

CONSEQUENCE-DRIVEN REASONING FOR HORN-SHIQ ONTOLOGIES

Yevgeny Kazakov

Oxford University Computing Laboratory

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OUTLINE

1 INTRODUCTION

2 CONSEQUENCE-BASED PROCEDURES



RESULTS OVERVIEW

- Classification times for some well-known large ontologies:

	GO	NCI	Galen v.0	Galen v.7	SNOMED
Concepts:	20465	27652	2748	23136	389472
FACT++	15.24	6.05	465.35	—	650.37
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PELLET	72.02	26.47	—	—	—
CEL	1.84	5.76	—	—	1185.70



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CB	1.17	3.57	0.32	9.58	49.44
Speed-Up:	1.57X	1.61X	143X	∞	13.15X

- The improvement is obtained using a new **consequence-based** reasoning procedure

available at:

cb-reasoner.googlecode.com



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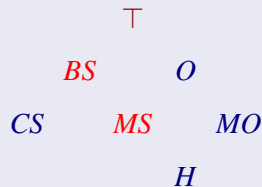
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TAXONOMY





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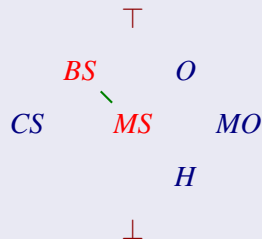
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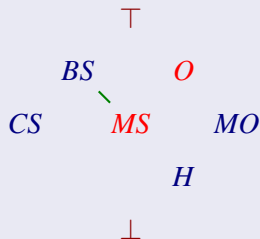
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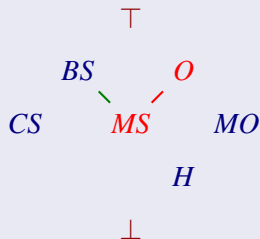
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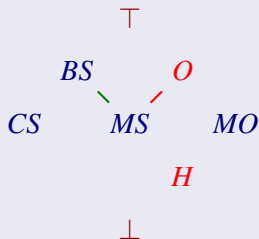
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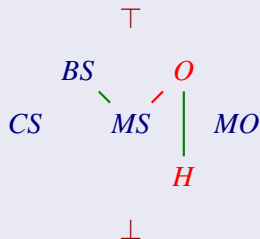
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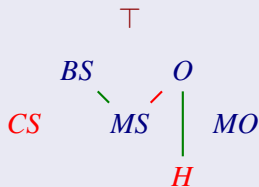
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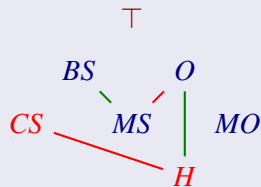
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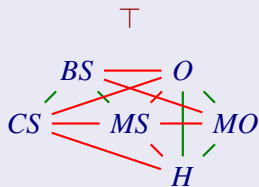
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- 1 Classification requires **enumeration**:
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 - Axioms $C \sqsubseteq D$ in general result in a disjunction $\top \sqsubseteq \neg C \sqcup D$
 - Non-determinism can be reduced using **absorption** and **hyper-tableaux** rules.



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- 3 **Large and highly cyclic models** caused by existential dependencies $C \sqsubseteq \exists R.D$ (**especially apparent for Galen**)

ONTOLOGY

Heart $\sqsubseteq \exists$ isPartOf.CirculatorySystem
CirculatorySystem $\sqsubseteq \exists$ hasPart.LeftLung
LeftLung $\sqsubseteq \exists$ isPartOf.RespiratorySystem
RespiratorySystem $\sqsubseteq \exists$ hasPart.Trachea
and so on . . .



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 \mathcal{EL} FAMILY OF DLs

- \mathcal{EL} [Baader et al., IJCAI 2003, 2005] is a **lightweight DL**:
 - concepts are constructed using \top , $C \sqcap D$, and $\exists R.C$
 - axioms are $C \sqsubseteq D$ and $C \equiv D$
 - \mathcal{EL}^{++} adds \perp , $R_1 \cdots R_n \sqsubseteq R$, nominals, safe datatypes

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- Most of the axioms in Galen are expressed in \mathcal{ELH}

EXAMPLE

$\text{KidneyExamination} \equiv \text{ClinicalAct} \sqcap$
 $\exists \text{hasSubprocess} . (\text{ExaminingProcess} \sqcap \exists \text{involves} . \text{Kidney})$

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(language = \mathcal{ELH})



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- Performs the full classification “in one pass”
- Derives only subsumptions that are implied ($< 1\%$ of all)
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The “only” disadvantage:

- the language may be too restricted
(tractable reasoning is the primary focus)



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inverse roles and **role functionality**:

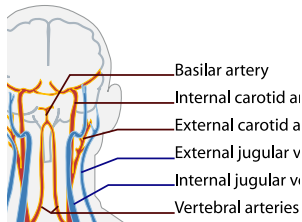
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BasilarArtery $\sqsubseteq \exists \text{isBranchOf.VertebralArtery}$

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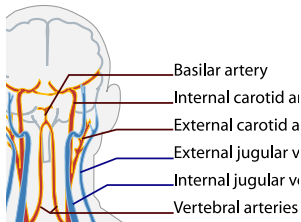
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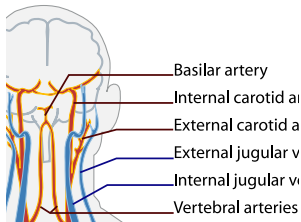


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- We are not scared of the high complexity!

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	positive	negative
universal restriction:	$A \sqsubseteq \forall R.B$	—

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EXAMPLE

- $A \sqsubseteq \forall R.(\neg B)$ $\exists R^-.A \sqsubseteq \leq 1 R.(B \sqcup C)$ are OK
- $A \sqsubseteq B \sqcup C$ $\forall R.A \sqsubseteq B$ $A \sqsubseteq \leq 2 R.B$ are **not** OK



NEW INFERENCE RULES

- Interactions between existential and universal restrictions:

$$\mathbf{1} \quad \frac{A \sqsubseteq \exists R.B \quad B \sqsubseteq \forall R^{-}.C}{A \sqsubseteq C} \quad \Leftrightarrow \quad \mathbf{CR5} \quad \frac{A \sqsubseteq \exists R.B \quad \exists R.B \sqsubseteq C}{A \sqsubseteq C}$$



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▶ all rules

- The general form of derived axioms:

$$\prod A_i \sqsubseteq B \quad \prod A_i \sqsubseteq \exists R. \prod B_j$$

– exponentially-many in worst case (which is optimal)



RESULTS

- A novel classification procedure for Horn *SHIQ* ontologies
- Key advantages over model-building procedures:
 - 1 deterministic
 - 2 full classification in one pass
 - 3 no problems with existential dependencies $C \sqsubseteq \exists R.D$
 - 4 optimal for Horn-*SHIQ* and *ELH* (pay-as-you-go).
- The implementation exhibits a significant speedup:

	GO	NCI	Galen v.0	Galen v.7	SNOMED
FACT++	15.24	6.05	465.35	—	650.37
HERMIT	199.52	169.47	45.72	—	—
PELLET	72.02	26.47	—	—	—
CEL	1.84	5.76	—	—	1185.70
CB	1.17	3.57	0.32	9.58	49.44
Speed-Up:	1.57X	1.61X	143X	∞	13.15X

available at:

cb-reasoner.googlecode.com



THE INFERENCE RULES FOR HORN *SHIQ*

$$\frac{}{M \sqcap A \sqsubseteq A}$$

$$\frac{}{M \sqsubseteq \top}$$

$$\frac{M \sqsubseteq A_1 \dots M \sqsubseteq A_n}{M \sqsubseteq C} : \prod_{i=1}^n A_i \sqsubseteq C \in \mathcal{O}$$

$$\frac{M \sqsubseteq \exists R.N \quad N \sqsubseteq \perp}{M \sqsubseteq \perp}$$

$$\frac{M \sqsubseteq \exists R_1.N \quad M \sqsubseteq \forall R_2.A}{M \sqsubseteq \exists R_1.(N \sqcap A)} : R_1 \sqsubseteq_{\mathcal{O}} R_2$$

$$\frac{M \sqsubseteq \exists R_1.N \quad N \sqsubseteq \forall R_2.A}{M \sqsubseteq A} : R_1 \sqsubseteq_{\mathcal{O}} R_2^-$$

$$M \sqsubseteq \exists R_1.N_1 \quad N_1 \sqsubseteq B$$

$$M \sqsubseteq \exists R_2.N_2 \quad N_2 \sqsubseteq B$$

$$M \sqsubseteq \leq 1 S.B$$

$$\frac{}{M \sqsubseteq \exists R_1.(N_1 \sqcap N_2)} : \begin{array}{l} R_1 \sqsubseteq_{\mathcal{O}} S \\ R_2 \sqsubseteq_{\mathcal{O}} S \end{array}$$

$$M \sqsubseteq \exists R_1.N_1 \quad M \sqsubseteq B$$

$$N_1 \sqsubseteq \exists R_2.(N_2 \sqcap A)$$

$$N_1 \sqsubseteq \leq 1 S.B \quad N_2 \sqcap A \sqsubseteq B$$

$$\frac{}{M \sqsubseteq A \quad M \sqsubseteq \exists R_2^- . N_1} : \begin{array}{l} R_1 \sqsubseteq_{\mathcal{O}} S^- \\ R_2 \sqsubseteq_{\mathcal{O}} S \end{array}$$

Where $M, N = \prod A_i$