simplicity, here we describe its discrete version.

Let (X, c) be a Markov chain, where X is a countable set of states, and $c: X \rightarrow \mathcal{D}_{\leq 1}X$ is a transition kernel that assigns to each state $x \in X$ a probability subdistribution $c(x) \in \mathcal{D}_{\leq 1}X$. Here $\mathcal{D}_{\leq 1}X = \{d: X \rightarrow [0, 1] \mid \sum_{x \in X} d(x) \leq 1\}$ denotes the set of probability subdistributions over X . For $Z \subseteq X$, let $c(x)(Z)$ denote the probability with which a successor of x is chosen from Z ; that is, $c(x)(Z) = \sum_{x' \in Z} c(x)(x')$. Since $c(x)$ is only a *sub*-distribution over X , the probability $c(x)(X)$ is ≤ 1 rather than $= 1$. The remaining probability $1 - c(x)(X)$ can be thought of as the probability of x getting stuck.

Recall from [4] that an equivalence relation $R \subseteq X^2$ is a (*probabilistic*) *bisimulation* if, for any $(x, y) \in R$ and each R -closed subset $Z \subseteq X$, $c(x)(Z) = c(y)(Z)$ holds.

The game introduced in [13] is in Table I. It is shown in [13]

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TABLE I
THE GAME FOR PROBABILISTIC BISIMILARITY FROM [13]

position	player	possible moves
$(x, y) \in X^2$	S	$Z \subseteq X$ s.t. $c(x)(Z) \neq c(y)(Z)$
$Z \subseteq X$	D	$(x', y') \in X^2$ s.t. $x' \in Z \wedge y' \notin Z$

that Duplicator is winning in the game at (x, y) if and only if x and y are bisimilar, in the sense of [4] (recalled above). It is not hard to find an intuitive correspondence between the game in Table I and the definition of bisimulation [4]: Spoiler challenges the bisimilarity claim between x, y by exhibiting Z such that $c(x)(Z) = c(y)(Z)$ is violated; Duplicator makes a counterargument by claiming that Z is in fact not bisimilarity-closed, exhibiting a pair of states (x', y') that Duplicator claims are bisimilar.

3) *Games for Probabilistic Bisimulation Metric*: Our following observation marked the beginning of the current work: the game for (qualitative) bisimilarity for probabilistic systems (from [13], Table I) can be almost literally adapted to (quantitative) *bisimulation metric* for probabilistic systems. This metric was first introduced in [5].

For simplicity we focus on the discrete setting; we also restrict to pseudometrics bounded by 1. Let (X, c) be a Markov chain with a countable state space X . The *bisimulation metric* $d_{(X,c)}: X^2 \rightarrow [0, 1]$ is defined to be the smallest pseudometric (with respect to the pointwise order) that makes the transition kernel

$$c: (X, d_{(X,c)}) \longrightarrow (\mathcal{D}_{\leq 1}X, \mathcal{K}(d_{(X,c)}))$$

non-expansive with respect to the specified pseudometrics. Here $\mathcal{K}(d_{(X,c)})$ is the so-called *Kantorovich metric* over $\mathcal{D}_{\leq 1}X$ induced by the pseudometric $d_{(X,c)}$ over X . It is defined as follows. For $\mu, \nu \in \mathcal{D}_{\leq 1}X$,

$$\mathcal{K}(d_{(X,c)})(\mu, \nu) = \sup_f |E_\mu[f] - E_\nu[f]|, \quad (1)$$

where in the above sup, f ranges over all non-expansive functions from $(X, d_{(X,c)})$ to $([0, 1], d_{[0,1]})$, $d_{[0,1]}$ denotes the usual Euclidean metric, and $E_\mu[f]$ is the expectation $\sum_{x \in X} f(x) \cdot \mu(x)$ of f with respect to μ .

Our observation is that the bisimulation metric $d_{(X,c)}$ is characterized by the game in Table II: Duplicator is winning at (x, y, ε) if and only if $d_{(X,c)}(x, y) \leq \varepsilon$.

The game seems to be new, although its intuition is similar to the one for Table I. Note that the formula (1) appears in the condition of Spoiler’s moves. Spoiler challenges by exhibiting a “predicate” f that suggests violation of the non-expansiveness of c ; and Duplicator makes a counterargument that f is in fact not non-expansive and thus invalid.

4) *Towards a Unifying Framework*: The last two games (Table I from [13] and Table II that seems new) motivate a general framework that embraces both. There are some clear analogies: the games are about *indistinguishability* of states x, y under a class of *observations* (Z and f respectively), and the *predicates* usable in those observations are subject to

TABLE II
THE GAME FOR (PROBABILISTIC) BISIMULATION METRIC,
ADAPTING [13]

position	P	possible moves
$(x, y, \varepsilon) \in X^2 \times [0, 1]$	S	$f: X \rightarrow [0, 1]$ such that $ E_{c(x)}[f] - E_{c(y)}[f] > \varepsilon$
$f: X \rightarrow [0, 1]$	D	$(x', y', \varepsilon') \in X^2 \times [0, 1]$ such that $ f(x') - f(y') > \varepsilon'$

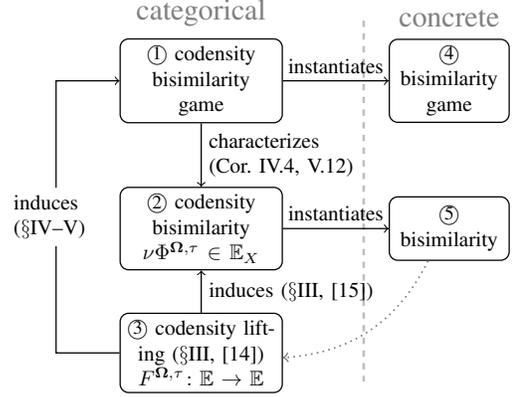


Fig. 1. Our Codensity-Based Framework for Bisimilarity and Games

certain preservation properties (bisimilarity-closedness in the former, and non-expansiveness in the latter).

B. A Codensity-Based Framework for Bisimilarity and Games

The main contribution of the current paper is a categorical framework that derives a variety of bisimilarity notions and corresponding game notions. The correspondence is proved once and for all on the categorical level of generality. It covers the three examples introduced earlier in §I-A, much like the recent categorical framework in [10] does. However, our fibration-based formalization has another dimension of generality. For example, besides relations and metrics, our examples include what we call *bisimulation topology*.

The overview of our categorical framework is in the left half of Fig. 1. We build on our previous works [14] and [15]. In [14] a general construction called *codensity lifting* is introduced: given a fibration $p: \mathbb{E} \rightarrow \mathbb{C}$ and parameters (Ω, τ) that embody the kind of *observations* we can make, a functor $F: \mathbb{C} \rightarrow \mathbb{C}$ is lifted to $F^{\Omega, \tau}: \mathbb{E} \rightarrow \mathbb{E}$. In [15], codensity lifting is used to introduce a generic family of bisimulation notions called *codensity bisimilarity*—see ②. In this paper, we extend these previous results by

- introducing the notion of *codensity bisimilarity game* (①) that comes in two variants (*untrimmed* (§IV) and *trimmed* (§V)),
- establishing the correspondence between codensity bisimulations (②) and games (①) on a fibrational level of generality, and
- working out several concrete examples (④, ⑤).

In general, devising a game notion (④) directly from a bisimilarity notion (⑤) is far from trivial. Indeed, doing so

for an individual bisimilarity notion has itself been deemed a scientific novelty [13], [16]. Our codensity-based framework (in the left half of Fig. 1) can automate *part of* this process in the following precise sense.

We derive concrete notions of bisimilarity (⑤) and bisimilarity game (④) as instances; then the correspondence between the two is guaranteed by the categorical general result between ① and ②.

We note, however, that this is no panacea. When one starts with a given concrete notion of bisimilarity (⑤), their next task would be to identify the right choice of the parameters $\mathbb{E} \xrightarrow{P} \mathbb{C}, \Omega, \tau$ for the codensity lifting (③). This task is not easy in general: we needed to get our hands dirty working out the examples in this paper, [14], and [15]. Nevertheless, we believe that the required passage from ⑤ to ③ is much easier than the direct derivation from ⑤ to ④, with our categorical framework providing templates of bisimilarity games (see Tables VII, IX and X). After all, our framework identifies which part of the path from ⑤ to ④ can be automated, and which part remains to be done individually. This is much like what many other categorical frameworks offer, as meta-level theories.

As an additional benefit, our categorical framework can be used to *discover* new bisimilarity notions (⑤), starting from (choices of parameters for) ③. We believe those derived new bisimilarity notions are useful, since our categorical theory embodies sound intuitions about observation, predicate transformation, and indistinguishability—see e.g. §II-B.

C. Contributions

Our main technical contributions are as follows.

- We introduce a categorical framework that uniformly describes various bisimulation notions (including metrics, preorders and topologies) and the corresponding game notions (Fig. 1). The framework is based on coalgebras, fibrations, and codensity liftings in particular [14]. Our general game notion comes in two variants.
 - The first (the *untrimmed* codensity game: §IV) arises naturally in a fibration, using its objects and arrows as possible moves. The untrimmed game is theoretically clean, but it tends to have a huge arena.
 - We therefore introduce a method that restricts these arenas, leading to the (*trimmed*) codensity bisimilarity game (§V). The reduction method is also described in general fibrational terms, specifically using fibered separators and generating sets.
- From the general framework, we derive several concrete examples of bisimilarity and its related notions (④ and ⑤ in Fig. 1). They are listed in Table VI. Among them, a few bisimilarity notions seem new (especially the *bisimulation topology* examples), and several game notions also seem new.
- We discuss the *transfer of codensity bisimilarity* by suitable fibered functors (§VII). As an example usage, we give an abstract proof of the fact that (usual) bisimilarity for Kripke frames is necessarily an equivalence (Example VII.2).

- Additionally, we conduct investigations of the game notion in [13] (Table I) in concrete, non-categorical terms. For one, we obtain its variation for bisimulation metric (as we showed in Table II). We also give a direct proof of the equivalence to another game notion for probabilistic bisimilarity, previously introduced in [16], by exhibiting a mutual translation of winning strategies. See the extended version [17].

D. Related Work

Besides the one in [13], another game characterization of probabilistic bisimulation has been given in [16]. It is described later in §II (Table III). The latter game has a bigger arena than the one in [13]: in [16] both players have to play a subset $Z \subseteq X$, while in [13] only Spoiler does so.

The work that is the closest to ours is the recent work [10] that studies bisimilarity games in a categorical setting. Their formalization uses (co)algebras (following the (co)algebraic generalization of the Kantorovich metric introduced in [8]), and therefore embraces a variety of different branching types. The major differences between the two works are as follows.

- Our current work is fibration-based (in particular \mathbf{CLat}_{\square} -fibrations), while [10] is not. As a consequence, ours accommodates an additional dimension of generality by changing fibrations, which correspond to different indistinguishability notions (relation, metric, topology, preorder, measurable structures, etc.). In contrast, the works [10] and [8] deal exclusively with two settings: binary relations and pseudometrics.
- A relationship to *modal logic* is beautifully established in [10], while it is not done in this work. We expect our fibrational framework can accommodate modal logic too: fibrations have been used for categorical modeling of logics [18]. We leave this aspect to future work.
- The categorical generalization [10] is based on the game notion in [16], while ours is based on that in [13]. Therefore, for some bisimulation notions (including the bisimulation metric), we obtain a game notion with a smaller arena. Compare Table II (an instance of ours) and Table IV (an instance of [10]).

There are a number of categorical studies of bisimilarity notions; notable mentions include open map-based approaches [19] and coalgebraic ones [20], [21]. The fibrational approach we adopt also uses coalgebras; it was initiated in [6] and pursued, e.g., in [7], [9], [11], and [15]. For example, in the recent work [11], fibrational generality is exploited to study up-to techniques for bisimilarity metric. They use the *Wasserstein lifting* of functors introduced in [8] instead of the codensity lifting that we use (it generalizes the *Kantorovich lifting* in [8], see Example III.4). It is known [8] that the Wasserstein and Kantorovich liftings can differ in general, while they coincide for some specific functors such as the distribution functor.

Some of our new examples are topological: we derive what we call *bisimulation topology* and a game notion that characterizes it. The relation between these notions and the

existing works on bisimulation and topology (including [22], [23]) is left as future work.

E. Organization

In §II, we present preliminaries on a general theory of games (we can restrict to *safety* games), and on fibrations. For the latter, we focus on a class called \mathbf{CLat}_{\top} -fibrations, and argue that they offer an appropriate categorical abstraction of sets equipped with indistinguishability structures. In §III, we present codensity lifting and codensity bisimilarity (②, ③ in Fig. 1). The material is based on [15], but we introduce some auxiliary notions needed for the correspondence with games. Our first game notion (the *untrimmed* one) is introduced in §IV; in §V, we cut down the arenas and obtain *trimmed* codensity bisimilarity game. The theory is further extended in §VI–VII: in §VI we accommodate multiple observation domains, and in §VII we discuss the transfer of codensity bisimilarities by full-faithful fibered functors preserving meets. These categorical observations give rise to the concrete examples in §VIII.

Some proofs and details are deferred to the appendix found in the extended version [17].

II. PRELIMINARIES

We write $\mathcal{P}: \mathbf{Set} \rightarrow \mathbf{Set}$ for the covariant powerset functor, and 2 for the two-point set $2 = \{\perp, \top\}$. We define the function $\diamond: \mathcal{P}2 \rightarrow 2$ called *may-modality* by $\diamond S = \top$ if and only if $\top \in S$. We write Eq_I for the diagonal (equality) relation over a set I .

A. Safety Games

Here we recall some standard game-theoretic notions and results. In capturing bisimilarity-like notions, we can restrict ourselves to *safety games*—they have a simple winning condition where every infinite play is won by the same player (namely Duplicator). This winning condition reflects the characterization of bisimilarity-like notions by suitable *greatest* fixed points; the correspondence generalizes, for example, to the one between parity games and nested alternating fixed points—see [24]. The term “safety game” occurs, e.g., in [25], [26].

Safety games are played between two players; in this paper, they are called *Duplicator* (D) and *Spoiler* (S). We restrict to those games in which Duplicator and Spoiler alternate turns.

Definition II.1 (safety game). A (*safety game*) arena is a triple $\mathcal{G} = (Q_D, Q_S, E)$ of a set Q_D of *Duplicator’s positions*, a set Q_S of *Spoiler’s positions*, and a *transition relation* $E \subseteq (Q_D \times Q_S) \cup (Q_S \times Q_D)$. Hence \mathcal{G} is a bipartite graph. We require that Q_D and Q_S are disjoint, and that $Q_D \cup Q_S \neq \emptyset$. We write $Q = Q_D \cup Q_S$.

For a position $q \in Q$, the elements of the set $\{q' \in Q \mid (q, q') \in E\}$ are called the *possible moves* at q . Unlike some works, we allow positions that have no possible moves at them.

A *play* in an arena $\mathcal{G} = (Q_D, Q_S, E)$ is a (finite or infinite) sequence of positions $q_0 q_1 \dots$, such that $(q_{i-1}, q_i) \in E$ so long as q_i belongs to the sequence.

A play in \mathcal{G} is *won* by either player, according to the following conditions: 1) a finite play $q_0 \dots q_n$ is won by Spoiler (or by Duplicator) if $q_n \in Q_D$ (or $q_n \in Q_S$ respectively); and 2) every infinite play $q_0 q_1 \dots$ is won by Duplicator.

Definition II.2 (strategy, winning position). In an arena $\mathcal{G} = (Q_D, Q_S, E)$, a *strategy* of Duplicator is a partial function $\sigma_D: Q^* \times Q_D \rightarrow Q_S$; we require that $\sigma_D(\vec{q}q) = q'$ implies $(q, q') \in E$. A *strategy* of Spoiler is defined similarly, as a partial function $\sigma_S: Q^* \times Q_S \rightarrow Q_D$ that returns a possible move at the last position in the history.

Given an initial position $q \in Q$ and two strategies σ_D and σ_S for Duplicator and Spoiler respectively, the *play* from q induced by (σ_D, σ_S) is defined in a natural inductive manner. The induced play is denoted by $\pi^{\sigma_D, \sigma_S}(q)$.

A position $q \in Q$ is said to be *winning* for Duplicator if there exists a strategy σ_D of Duplicator such that, for any strategy σ_S of Spoiler, the induced play $\pi^{\sigma_D, \sigma_S}(q)$ is won by Duplicator.

In what follows, for simplicity, we restrict the initial position q of a play $\pi^{\sigma_D, \sigma_S}(q)$ to be in Q_S . (Note that Spoiler’s position can be winning for Duplicator.)

Winning positions of safety games are witnessed by *invariants* (Prop. II.4), which is a well-known fact.

Definition II.3 (invariant). Let $\mathcal{G} = (Q_D, Q_S, E)$ be an arena. A subset $P \subseteq Q_S$ is called an *invariant* for Duplicator if, for each $q \in P$ and any possible move $q' \in Q_D$ at q , there exists a possible move q'' at q' that is in P . That is, $\forall q \in P. \forall q' \in Q_D. ((q, q') \in E \Rightarrow \exists q'' \in Q_S. (q', q'') \in E \wedge q'' \in P)$.

Proposition II.4. 1) *Any position $q \in P$ in an invariant P for Duplicator is winning for Duplicator.*
 2) *Invariants are closed under arbitrary union. Therefore, there exists a largest invariant for Duplicator.*
 3) *The largest invariant for Duplicator coincides with the set of winning positions for Duplicator in Q_S .* \square

Examples of safety games have been given in Tables I–II. We present two other examples (Tables III–IV).

Example II.5 (alternative games for probabilistic bisimilarity and bisimulation metric). In [16], a game notion that characterizes (qualitative) probabilistic bisimilarity is presented. It is in Table III, presented in a slightly adapted form.

This game notion is categorically generalized in [10]; the generalization has freedom in the choice of coalgebra functors (i.e. branching types), as well as in the choice between relations and metrics. The instance of this general game notion for bisimulation metric is shown in Table IV.

The two games (Tables III–IV) characterize the same bisimilarity-like notions as the games in Tables I–II, respectively; so they are equivalent. We can go further and give a direct equivalence proof by mutually translating winning strategies. Such a proof is not totally trivial; we do so for the pair for probabilistic bisimilarity. See [17, Appendix A].

We note that the game in Table II (an instance of our current framework) is simpler than Table IV (an instance of [10]).

Table II is not only structurally simpler (it has fewer rows), but its set of moves are smaller too, asking for functions $X \rightarrow [0, 1]$ only at one place.

TABLE III
THE GAME FOR PROBABILISTIC BISIMILARITY, FROM [16]

We omit labels that are needed to distinguish $(x, y) \in X^2$ (an S-position) from $(s, t) \in X^2$ (a D-position).

position	pl.	possible moves
$(x, y) \in X^2$	S	$(s, t) \in X^2$ s.t. $\{s, t\} = \{x, y\}$
$(s, t) \in X^2$	D	(Z, Z') s.t. $Z \subseteq Z' \subseteq X$ and $c(s)(Z) \leq c(t)(Z')$
$(Z, Z') \in (\mathcal{P}X)^2$	S	$(Z, y') \in \mathcal{P}X \times X$ s.t. $y' \in Z' \setminus Z$
$(Z, y') \in \mathcal{P}X \times X$	D	$(x', y') \in X^2$ s.t. $x' \in Z$

TABLE IV
THE GAME FOR BISIMULATION METRIC, FROM [10]

position	pl.	possible moves
$(x, y, \varepsilon) \in X^2 \times [0, 1]$	S	$(s, t) \in X^2$ s.t. $\{s, t\} = \{x, y\}$, and $f: X \rightarrow [0, 1]$
$(s, t, f, \varepsilon) \in X^2 \times [0, 1]^X \times [0, 1]$	D	$g: X \rightarrow [0, 1]$ such that $\max\{0, E_{c(s)}[f] - E_{c(t)}[g]\} \leq \varepsilon$
$(f, g, \varepsilon) \in ([0, 1]^X)^2$	S	$(i, j) \in ([0, 1]^X)^2$ such that $\{i, j\} = \{f, g\}$, and $x' \in X$
$(x', i, j, \varepsilon) \in X \times ([0, 1]^X)^2 \times [0, 1]$	D	$(x', y', \varepsilon') \in X^2 \times [0, 1]$ such that $i(x') \leq j(y')$, and $\varepsilon' = j(y') - i(x')$

Our categorical framework based on codensity liftings (presented in later sections) covers Tables I–II but not Tables III–IV. Accommodation of the latter two is future work.

B. \mathbf{CLat}_{\sqcap} -fibrations

1) *Definition and Properties:* Here we sketch a basic theory of fibrations—see, e.g., [18] for a comprehensive account. In particular, we focus on a class of poset fibrations called \mathbf{CLat}_{\sqcap} -fibrations. We observe that the simple axiomatics of the class adequately capture all the examples of interest—and hence the mathematical essences of the logical phenomena that we wish to model.

Our exposition here is largely based on that in [15]. However, in this paper we introduce new notation and terminology (such as *indistinguishability order* and *decent map*)—see §II-B2. They help to further clarify the intuitions.

A formal definition is as follows. (See [17, Appendix B] for a rather gentle introduction to \mathbf{CLat}_{\sqcap} -fibration.)

Definition II.6 (\mathbf{CLat}_{\sqcap} -fibration). A \mathbf{CLat}_{\sqcap} -fibration is a fibration $p: \mathbb{E} \rightarrow \mathbb{C}$ such that each fiber \mathbb{E}_X (for each $X \in \mathbb{C}$) is a complete lattice, and each pullback functor $f^*: \mathbb{E}_Y \rightarrow \mathbb{E}_X$ (for each $f: X \rightarrow Y$ in \mathbb{C}) preserves all meets \sqcap .

Via the Grothendieck construction, a \mathbf{CLat}_{\sqcap} -fibration is in a bijective correspondence with a functor $F_{\mathbb{E}}: \mathbb{C}^{\text{op}} \rightarrow \mathbf{CLat}_{\sqcap}$, where \mathbf{CLat}_{\sqcap} is the category of complete lattices and functions preserving all meets—see [18] and [7], as well as [17, Appendix B]. The functor $F_{\mathbb{E}}$ assigns

- a complete lattice \mathbb{E}_X (called the *fiber* over X) to each $X \in \mathbb{C}$, and
- a function $f^*: \mathbb{E}_Y \rightarrow \mathbb{E}_X$ preserving all meets to each $f: X \rightarrow Y$ in \mathbb{C} . The map f^* is called a *pullback*; it is also called a *pullback functor* since, in the general theory of fibrations, a fiber \mathbb{E}_X is a category rather than a poset.

Although the *indexed category* presentation $F_{\mathbb{E}}: \mathbb{C}^{\text{op}} \rightarrow \mathbf{CLat}_{\sqcap}$ may be more intuitive at first, we shall stick to the *fibration* presentation $p: \mathbb{E} \rightarrow \mathbb{C}$ since we will eventually need some global structures in the *total category* \mathbb{E} . It turns out that \mathbf{CLat}_{\sqcap} -fibrations are special kinds of *topological functor* [27] such that each fiber category is a poset. Topological functors are a well-studied topic, and many examples and results are available; a good summary is found in [28].

The use of poset fibrations is common in categorical modeling of logics [7], [9]. \mathbf{CLat}_{\sqcap} -fibrations additionally require fibered small meets; this simple assumption turns out to be a mathematically powerful one.

Proposition II.7. *Let $p: \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration.*

- 1) p is split, and faithful as a functor.
- 2) Each arrow $f: X \rightarrow Y$ has its pushforward $f_*: \mathbb{E}_X \rightarrow \mathbb{E}_Y$, so that an adjunction $f_* \dashv f^*$ is formed. This is a consequence of Freyd's adjoint functor theorem; it makes p a bifibration [18].
- 3) $p^{\text{op}}: \mathbb{E}^{\text{op}} \rightarrow \mathbb{C}^{\text{op}}$ is also a \mathbf{CLat}_{\sqcap} -fibration.
- 4) The change-of-base [18, Lemma 1.5.1] of p along any functor $H: \mathbb{D} \rightarrow \mathbb{C}$ is also a \mathbf{CLat}_{\sqcap} -fibration.
- 5) If \mathbb{C} is (co)complete, then the total category \mathbb{E} is also (co)complete. This follows from [18, Prop. 9.2.1]. \square

2) *Notation, Terminology and Intuitions:* Our view of a \mathbf{CLat}_{\sqcap} -fibration $p: \mathbb{E} \rightarrow \mathbb{C}$ is that it equips objects of \mathbb{C} with what we call *indistinguishability structures*. This suits our purpose, since various bisimilarity-like notions are all about degrees of indistinguishability between (the behaviors of) states of a system. We present examples later in §II-B3.

Notation II.8 (indistinguishability predicate/order). Let $p: \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration. An object $P \in \mathbb{E}_X$ in the fiber category \mathbb{E}_X (i.e. an element of the complete lattice \mathbb{E}_X) is called an *indistinguishability predicate* over X . Our view is that P is an additional structure on X ; therefore, as a convention, an object $P \in \mathbb{E}_X$ shall also be denoted by $(X, P) \in \mathbb{E}_X$.

Each fiber \mathbb{E}_X is a complete lattice; its order is denoted by \sqsubseteq and called the *indistinguishability order* over X . Intuitively, $P \sqsubseteq Q$ means that Q has a greater degree of indistinguishability than P —that is, Q is coarser than P , and P is more discriminating than Q .

The supremum and infimum with respect to the indistinguishability order \sqsubseteq are denoted by \sqcup and \sqcap respectively.

Definition II.9 (decent map). Let $p: \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration, $f: X \rightarrow Y$ be an arrow in \mathbb{C} , $(X, P) \in \mathbb{E}_X$ and $(Y, Q) \in \mathbb{E}_Y$ be objects in the fibers. We say that f is *decent* (from P to Q) if there exists a necessarily unique arrow $\hat{f}: P \rightarrow Q$ in \mathbb{E} such that $p\hat{f} = f$. We write

$f: (X, P) \dot{\rightarrow} (Y, Q)$ in this case. The following equivalences follow.

$$f: (X, P) \dot{\rightarrow} (Y, Q) \iff P \sqsubseteq f^*Q \iff f_*P \sqsubseteq Q$$

We write $f: (X, P) \dot{\rightarrow} (Y, Q)$ if f is *not* decent.

The notion of decency is a fibered generalization of continuity, non-expansiveness, relation-preservation, etc. Decency $f: (X, P) \dot{\rightarrow} (Y, Q)$ means f *respects* indistinguishability, carrying P -indistinguishable elements to Q -indistinguishable ones.

3) *Examples:* As shown in Table V, various well-known categories can be seen as categories that equip sets with certain indistinguishability structures. The evident forgetful functors from the total categories (**Top**, **Meas**, etc.) to **Set** in Table V are all \mathbf{CLat}_{\sqcap} -fibrations.

Specifically, **Top** is the category of topological spaces and continuous maps; **Meas** is that of measurable spaces and measurable maps; \mathbf{PMet}_1 is that of 1-bounded pseudometric spaces (where a *pseudo-metric* is a metric without the condition $d(x, y) = 0 \Rightarrow x = y$) and non-expansive maps; **ERel** is that of sets with endorelations ($X, R \subseteq X^2$) and relation-preserving maps; **Pre** is that of preordered sets and monotone maps; and **EqRel** is that of sets with equivalence relations and relation-preserving maps—see [15] for details.

Note that, in **Top** and **Meas**, the indistinguishability order is the opposite of the inclusion order. Therefore the meet of a family of indistinguishability structures computed as the one generated from the *union* of the family.

Another class of examples is given as follows: for any well-powered category \mathbb{B} admitting small limits, the subobject fibration of \mathbb{B} is a \mathbf{CLat}_{\sqcap} -fibration. All the algebraic categories over **Set** and Grothendieck topoi fall into this class. On the other hand, the forgetful functors from algebraic categories over **Set** are rarely (\mathbf{CLat}_{\sqcap} -)fibrations.

III. CODENSITY BISIMILARITY

We introduce *codensity lifting* (③ in Fig. 1) and *codensity bisimilarity* (②) based on [15]. These turn out to subsume many bisimilarity-like notions in the literature. The material in §III-A–III-B is largely from [15]; §III-C is new, paving the way to codensity bisimilarity games presented in later sections.

A. Codensity Lifting

Definition III.1 (codensity lifting $F^{\Omega, \tau}$ [15]). Let $p: \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration, and $F: \mathbb{C} \rightarrow \mathbb{C}$ be a functor. A *parameter of codensity lifting* of F along p is a pair of

- a \mathbb{C} -arrow $\tau: F\Omega \rightarrow \Omega$ (i.e. an F -algebra) called *modality* [30], [31] and
- an \mathbb{E} -object Ω above Ω called *observation domain*.

The *codensity lifting* of $F: \mathbb{C} \rightarrow \mathbb{C}$ with parameter (Ω, τ) is the endofunctor $F^{\Omega, \tau}: \mathbb{E} \rightarrow \mathbb{E}$ defined as follows. On objects,

$$F^{\Omega, \tau}P = \prod_{k \in \mathbb{E}(P, \Omega)} (\tau \circ F(p(k)))^* \Omega.$$

Its action on arrows is as follows. It is not hard to see that, for each arrow $l: P \rightarrow Q$ in \mathbb{E} , the arrow $F(p(l))$ is decent from

$F^{\Omega, \tau}P$ to $F^{\Omega, \tau}Q$. Then we define $F^{\Omega, \tau}l: F^{\Omega, \tau}P \rightarrow F^{\Omega, \tau}Q$ to be the unique arrow in \mathbb{E} above $F(p(l))$.

An alternative description is possible. When \mathbb{E} has powers \pitchfork and p preserves them, $F^{\Omega, \tau}$ is characterized as the following pullback in the fibration $[\mathbb{E}, p]: [\mathbb{E}, \mathbb{E}] \rightarrow [\mathbb{E}, \mathbb{C}]$:

$$\begin{array}{ccc} [\mathbb{E}, \mathbb{E}] & & F^{\Omega, \tau} \dashrightarrow \mathbb{E}(-, \Omega) \pitchfork \Omega \\ \downarrow [\mathbb{E}, p] & & \\ [\mathbb{E}, \mathbb{C}] & & F \circ p \xrightarrow{\alpha} \mathbb{E}(-, \Omega) \pitchfork \Omega \end{array}$$

where $\alpha_P = \langle \tau \circ F(p(k)) \rangle_{k \in \mathbb{E}(P, \Omega)}$ is the tupling. A similar characterization of codensity liftings of monads is in [14].

Table VI lists concrete examples of codensity liftings, with various fibrations p , functors F , and parameters (Ω, τ) . Some of them coincide with known notions. For example, the entry 5 of the table says that the functor $(\mathcal{D}_{\leq 1})^{\Omega, \tau}$, with the designated Ω and τ , carries a metric space (X, d) to the set $\mathcal{D}_{\leq 1}X$ equipped with the well-known Kantorovich metric $\mathcal{K}(d)$ induced by d —see (1).

Besides the functors listed in the table, there are some natural ways to systematically lift polynomial functors, by defining $\tau: F\Omega \rightarrow \Omega$ in an inductive manner—see, e.g., [11].

Example III.2. Let us closely look at the entry 4 of Table VI. There we codensity-lift the covariant powerset functor \mathcal{P} along the \mathbf{CLat}_{\sqcap} -fibration $\mathbf{EqRel} \rightarrow \mathbf{Set}$. We use the parameter $((2, \text{Eq}_2), \diamond)$, where $\diamond: \mathcal{P}2 \rightarrow 2$ is the modality given in §II.

We shall abbreviate $(2, \text{Eq}_2)$ by Eq_2 —a notational convention that is used throughout the paper.

Then $\mathcal{P}^{\text{Eq}_2, \diamond}(X, R)$ relates $S, T \in \mathcal{P}X$ if and only if

$$\begin{aligned} \forall k: X \rightarrow 2. & \left((\forall (x, y) \in R. k(x) = k(y)) \right) \\ \implies & \left((\exists x \in S. k(x) = \top) \iff (\exists x \in T. k(x) = \top) \right). \end{aligned}$$

Straightforward calculation shows that this is equivalent to

$$(\forall x \in S. \exists y \in T. (x, y) \in R) \wedge (\forall y \in T. \exists x \in S. (x, y) \in R).$$

This lifting is the restriction of the standard relational lifting of \mathcal{P} along $\mathbf{ERel} \rightarrow \mathbf{Set}$, which is used for the usual bisimulation notion for Kripke frames, to \mathbf{EqRel} .

Example III.3. In the entry 3 of Table VI, we codensity-lift \mathcal{P} along the \mathbf{CLat}_{\sqcap} -fibration $\mathbf{ERel} \rightarrow \mathbf{Set}$ (instead of $\mathbf{EqRel} \rightarrow \mathbf{Set}$) with the parameter $((2, \text{Eq}_2), \diamond)$.

The characterization of $\mathcal{P}^{\text{Eq}_2, \diamond}(X, R)$ is slightly involved. Its relation part relates $S, T \in \mathcal{P}X$ if and only if

$$\begin{aligned} (\forall x \in S. \exists y \in T. (x, y) \in R^{\text{eq}}) \wedge \\ (\forall y \in T. \exists x \in S. (x, y) \in R^{\text{eq}}), \end{aligned}$$

where R^{eq} denotes the equivalence closure of R .

It is not clear at this stage whether the codensity bisimilarities induced by the above liftings (Examples III.2–III.3, i.e. the entries 4 & 3 of Table VI) coincide with the usual bisimilarity notion for Kripke frames. This is because of the involvement of mandatory equivalence closures—specifically by the use of \mathbf{EqRel} in Example III.2, and by the occurrence of $(_)^{\text{eq}}$ in Example III.3. Later, in Example VII.2, we prove that both

TABLE V
CLat $_{\square}$ -FIBRATIONS

fibration	indistinguishability structure	decent map	$P \sqsubseteq Q$	$\prod P_i$
Top \rightarrow Set	topology	continuous func.	$P \supseteq Q$	generated from $\bigcup P_i$
Meas \rightarrow Set	σ -field	measurable func.	$P \supseteq Q$	generated from $\bigcup P_i$
PMet $_1$ \rightarrow Set	pseudometric	non-expansive func.	$\forall x, y. P(x, y) \geq Q(x, y)$	$(x, y) \mapsto \sup_i P_i(x, y)$
ERel \rightarrow Set	endorelation	relation preserving func.	$P \subseteq Q$	$\bigcap P_i$
Pre \rightarrow Set	preorder	monotone func.	$P \subseteq Q$	$\bigcap P_i$
EqRel \rightarrow Set	equivalence relation	relation preserving func.	$P \subseteq Q$	$\bigcap P_i$

TABLE VI
CODENSITY LIFTING OF FUNCTORS

	fibration $p : \mathbb{E} \rightarrow \mathbb{C}$	functor $F : \mathbb{C} \rightarrow \mathbb{C}$	obs. dom. Ω	modality τ	lifting $F^{\Omega, \tau}$ of F
1	Pre \rightarrow Set	powerset \mathcal{P}	$(2, \leq)$	$\diamond : \mathcal{P}2 \rightarrow 2$	lower preorder [14]
2	Pre \rightarrow Set	powerset \mathcal{P}	$(2, \geq)$	$\diamond : \mathcal{P}2 \rightarrow 2$	upper preorder [14]
3	ERel \rightarrow Set	powerset \mathcal{P}	$(2, \text{Eq}_2)$	$\diamond : \mathcal{P}2 \rightarrow 2$	(for bisimulation, see Ex. III.3 & VII.2)
4	EqRel \rightarrow Set	powerset \mathcal{P}	$(2, \text{Eq}_2)$	$\diamond : \mathcal{P}2 \rightarrow 2$	(for bisimulation, see Ex. III.2 & VII.2)
5	PMet $_1$ \rightarrow Set	subdistrib. $\mathcal{D}_{\leq 1}$	$([0, 1], d_{[0, 1]})$	$e : \mathcal{D}_{\leq 1}[0, 1] \rightarrow [0, 1]$	Kantorovich metric
6	PMet $_1$ \rightarrow Set	powerset \mathcal{P}	$([0, 1], d_{[0, 1]})$	$\text{inf} : \mathcal{P}[0, 1] \rightarrow [0, 1]$	Hausdorff pseudometric (cf. [17, Appx. C])
7	$U^*(\mathbf{PMet}_1) \rightarrow \mathbf{Meas}$	sub-Giry $\mathcal{G}_{\leq 1}$	$([0, 1], d_{[0, 1]})$	$e : \mathcal{G}_{\leq 1}[0, 1] \rightarrow [0, 1]$	Kantorovich metric
8 \dagger	Pre \rightarrow Set	powerset \mathcal{P}	$(2, \leq), (2, \geq)$	$\diamond : \mathcal{P}2 \rightarrow 2$	convex preorder [14]
9 \dagger	EqRel \rightarrow Set	subdistrib. $\mathcal{D}_{\leq 1}$	$(2, \text{Eq}_2)$	$(\tau_r : \mathcal{D}_{\leq 1}2 \rightarrow 2)_{r \in [0, 1]}$	(for prob. bisim., see §VIII-G)
10 \dagger	Top \rightarrow Set	$2 \times (_)^\Sigma$	Sierpinski space	(see Ex. VI.5)	(for bisim. topology, see Ex. VI.5)

The fibration $U^*(\mathbf{PMet}_1) \rightarrow \mathbf{Meas}$ is obtained as a change-of-base, pulling back $\mathbf{PMet}_1 \rightarrow \mathbf{Set}$ along $U : \mathbf{Meas} \rightarrow \mathbf{Set}$. $d_{[0, 1]}$ denotes the Euclidean metric on the unit interval $[0, 1]$. The modality \diamond is introduced in the beginning of §II. The functions $e : \mathcal{D}_{\leq 1}[0, 1] \rightarrow [0, 1]$ and $e : \mathcal{G}_{\leq 1}[0, 1] \rightarrow [0, 1]$ both return expected values. The lower, upper and convex preorders are known for powerdomains; see e.g. [29]. The function $\tau_r : \mathcal{D}_{\leq 1}2 \rightarrow 2$ is defined by $\tau_r(p) = \top$ if $p(\top) \geq r$, and $\tau_r(p) = \perp$ otherwise.

The examples marked with \dagger involve multiple modalities and observation domains. The extension that allows such is described later in §VI.

of the codensity bisimilarities indeed coincide with the usual bisimilarity notion. The proof relies crucially on transfer of codensity liftings via fibered functors.

Example III.4. Here we follow [15, Example 3] and show that codensity lifting generalizes the *Kantorovich lifting* of functors introduced in [8]. Take $\mathbf{PMet}_1 \rightarrow \mathbf{Set}$ as the \mathbf{CLat}_{\square} -fibration p in Def. III.1. As Ω , we take $\Omega = [0, 1]$ with the usual Euclidean metric $d_{[0, 1]}$. There is freedom in the choice of a modality $\tau : F\Omega \rightarrow \Omega$ —this corresponds to what is called an *evaluation function* in [8]. This way we recover the Kantorovich lifting in [8] as $F^{\Omega, \tau}$.

B. Codensity Bisimilarity

In [15], *codensity bisimulation* and *bisimilarity* are introduced. Recall that a *coalgebra* $c : X \rightarrow FX$ is a categorical presentation of state-based transition systems, such as automata, Markov chains, etc.—see, e.g., [20], [21], and also §VIII.

Definition III.5 (codensity bisimulation). Assume the setting of Def. III.1. Let $c : X \rightarrow FX$ be an F -coalgebra. An object $P \in \mathbb{E}_X$ is a $((\Omega, \tau)$ -codensity bisimulation over c if $c : (X, P) \dot{\rightarrow} (FX, F^{\Omega, \tau}P)$; that is, c is decent with respect to the designated indistinguishability structures on X and FX .

We move on to the characterization of codensity bisimulations as (post-)fixed points of suitable predicate transformers.

Definition III.6 (predicate transformer $\Phi^{\Omega, \tau}$). Assume the setting of Def. III.5. We define a *predicate transformer* $\Phi_c^{\Omega, \tau} : \mathbb{E}_X \rightarrow \mathbb{E}_X$ with respect to c and $F^{\Omega, \tau}$ by:

$\mathbb{E}_X \rightarrow \mathbb{E}_X$ with respect to c and $F^{\Omega, \tau}$ by:

$$\Phi_c^{\Omega, \tau} P = c^*(F^{\Omega, \tau}P), \text{ that is, } \prod_{k \in \mathbb{E}(P, \Omega)} (\tau \circ F(p(k)) \circ c)^* \Omega. \quad (2)$$

Theorem III.7. Assume the setting of Def. III.5. For any $P \in \mathbb{E}_X$, the following are equivalent.

- 1) $c : (X, P) \dot{\rightarrow} (FX, F^{\Omega, \tau}P)$; that is, P is a codensity bisimulation over c (Def. III.5).
- 2) $P \sqsubseteq \Phi_c^{\Omega, \tau}P$.
- 3) For each $k \in \mathbb{C}(X, \Omega)$, $k : (X, P) \dot{\rightarrow} (\Omega, \Omega)$ implies $\tau \circ Fk \circ c : (X, P) \dot{\rightarrow} (\Omega, \Omega)$. \square

The predicate transformer $\Phi_c^{\Omega, \tau}$ is a monotone map from the complete lattice \mathbb{E}_X to itself. Therefore, by the Knaster-Tarski theorem, the greatest post-fixed point of $\Phi_c^{\Omega, \tau}$ exists and it is the greatest fixed point of $\Phi_c^{\Omega, \tau}$.

Definition III.8 (codensity bisimilarity $\nu\Phi_c^{\Omega, \tau}$). Assume the setting of Def. III.5. The greatest codensity bisimulation, whose existence is guaranteed by the above arguments, is called the *codensity bisimilarity*. It is denoted by $\nu\Phi_c^{\Omega, \tau}$.

Some bisimilarity notions, including bisimilarity of deterministic automata (§VIII-B), are accommodated in the generalized framework with multiple observation domains—see §VI.

Example III.9 (bisimulation metric). Consider the \mathbf{CLat}_{\square} -fibration $\mathbf{PMet}_1 \rightarrow \mathbf{Set}$ and the subdistribution functor $\mathcal{D}_{\leq 1} : \mathbf{Set} \rightarrow \mathbf{Set}$. Recall that $\mathcal{D}_{\leq 1}(X) = \{p : X \rightarrow [0, 1] \mid \sum_{x \in X} p(x) \leq 1\}$. As a parameter of codensity lifting, we take $(\Omega, \tau) = (([0, 1], d_{[0, 1]}), e : \mathcal{D}_{\leq 1}[0, 1] \rightarrow [0, 1])$, where

TABLE VII
UNTRIMMED CODENSITY BISIMILARITY GAME

position	pl.	possible moves
$P \in \mathbb{E}_X$	S	$k \in \mathbb{C}(X, \Omega)$ s.t. $\tau \circ Fk \circ c : (X, P) \dashrightarrow (\Omega, \Omega)$
$k \in \mathbb{C}(X, \Omega)$	D	$P' \in \mathbb{E}_X$ s.t. $k : (X, P') \dashrightarrow (\Omega, \Omega)$

TABLE VIII
UNTRIMMED CODENSITY GAME FOR BISIMULATION METRIC

position	pl.	possible moves
$d \in (\mathbf{PMet}_1)_X$	S	$k \in \mathbf{Set}(X, [0, 1])$ s.t. $e \circ Fk \circ c \notin \mathbf{PMet}_1(d, d_{[0,1]})$
$k \in \mathbf{Set}(X, [0, 1])$	D	$d' \in (\mathbf{PMet}_1)_X$ s.t. $k \notin \mathbf{PMet}_1(d', d_{[0,1]})$

e is the *expectation function* $e(p) = \sum_{r \in [0,1]} r \cdot p(r)$ and $d_{[0,1]}$ is the Euclidean metric. Let $c: X \rightarrow \mathcal{D}_{\leq 1}X$ be a coalgebra, identified with a Markov chain.

The codensity bisimilarity in this setting coincides with bisimulation metric from [5] (see also §I-A3). This fact is not hard to check directly; one can also derive the coincidence via Example III.4 and the observations in [8].

C. Joint Codensity Bisimulation

We introduce a notion of *joint codensity bisimulation*. This minor variation of codensity bisimulation becomes useful in the proof of soundness and completeness of our game notion (§IV).

Definition III.10 (joint codensity bisimulation). Assume the setting of Def. III.5. Let $\mathcal{V} \subseteq |\mathbb{E}_X|$; joins in \mathbb{E}_X are denoted by \sqcup . We say that \mathcal{V} is a *joint codensity bisimulation* over c if $\sqcup_{P \in \mathcal{V}} P$ is a codensity bisimulation over c .

For instance, the set of all codensity bisimulations is a joint codensity bisimulation, because the join $\nu\Phi_c^{\Omega, \tau}$ is the largest bisimulation (a consequence of the Knaster–Tarski theorem).

Lemma III.11. *In the setting of Def. III.5, the downset $\downarrow(\nu\Phi_c^{\Omega, \tau})$ is the largest joint codensity bisimulation (with respect to the inclusion order).* \square

IV. UNTRIMMED GAMES FOR CODENSITY BISIMILARITY

As the first main technical contribution, we introduce what we call the *untrimmed* version of codensity bisimilarity game. It is mathematically simple but its game arenas can become much bigger than necessary. The *trimmed* version of games—with smaller arenas—will be introduced later in §V, after developing necessary categorical infrastructure.

Assume the setting of Def. III.5 for the rest of the section.

Definition IV.1 (untrimmed codensity bisimilarity game). The *untrimmed codensity bisimilarity game* is the safety game played by two players D and S, shown in Table VII.

Lemma IV.2. *Let $\mathcal{V} \subseteq |\mathbb{E}_X|$. The following are equivalent.*

- 1) \mathcal{V} is an invariant for Duplicator (Def. II.3) in the safety game in Table VII.
- 2) \mathcal{V} is a joint codensity bisimulation over c . \square

Theorem IV.3. *The following coincide.*

- 1) The set of all winning positions for D.
- 2) The downset $\downarrow(\nu\Phi_c^{\Omega, \tau})$ of the codensity bisimilarity.

Proof. On the one hand, the set of all winning positions for D is the largest invariant for D, by Prop. II.4. On the other hand, the downset $\downarrow(\nu\Phi_c^{\Omega, \tau})$ is the largest joint codensity bisimulation over c . Thus, the statement follows from Lem. IV.2. \square

We conclude that our game characterizes the codensity bisimilarity $\nu\Phi_c^{\Omega, \tau}$ (Def. III.8).

Corollary IV.4 (soundness and completeness of untrimmed codensity games). *$P \in \mathbb{E}_X$ is a winning position for D if and only if $P \sqsubseteq \nu\Phi_c^{\Omega, \tau}$.* \square

Example IV.5. Recall Example III.9. Using the untrimmed codensity bisimilarity game, we can characterize the bisimulation metric from [5]. Our general definition (Def. IV.1) instantiates to the one in Table VIII, which is however more complicated than the game we exhibited in the introduction (Table II). For example, in Table VIII, Duplicator’s move is a pseudometric $d: X^2 \rightarrow [0, 1]$ rather than a triple (x, y, ε) .

V. TRIMMED CODENSITY GAMES FOR BISIMILARITY

Our previous untrimmed game (Table VII) is pleasantly simple from a theoretical point of view. However, as we saw in Example IV.5, its instances tend to have a much bigger arena than some known game notions.

Here we push our theory a step further, and present a fibrational construction that allows us to *trim* our games. We note that our construction still remains on the fibrational level of abstraction.

A. Generating Sets and Fibered Separators in a Fibration

We start by building some fibrational infrastructure. The following notion is a natural extension of the corresponding lattice-theoretic one. Unlike in the case of algebraic lattices, we do not assume the compactness of elements in \mathcal{G} .

Definition V.1 (generating set). Let $p: \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\square} -fibration and $X \in \mathbb{C}$ be an object. We say that a set $\mathcal{G} \subseteq |\mathbb{E}_X|$ is a *generating set* of the fiber \mathbb{E}_X if, for any $P \in \mathbb{E}_X$, there exists $\mathcal{A} \subseteq \mathcal{G}$ such that $\sqcup \mathcal{A} = P$.

Example V.2. Consider the \mathbf{CLat}_{\square} -fibration $\mathbf{EqRel} \rightarrow \mathbf{Set}$ and $X \in \mathbf{Set}$. For any $x, y \in X$, we define the equivalence relation $E_{x,y}$ to be the least one equating x, y , that is, $(z, w) \in E_{x,y}$ if and only if $(z = w \vee \{z, w\} = \{x, y\})$. Then the set $\mathcal{G} = \{E_{x,y} \mid x, y \in X\}$ of elements of the fiber \mathbf{EqRel}_X is a generating set.

Example V.3. Recall Example III.9. For $x, y \in X$ ($x \neq y$) and $r \in [0, 1]$, the pseudometric $d_{x,y,r}$ over X is defined by

$$d_{x,y,r}(z, w) = \begin{cases} 0 & z = w \\ r & \{z, w\} = \{x, y\} \\ 1 & \text{otherwise.} \end{cases}$$

Then the set of pseudometrics $\{d_{x,y,r} \mid x, y \in X, x \neq y, r \in [0, 1]\}$ is a generating set of the fiber $(\mathbf{PMet}_1)_X$.

One natural question is how to find such a generating set in the fiber over the state space. Below we show that a generating set of the fiber of a special object (called *fibred separator*) induces a generating set of each fiber by push-forward.

Definition V.4 (fibred separator). Let $p : \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration. We say that $S \in \mathbb{C}$ is a *fibred separator* if, for any $X \in \mathbb{C}$ and $P, Q \in \mathbb{E}_X$, we have

$$(\forall f \in \mathbb{C}(S, X). f^*P = f^*Q) \implies P = Q.$$

Theorem V.5. Let $S \in \mathbb{C}$ be a fibred separator of a \mathbf{CLat}_{\sqcap} -fibration $p : \mathbb{E} \rightarrow \mathbb{C}$, and $\mathcal{G} \subseteq |\mathbb{E}_S|$ be a generating set of \mathbb{E}_S . For any $X \in \mathbb{C}$, the following set is a generating set of \mathbb{E}_X :

$$\{f_*P \mid P \in \mathcal{G}, f \in \mathbb{C}(S, X)\}.$$

Here f_* denotes the pushforward along f (§II-B). \square

In fact, it was Thm. V.5 behind Examples V.2–V.3: in both cases, $2 \in \mathbf{Set}$ turns out to be a fibred separator for the fibrations in question ($\mathbf{EqRel} \rightarrow \mathbf{Set}$ and $\mathbf{PMet}_1 \rightarrow \mathbf{Set}$), and the presented generating sets are obtained via push-forward.

The following result is useful in finding fibred separators—see §VIII-F.

Proposition V.6 (change-of-base and fibred separators). Let $p : \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration, $R : \mathbb{D} \rightarrow \mathbb{C}$ be a functor with a left adjoint $L : \mathbb{C} \rightarrow \mathbb{D}$, and $S \in \mathbb{C}$ be a fibred separator for p . Then $LS \in \mathbb{D}$ is a fibred separator of the change-of-base fibration R^*p . \square

B. \mathcal{G} -Joint Codensity Bisimulation

We use generating sets to restrict moves in codensity games.

Definition V.7. In the setting of Def. III.5, let \mathcal{G} be a generating set of \mathbb{E}_X . A \mathcal{G} -joint codensity bisimulation over $c : X \rightarrow FX$ is a joint codensity bisimulation \mathcal{V} over c such that $\mathcal{V} \subseteq \mathcal{G}$.

Lemma V.8 (key lemma). Assume the setting of Def. III.5, and let \mathcal{G} be a generating set of \mathbb{E}_X . The intersection $(\downarrow(\nu\Phi_c^{\Omega, \tau})) \cap \mathcal{G}$ of the downset $\downarrow(\nu\Phi_c^{\Omega, \tau})$ and the generating set \mathcal{G} is the largest \mathcal{G} -joint codensity bisimulation.

Proof. Since \mathcal{G} is a generating set, the union of all elements of $\downarrow(\nu\Phi_c^{\Omega, \tau}) \cap \mathcal{G}$ is equal to $\nu\Phi_c^{\Omega, \tau}$. Thus, $\downarrow(\nu\Phi_c^{\Omega, \tau}) \cap \mathcal{G}$ is a \mathcal{G} -joint codensity bisimulation.

For any \mathcal{G} -joint codensity bisimulation \mathcal{V} , we have already shown $\mathcal{V} \subseteq \downarrow(\nu\Phi_c^{\Omega, \tau})$. We also have $\mathcal{V} \subseteq \mathcal{G}$ by definition. These imply $\mathcal{V} \subseteq \downarrow(\nu\Phi_c^{\Omega, \tau}) \cap \mathcal{G}$. \square

C. Trimmed Codensity Bisimilarity Games

The above structural results lead to our second game notion.

Definition V.9 (trimmed codensity bisimilarity game). Assume the setting of Def. III.5, and that $\mathcal{G} \subseteq \mathbb{E}_X$ is a generating set. The *codensity bisimilarity game* is the safety game played by two players D and S, shown in Table IX.

Assume the setting of Def. V.9 for the rest of the section.

TABLE IX
TRIMMED CODENSITY BISIMILARITY GAME

position	pl.	possible moves
$P \in \mathcal{G}$	S	$k \in \mathbb{C}(X, \Omega)$ s.t. $\tau \circ Fk \circ c : (X, P) \dashrightarrow (\Omega, \Omega)$
$k \in \mathbb{C}(X, \Omega)$	D	$P' \in \mathcal{G}$ s.t. $k : (X, P') \dashrightarrow (\Omega, \Omega)$

Lemma V.10. Let $\mathcal{V} \subseteq |\mathbb{E}_X|$. The following are equivalent:

- 1) \mathcal{V} is an invariant for D (Def. II.3) in the game in Table IX.
- 2) \mathcal{V} is a \mathcal{G} -joint codensity bisimulation over c . \square

Theorem V.11. The following sets coincide.

- 1) The set of D-winning positions in the game in Table IX.
- 2) The intersection $(\downarrow(\nu\Phi_c^{\Omega, \tau})) \cap \mathcal{G}$ of the downset of the codensity bisimilarity over c and the generating set \mathcal{G} . \square

We conclude that our second game characterizes the codensity bisimilarity $\nu\Phi_c^{\Omega, \tau}$ (Def. III.8) too.

Corollary V.12 (soundness and completeness of trimmed codensity games). In Def. V.9, $P \in \mathcal{G}$ is a winning position for Duplicator if and only if $P \sqsubseteq \nu\Phi_c^{\Omega, \tau}$. \square

VI. MULTIPLE OBSERVATION DOMAINS

We extend the theory so far and accommodate multiple observation domains and modalities. This extension is needed for some examples, such as those marked with \dagger in Table VI.

We consider the class $\mathbf{Lift}(F, p)$ of liftings of an endofunctor $F : \mathbb{C} \rightarrow \mathbb{C}$ along a \mathbf{CLat}_{\sqcap} -fibration $p : \mathbb{E} \rightarrow \mathbb{C}$. It comes with a natural pointwise partial order:

$$G \sqsubseteq H \iff \forall X \in \mathbb{E}. GX \sqsubseteq HX \quad (G, H \in \mathbf{Lift}(F, p)), \quad (3)$$

and the partially ordered class $\mathbf{Lift}(F, p)$ admits meets of arbitrary size. As done in the original codensity lifting of endofunctors in [15] (and monads in [14]), we extend the codensity lifting so that it takes a family of parameters $\{(\Omega_A, \tau_A)\}_{A \in \mathbb{A}}$, and returns the *intersection* of the codensity liftings of F with these parameters.

Definition VI.1 (codensity lifting of a functor with multiple parameters [15]). Let $F : \mathbb{C} \rightarrow \mathbb{C}$ be a functor, $p : \mathbb{E} \rightarrow \mathbb{C}$ be a \mathbf{CLat}_{\sqcap} -fibration, \mathbb{A} be a class, and $\{(\Omega_A, \tau_A)\}_{A \in \mathbb{A}}$ be an \mathbb{A} -indexed family of parameters (of the codensity lifting of F along p), which is denoted simply by (Ω, τ) . The (multiple-parameter) *codensity lifting* of F with (Ω, τ) is the endofunctor $F^{\Omega, \tau} : \mathbb{E} \rightarrow \mathbb{E}$ defined by the intersection of the codensity liftings:

$$F^{\Omega, \tau} P = \prod_{A \in \mathbb{A}} F^{\Omega_A, \tau_A} P,$$

that is, $\prod_{A \in \mathbb{A}, k \in \mathbb{E}(P, \Omega_A)} (\tau_A \circ F(p(k)))^*(\Omega_A).$

The rest of the theoretical development is completely parallel to the one in the previous sections. Therefore we only

TABLE X
TRIMMED CODENSITY BISIMILARITY GAME WITH MULTIPLE
OBSERVATIONS

position	pl.	possible moves
$P \in \mathcal{G}$	S	$A \in \mathbb{A}$ and $k \in \mathbb{C}(X, \Omega_A)$ s.t. $\tau_A \circ Fk \circ c : (X, P) \dashrightarrow (\Omega_A, \Omega_A)$
$A \in \mathbb{A}$ and $k \in \mathbb{C}(X, \Omega_A)$	D	$P' \in \mathcal{G}$ s.t. $k : (X, P') \dashrightarrow (\Omega_A, \Omega_A)$

present key definitions and the main result (Cor. VI.4). The omitted definitions and results can be recovered from the ones in §III–V, by replacing a single-parameter codensity lifting (Def. III.1) by a multi-parameter one (Def. VI.1).

Definition VI.2 (codensity bisimulation). Assume the setting of Def. VI.1. Let $c : X \rightarrow FX$ be an F -coalgebra. An object $P \in \mathbb{E}_X$ is a *codensity bisimulation* over c if $c : (X, P) \dashrightarrow (FX, F^{\Omega, \tau} P)$; that is, $c : X \rightarrow FX$ is decent with respect to the designated indistinguishability structures.

Definition VI.3 (codensity bisimilarity game). In the setting of Def. VI.2, let \mathcal{G} be a generating set of \mathbb{E}_X . The *codensity bisimilarity game* is the safety game, played by two players D and S, shown in Table X.

Corollary VI.4 (soundness and completeness of codensity games). Assume the setting of Def. VI.3. $P \in \mathbb{E}_X$ is a winning position for Duplicator if and only if $P \sqsubseteq \nu \Phi_c^{\Omega, \tau}$. \square

Example VI.5 (bisimulation topology for deterministic automata). Here we describe the topological example in Table V. Consider the \mathbf{CLat}_{\square} -fibration $\mathbf{Top} \rightarrow \mathbf{Set}$ and the functor $A_{\Sigma} = 2 \times (_)^{\Sigma} : \mathbf{Set} \rightarrow \mathbf{Set}$, where Σ is a fixed alphabet. Coalgebras for this functor are deterministic automata over Σ ; see e.g. [20], [21].

We take the following data as a parameter of codensity lifting (cf. Def. VI.1): $\mathbb{A} = \{\varepsilon\} \cup \Sigma$, Ω_{α} is the Sierpinski space for each $\alpha \in \mathbb{A}$, and the modalities $\tau_{\varepsilon}, \tau_a : A_{\Sigma} 2 \rightarrow 2$ (where $a \in \Sigma$) are defined by

$$\tau_{\varepsilon}(t, \rho) = t \quad \text{and} \quad \tau_a(t, \rho) = \rho(a).$$

Recall that the Sierpinski space is the set $2 = \{\perp, \top\}$ with the topology $\{\emptyset, \{\top\}, 2\}$; this observation domain models the situation where acceptance of a word is only *semi*-decidable.

Let $c : X \rightarrow A_{\Sigma} X$ be a deterministic automata. The above choice of parameters leads to the following codensity bisimilarity: the state space X is equipped with the topology generated by the following family of open sets.

$$\{x \in X \mid w \text{ is accepted from } x\} \subseteq X, \quad \text{for each } w \in \Sigma^*$$

One can extract various information from this *bisimulation topology* via standard topological constructs. For example, the specialization order of this topology coincides with the language inclusion order.

For illustration by comparison, consider changing the observation domain from the Sierpinski space to the discrete 2-point set. The bisimulation topology over X is now generated by

$$\{x \in X \mid w \text{ is accepted from } x\} \text{ and } \{x \in X \mid w \text{ is not accepted from } x\}, \quad \text{for each } w \in \Sigma^*.$$

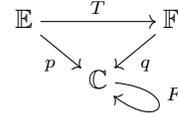
We can now observe rejection of a word, too, because $\{\perp\} \subseteq 2$ is open. The specialization order of this topology is the language equivalence, and it satisfies the R0 separation axiom (while the last Sierpinski example does not).

We take these examples of bisimulation topology as a process-semantic incarnation of the ‘‘observability via topology, computability via continuity’’ paradigm from domain theory. The definition of codensity bisimulation (cf. Def. III.1) fits well with this intuition, too: a continuous map $k : (X, P) \dashrightarrow \Omega$ in Def. III.1 is a ‘‘computable observation’’; accordingly, an open set of the bisimulation topology is a property that is decided by finitely many of those computable observations.

VII. TRANSFER OF CODENSITY BISIMILARITIES

In our formulation, for the same endofunctor $F : \mathbb{C} \rightarrow \mathbb{C}$, we can use various \mathbf{CLat}_{\square} -fibrations and parameters (Ω, τ) to equip F -coalgebras with different bisimilarity-like notions. Some relations among those codensity bisimilarities can be categorically captured by the following theorem.

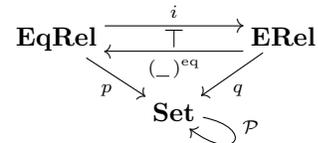
Theorem VII.1 (transfer of codensity bisimilarity). Let $p : \mathbb{E} \rightarrow \mathbb{C}$ and $q : \mathbb{F} \rightarrow \mathbb{C}$ be \mathbf{CLat}_{\square} -fibrations, $F : \mathbb{C} \rightarrow \mathbb{C}$ be an endofunctor, $c : X \rightarrow FX$ be an F -coalgebra, $T : \mathbb{E} \rightarrow \mathbb{F}$ be a full and faithful fibered functor from p to q preserving fibered meets, and $\{(\Omega_A, \tau_A)\}_{A \in \mathbb{A}}$ be an \mathbb{A} -indexed family of parameters for codensity lifting of F along p .



In this setting, $\{(T\Omega_A, \tau_A)\}_{A \in \mathbb{A}}$ is an \mathbb{A} -indexed family of parameters for codensity lifting of F along q , and we have $\nu \Phi_c^{T\Omega, \tau} = T(\nu \Phi_c^{\Omega, \tau})$. \square

Example VII.2. We show that the codensity bisimilarities in Examples III.2 & III.3 are indeed the usual bisimilarity notions for Kripke frames. Recall that they are build on the two \mathbf{CLat}_{\square} -fibrations $\mathbf{EqRel} \rightarrow \mathbf{Set}$ and $\mathbf{ERel} \rightarrow \mathbf{Set}$.

We first note that the inclusion functor $i : \mathbf{EqRel} \rightarrow \mathbf{ERel}$ is a reflection, having the equivalence closure $(_)^{\text{eq}} : \mathbf{ERel} \rightarrow \mathbf{EqRel}$ as the left adjoint. It follows that i is meet-preserving. Moreover, i is fibered.



We introduce shorthands $\dot{\mathcal{P}}_2, \dot{\mathcal{P}}_3$ for the liftings in Examples III.2 & III.3:

$$\begin{aligned}\dot{\mathcal{P}}_2 &= \mathcal{P}^{\text{Eq}_2, \diamond} : \mathbf{EqRel} \rightarrow \mathbf{EqRel} \quad (\text{Example III.2}), \\ \dot{\mathcal{P}}_3 &= \mathcal{P}^{\text{Eq}_2, \diamond} : \mathbf{ERel} \rightarrow \mathbf{ERel} \quad (\text{Example III.3})\end{aligned}$$

Now, for the sake of our proof, let us introduce a relational lifting $\dot{\mathcal{P}}_1 : \mathbf{ERel} \rightarrow \mathbf{ERel}$ of \mathcal{P} along $\mathbf{ERel} \rightarrow \mathbf{Set}$, for which it is obvious that the corresponding bisimilarity notion is the usual bisimilarity for Kripke frames. We do so in concrete terms, instead of as a codensity lifting:

$$(S, T) \in \dot{\mathcal{P}}_1(R) \iff (\forall x \in S. \exists y \in T. (x, y) \in R) \wedge (\forall y \in T. \exists x \in S. (x, y) \in R).$$

We note that $\dot{\mathcal{P}}_2$ is the restriction of $\dot{\mathcal{P}}_1$ from \mathbf{ERel} to \mathbf{EqRel} along i . Note also that $\dot{\mathcal{P}}_3 = \dot{\mathcal{P}}_1 \circ i \circ (_)^{\text{eq}}$.

Let $c : X \rightarrow \mathcal{P}X$ be a Kripke frame and $\Phi_i = c^* \circ \dot{\mathcal{P}}_i$ ($i = 1, 2, 3$) be the predicate transformer corresponding to each lifting. Theorem VII.1 states that $\nu\Phi_3 = i(\nu\Phi_2)$.

Furthermore, by $\dot{\mathcal{P}}_1 \sqsubseteq \dot{\mathcal{P}}_3$ (where \sqsubseteq is the order in (3)), we have $\nu\Phi_1 \sqsubseteq \nu\Phi_3$. From $i \circ \dot{\mathcal{P}}_2 = \dot{\mathcal{P}}_1 \circ i$ and fiberedness of c , we can see that $i(\nu\Phi_2)$ is a fixed point of Φ_1 , which yields $i(\nu\Phi_2) \sqsubseteq \nu\Phi_1$ by the Knaster–Tarski theorem. The three (in)equalities so far allow us to conclude $\nu\Phi_3 = i(\nu\Phi_2) = \nu\Phi_1$, stating that the conventional bisimilarity $\nu\Phi_1$ is equal to the codensity bisimilarities in Examples III.2 & III.3. As a consequence, the conventional bisimilarity $\nu\Phi_1$ is necessarily an equivalence relation.

VIII. EXAMPLES

A. Kripke Frames and (Conventional) Bisimilarity

We consider $\mathbf{EqRel} \rightarrow \mathbf{Set}$ as an underlying \mathbf{CLat}_{Γ} -fibration, and a Kripke frame $c : X \rightarrow \mathcal{P}X$ (as a \mathcal{P} -coalgebra). We further use the codensity lifting $\mathcal{P}^{\text{Eq}_2, \diamond}$ (Example III.2) and the generating set described in Example V.2 to trim games; the resulting game is shown in Table XI. As shown in Example VII.2, the codensity bisimilarity $\nu\Phi_c^{\text{Eq}_2, \diamond}$ (which is $\nu\Phi_2$ in Example VII.2) coincides with conventional bisimilarity on c using the standard relational lifting of \mathcal{P} to \mathbf{ERel} .

Theorem VIII.1. *(x, y) is a D-winning position if and only if $(x, y) \in \nu\Phi_c^{\text{Eq}_2, \diamond}$, if and only if x and y are bisimilar. \square*

Extension of this result to labeled transition systems and Kripke models (with valuations) is straightforward, using suitable choice of endofunctors—see also §VIII-B.

B. Deterministic Automata and Their Language Equivalence

We use the functor $A_\Sigma : \mathbf{Set} \rightarrow \mathbf{Set}$ from Example VI.5, for which a coalgebra is a deterministic automaton. Let us lift A_Σ along the fibration $\mathbf{EqRel} \rightarrow \mathbf{Set}$, with the same modalities $\tau_\varepsilon, \tau_a : A_\Sigma 2 \rightarrow 2$ as in Example VI.5 (where $a \in \Sigma$). Our observation domains are $\Omega_\alpha = (2, \text{Eq}_2)$ for all $\alpha \in \{\varepsilon\} \cup \Sigma$.

TABLE XI
CODENSITY BISIMILARITY GAME FOR CONVENTIONAL BISIMILARITY

position	pl.	possible moves
$(x, y) \in X \times X$	S	$k \in \mathbf{Set}(X, 2)$ s.t. $\exists x' \in c(x). k(x') = \top$ $\not\exists y' \in c(y). k(y') = \top$
$k \in \mathbf{Set}(X, 2)$	D	(x'', y'') s.t. $k(x'') \neq k(y'')$

TABLE XII
CODENSITY BISIMILARITY GAME FOR DETERMINISTIC AUTOMATA AND THEIR LANGUAGE EQUIVALENCE

position	pl.	possible moves
$(x, y) \in X \times X$	S	If $\pi_1(x) \neq \pi_1(y)$ then S wins If $\pi_1(x) = \pi_1(y)$ then $a \in \Sigma$ and $k \in \mathbf{Set}(X, 2)$ s.t. $k(\pi_2(x)(a)) \neq k(\pi_2(y)(a))$
$a \in \Sigma$ and $k \in \mathbf{Set}(X, 2)$	D	$(x'', y'') \in X \times X$ s.t. $k(x'') \neq k(y'')$

TABLE XIII
CODENSITY BISIMILARITY GAME FOR DETERMINISTIC AUTOMATA AND BISIMULATION TOPOLOGY

position	pl.	possible moves
$\mathcal{O} \in \mathbf{Top}_X$	S	$a \in \{\varepsilon\} \cup \Sigma$ and $k \in \mathbf{Set}(X, 2)$ s.t. $\tau_a \circ (A_\Sigma k) \circ c : (X, \mathcal{O}) \rightarrow (2, \Omega_a)$
$a \in \{\varepsilon\} \cup \Sigma$ and $k \in \mathbf{Set}(X, 2)$	D	$\mathcal{O}' \in \mathbf{Top}_X$ s.t. $k : (X, \mathcal{O}') \rightarrow (2, \Omega_a)$

TABLE XIV
CODENSITY BISIMILARITY GAME FOR NONDETERMINISTIC AUTOMATA AND THEIR BISIMILARITY

position	pl.	possible moves
$(x, y) \in X \times X$	S	If $\pi_1(x) \neq \pi_1(y)$ then S wins If $\pi_1(x) = \pi_1(y)$ then $a \in \Sigma$ and $k \in \mathbf{Set}(X, 2)$ s.t. $\exists x' \in \pi_2(x)(a). k(x') = \top$ $\not\exists y' \in \pi_2(y)(a). k(y') = \top$
$a \in \Sigma$ and $k \in \mathbf{Set}(X, 2)$	D	$(x'', y'') \in X \times X$ s.t. $k(x'') \neq k(y'')$

TABLE XV
CODENSITY BISIMILARITY GAME FOR PROBABILISTIC BISIMILARITY

position	pl.	possible moves
$(x, y) \in X \times X$	S	$r \in [0, 1]$ and $k \in \mathbf{Set}(X, 2)$ s.t. $c(x)(k^{-1}(\top)) \geq r > c(y)(k^{-1}(\top))$, or $c(y)(k^{-1}(\top)) \geq r > c(x)(k^{-1}(\top))$
$r \in [0, 1]$ and $k \in \mathbf{Set}(X, 2)$	D	(x'', y'') s.t. $k(x'') \neq k(y'')$

The resulting codensity lifting $(A_\Sigma)^{\Omega, \tau} : \mathbf{EqRel} \rightarrow \mathbf{EqRel}$ is concretely described as follows.

$$(A_\Sigma)^{\Omega, \tau}(R) = \left\{ \begin{array}{l} ((t_1, \rho_1), \\ (t_2, \rho_2)) \\ \in (A_\Sigma X)^2 \end{array} \left| \begin{array}{l} \forall k : X \rightarrow 2. \\ (\forall x, y \in X. (x, y) \in R \Rightarrow k(x) = k(y)) \\ \Rightarrow (t_1 = t_2) \wedge \\ (\forall a \in \Sigma. (k \circ \rho_1)(a) = (k \circ \rho_2)(a)) \end{array} \right. \right\}.$$

Let $c : X \rightarrow A_\Sigma X$ be a deterministic automaton. It is not hard to see that the codensity bisimilarity $\nu\Phi_c^{\Omega, \tau}$ coincides with language equivalence of deterministic automata. Our trimmed codensity game is shown in Table XII (in a slightly optimized form). The game therefore characterizes the language equivalence (Cor. VI.4).

C. Deterministic Automata and the Language Topology

We introduced two versions of *bisimulation topology* for deterministic automata in Example VI.5. They are in close correspondences with accepted languages; therefore we call them *language topologies*.

For the first topology in Example VI.5 (where Ω is the Sierpinski space, capturing that acceptance is only semi-decidable), the corresponding (untrimmed) codensity game is shown in Table XIII. It follows from our general results that the game notion is sound and complete.

We have not yet found a good way (e.g. generating sets) of trimming the game arena; this is left as future work.

D. Nondeterministic Automata and Bisimilarity

Let us now turn to nondeterministic automata, that is, N_Σ -coalgebras for the functor $N_\Sigma = 2 \times (\mathcal{P}_-)^\Sigma$. Much like the situation for DFAs, we lift this functor along the \mathbf{CLat}_\square -fibration $\mathbf{EqRel} \rightarrow \mathbf{Set}$ by codensity lifting with multiple observation domains, as follows. Let \mathbb{A} be the set $\{\varepsilon\} \cup \Sigma$. We set the parameter of codensity lifting as follows, where $a \in \Sigma$.

$$\Omega_\varepsilon = \Omega_a = \text{Eq}_{Q_2}, \tau_\varepsilon(t, \rho) = t, \tau_a(t, \rho) = \diamond(\rho(a)).$$

The resulting codensity lifting $(N_\Sigma)^{\Omega, \tau}: \mathbf{EqRel} \rightarrow \mathbf{EqRel}$ is concretely described as

$$(N_\Sigma)^{\Omega, \tau}(R) = \left\{ \begin{array}{l} ((t_1, \rho_1), \\ (t_2, \rho_2)) \\ \in (N_\Sigma X)^2 \end{array} \left| \begin{array}{l} \forall k: X \rightarrow 2. \\ (\forall x, y \in X. (x, y) \in R \Rightarrow k(x) = k(y)) \\ \Rightarrow (t_1 = t_2) \wedge \\ (\forall a \in \Sigma. \top \in (k \circ \rho_1)(a) \\ \Leftrightarrow \top \in (k \circ \rho_2)(a)) \end{array} \right. \right\}.$$

Let $c: X \rightarrow N_\Sigma X$ be a nondeterministic automaton. It is again not hard to see that the codensity bisimilarity $\nu\Phi_c^{\Omega, \tau}$ is the usual notion of bisimilarity of nondeterministic automata. Our trimmed codensity game is shown in Table XIV, in a slightly optimized form, and it captures bisimilarity.

A topological variant of the above story is possible, much like in §VIII-C.

E. Markov Chains and Bisimulation Metric

Recall Examples III.9, IV.5, and V.3. Markov chains are $\mathcal{D}_{\leq 1}$ -coalgebras. We use the \mathbf{CLat}_\square -fibration $\mathbf{PMet}_1 \rightarrow \mathbf{Set}$, taking pseudometrics as a notion of indistinguishability. With the lifting parameter we described in Example III.9, we get the bisimulation metric as the codensity bisimilarity. We can use the generating set described in Example V.3 to obtain a trimmed codensity game; the resulting game essentially coincides with the one in Table II in the introduction. Therefore, Cor. V.12 gives an abstract proof for the correctness of the game.

F. Continuous State Markov Chains and Bisimulation Metric

In order to accommodate continuous state Markov chains (for which measurable structures are essential), we consider an example that involves \mathbf{Meas} . Continuing §VIII-E, by the change-of-base along the forgetful functor $U: \mathbf{Meas} \rightarrow \mathbf{Set}$,

we get another \mathbf{CLat}_\square -fibration $U^*(\mathbf{PMet}_1) \rightarrow \mathbf{Meas}$. A continuous state Markov chain is a coalgebra $X \rightarrow \mathcal{G}_{\leq 1} X$ of the so-called *sub-Giry* functor over \mathbf{Meas} —see, e.g., [30].

Since the forgetful functor $\mathbf{Meas} \rightarrow \mathbf{Set}$ has a left adjoint, Prop. V.6 gives us a fibered separator for $U^*(\mathbf{PMet}_1) \rightarrow \mathbf{Meas}$. This gives us a game notion similar to that in §VIII-E.

G. Markov Chains and Probabilistic Bisimilarity

In order to define bisimilarity-like equivalence relation on Markov chains, we first lift $\mathcal{D}_{\leq 1}$ along the \mathbf{CLat}_\square -fibration $\mathbf{EqRel} \rightarrow \mathbf{Set}$. For that purpose, here we use the following multiple lifting parameters. The index set is $\mathbb{A} = [0, 1]$. For each $r \in \mathbb{A}$, we set $\Omega_r = (2, \text{Eq}_{Q_2})$, and define a *threshold modality* $\tau_r: \mathcal{D}_{\leq 1} 2 \rightarrow 2$ by $\tau_r(p) = \top$ if and only if $p(\top) \geq r$. Then for any $R \in \mathbf{EqRel}_X$, the relation part of the codensity lifting $\mathcal{D}_{\leq 1}^{\Omega, \tau}(X, R)$ relates $p, q \in \mathcal{D}_{\leq 1}(X)$ if and only if

$$\begin{aligned} \forall r \in [0, 1]. \forall k: X \rightarrow 2. & ((\forall (x, y) \in R. k(x) = k(y)) \\ \Rightarrow & (\sum_{x \in k^{-1}(\top)} p(x) \geq r \Leftrightarrow \sum_{x \in k^{-1}(\top)} q(x) \geq r)). \end{aligned}$$

Let us fix a Markov chain $c: X \rightarrow \mathcal{D}_{\leq 1} X$. All these data give rise to $\mathcal{D}_{\leq 1}^{\Omega, \tau}$ and $\nu\Phi_c^{\Omega, \tau}$ as in Definitions VI.2 and III.8. It is not hard to see that the resulting codensity bisimilarity coincides with probabilistic bisimilarity in [4]. Note, for example, that a relation-preserving map $k: (X, R) \rightarrow (2, \text{Eq}_{Q_2})$ coincides with an R -closed subset of X . The resulting trimmed codensity game is in Table XV. It is essentially the same as Table I (arising from [13]). The difference is that r is additionally present in Table XV; it is easy to realize that r plays no role in the game.

IX. CONCLUSIONS AND FUTURE WORK

Motivated by some recent works [8], [10], [11], [13], and especially by the similarity of the two games (Tables I and II), we introduced a fibrational framework that uniformly describes the correspondence between various bisimilarity notions and games. The fibrational abstraction allows us to accommodate some new examples, such as bisimulation topology. Moreover, the structural theory developed in §VI–VII provides new insights to the nature of bisimilarity, we believe, identifying the crucial role of observation maps ($k: X \rightarrow \Omega$ in Def. III.1) in bisimulation notions.

As future work, we intend to accommodate *modal logics* as is done in [10]. We are also interested in using games with more complex winning conditions (e.g. parity); they have been used for (bi)simulation notions for Büchi and parity automata [32]. Finally, we will pursue the algorithmic use of the current results.

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