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STRUCTURAL RESOURCES FOR Q(x)

• An attempt to formulate quantum mechanics in *symmetric monoidal categories*, i.e. using *CS methods*.

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- A novel take on *quantum logic* and *Q-automation*.
- *High-level methods* for quantum info. and comp.
- Intuitive purely graphical quantum reasoning.

— security related —

B. Coecke, B.-S. Wang, Q.-L. Wang, Y.-J. Wang and Q.-Y. Zhang (2010) *Graphical calculus for quantum key distribution*. ENTCS (QPL'09 volume).

B. Coecke and S. Perdrix (2010) *Environment and classical channels in categorical quantum mechanics*. Computer Science Logic, LNCS 6247. arXiv:1004.1598

A. Hillebrand (2011) *Quantum protocols involving multiparticle entanglement and their representations in the zx-calculus*. MSC thesis, University of Oxford, 2011. http://www.cs.ox.ac.uk/people/bob.coecke/Anne.pdf

CATEGORICAL Q.M. IN †-COMPACT CATEGORIES

S. Abramsky and B.C. – LiCS'04 – arXiv:0808.1023

— (physical) data in monoidal category — Systems:

A B C

Processes:

 $A \xrightarrow{f} A \qquad A \xrightarrow{g} B \qquad B \xrightarrow{h} C$

Compound systems:

 $A \otimes B$ I $A \otimes C \xrightarrow{f \otimes g} B \otimes D$

Temporal composition:

$$A \xrightarrow{h \circ g} C := A \xrightarrow{g} B \xrightarrow{h} C \qquad A \xrightarrow{1_A} A$$

— graphical notation —



— merely a new notation? —

$(g \circ f) \otimes (k \circ h) = (g \otimes k) \circ (f \otimes h)$



— graphical notation —

Thm. [Joyal & Street '91] An equational statement between expressions in symmetric monoidal categorical language holds if and only if it is derivable in the graphical notation via homotopy. — quantum metric —

$f: A \to B$



— quantum metric —

 $f^{\dagger} \colon B \to A$



— (pure) classical vs. quantum —













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Thm. [Selinger '08] An equational statement between expressions in dagger compact symmetric monoidal categorical language holds if and only if it is derivable in the category of finite dimensional Hilbert spaces, linear maps, tensor product, and adjoints.



\Rightarrow quantum teleportation

CLASSICAL STRUCTURES IN CATEGORICAL Q.M.

B.C. and D. Pavlovic – arXiv:quant-ph/0608035 B.C., D. Pavlovic and J. Vicary – arXiv:0810.0812 B.C., E.O. Paquette and D. Pavlovic – arXiv:0904.1997 B.C. and S. Perdrix – arXiv:1004.1598

quantum data cannot be copied nor deleted

quantum data cannot be copied nor deleted

classical data CAN be copied and deleted

NON-FEATURE: quantum data cannot be copied nor deleted

FEATURE:

classical data CAN be copied and deleted

NON-FEATURE: quantum data cannot be copied nor deleted

FEATURE:

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A commutative monoid is object A with morphisms



s.t.

A commutative monoid is object A with morphisms



s.t.

A cocommutative comonoid is object A with morphisms



s.t.



FdHilb:

$$\begin{array}{c} \swarrow \\ & & & \\ & & & \\ & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & &$$

Theorem. Special dagger commutative Frobenius algebras (†-SCFAs) in FHilb, that is,



exactly correspond with **orthonormal bases** on the underlying Hilbert space via the correspondence:

 $\{ |i\rangle \}_i \quad \longleftrightarrow \quad |i\rangle \mapsto |ii\rangle$

A **†SCFA** is a pair:



which is such that:

 $\bigcup_{i=1}^{n} \bigcup_{i=1}^{n} \bigcup_{i$ $\bigcirc = \checkmark \bigcirc = |$

A **†SCFA** is a family:



which is such that, for k > 0:



Prop. The 2/0-spider \sim compactness:







COMPLEMENTARITY IN CATEGORICAL Q.M.

B.C. and R. Duncan – ICALP'08 – arXiv:0906.4725

— complementarity —

Two bases

 $\{|0\rangle, \dots, |n\rangle\}$ and $\{|0\rangle, \dots, |n\rangle\}$ are **complementary** (or **unbiased**) if $|\langle i || j \rangle| = \frac{1}{\sqrt{\dim}}$ yielding equal transition probabilities.



— complementarity —

Thm. Complementarity means:


Z-spin:



X-spin:





i.e.

$$(\delta_Z^{\dagger} \otimes 1) \circ (1 \otimes \delta_X) = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = CNOT$$

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \circ \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = ?$$









— phases —

Thm. Unbiased states of an observable always constitute an Abelian group with conjugates as inverses.



such that:



— phases —

For qubits in FHilb with green $\equiv \{|0\rangle, |1\rangle\} \equiv Z$:

$$\mathbf{\dot{\alpha}} = \begin{pmatrix} 1\\ e^{i\alpha} \end{pmatrix}_{Z} \qquad \mathbf{\dot{\alpha}} = Z_{\alpha} = \begin{pmatrix} 1 & 0\\ 0 & e^{i\alpha} \end{pmatrix}_{Z}$$

$$-phases -$$
For qubits in FHilb with green $\equiv \{|0\rangle, |1\rangle\} \equiv Z$:

$$\oint_{C} = \begin{pmatrix} 1 \\ e^{i\alpha} \end{pmatrix}_{Z} \quad \oint_{Q} = Z_{\alpha} = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\alpha} \end{pmatrix}_{Z}$$

These are relative phases for Z, hence in X-Y:



$$-phases -$$
For qubits in FHilb with red $\equiv \{|+\rangle, |-\rangle\} \equiv X$:
$$\oint = \begin{pmatrix} 1 \\ e^{i\alpha} \end{pmatrix}_X \quad \oint = X_\alpha = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\alpha} \end{pmatrix}_X$$

These are relative phases for X, hence in Z-Y:



— phases —

Thm. Every linear map in \mathbf{FHilb}_2 can be expressed in the language of a pair of complementary observables and the corresponding phases, that is, it can be written down using only red and green decorated spiders.

$$\Lambda^{Z}(\gamma) \circ \Lambda^{X}(\beta) \circ \Lambda^{Z}(\alpha) =$$

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = \mathbf{0} = \mathbf{$$

-MBQC -

New results on resource requirements and complexity of translations between Q-computational models:



R. Duncan and S. Perdrix ICALP'10

quantomatic - Dixon / Duncan / Frot / Kissinger / Merry / Soloviev



http://sites.google.com/site/quantomatic/home

ENVIRONMENT AND CLASSICAL CHANNELS IN CATEGORICAL Q.M.

B.C. and S. Perdrix – CSL'10 – arXiv:1004.1598

Selinger QPL'05: pure <u>cats</u> to mixed <u>cats</u> construction

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C QPL'06: axiomatic account on mixed <u>cats</u> via:

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C and Perdrix '10: interaction of = with \checkmark & \checkmark :

- Decoherence
- Classical channels
- Complex control structure
- Elementary derivation of general protocols

 $\{\stackrel{\triangleq}{\uparrow}\}_A$ is **environment** iff for all (pure) f, g:



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and:

 $\frac{1}{\overline{\overline{z}}} = \bigcup^{\hat{\overline{z}}} \qquad \stackrel{A}{\overline{\overline{z}}} \stackrel{B}{\overline{\overline{z}}} = \stackrel{A \otimes B}{\frac{1}{\overline{\overline{z}}}}$





Prop 1: $\hat{=} = \hat{=}$





Prop 2:



Destructive measurement:

Destructive measurement:

Non-destructive measurement:



1st example: QUANTUM TELEPORTATION













Indeed measurement:



Indeed controlled unitary:



2nd example: QUANTUM KEY DISTRIBUTION

— key distribution —


— key distribution —



STRONG COMPLEMENTARITY

B.C., Ross Duncan, Quanlong Wang Anne Hillebrand **Def.** Strong complementarity := (scaled) bialgebra



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Prop. Strong complementarity \Rightarrow Complementarity

Def. Strong complementarity := (scaled) bialgebra



Prop. Strong complementarity \Rightarrow Complementarity

Conj. Strong complementarity, $\pi/2$ -phases, H-decomp. is complete with respect to stabilizer qubit theory.

R. Duncan and S. Perdrix (2009) *Graph states and the necessity of Euler decomposition*. CiE'09, LNCS 5635. arXiv:0902.0500

Alex Lang and B.C. (2011) *Trichromatic open digraphs for understanding qubits*. QPL'11 proceedings.



Prop. Strong complementarity \Rightarrow Complementarity

Conj. Strong complementarity, $\pi/2$ -phases, H-decomp. is complete with respect to stabilizer qubit theory.

Claim. Strong complementarity is more fundamental as a structural resource than complementarity.

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \circ \sigma \circ \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \circ \sigma \circ \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} = ?$$

























