Towards Scalable Synthesis of Stochastic Control Systems

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Abstract Formal synthesis approaches over stochastic systems have received significant attention in the past few years, in view of their ability to provide provably correct controllers for complex logical specifications in an automated fashion. Examples of complex specifications include properties expressed as formulae in linear temporal logic (LTL) or as automata on infinite strings. A general methodology to synthesize controllers for such properties resorts to symbolic models of the given stochastic systems. Symbolic models are finite abstractions of the given concrete systems with the property that a controller designed on the abstraction can be refined (or implemented) into a controller on the original system. Although the recent development of techniques for the construction of symbolic models has been quite encouraging, the general goal of formal synthesis over stochastic control systems is by no means solved. A fundamental issue with the existing techniques is the known "curse of dimensionality," which is due to the need to discretize state and input sets. Such discretization generally results in an exponential complexity over the number of state and input variables in the concrete system. In this work we propose a novel abstraction technique for incrementally stable stochastic control systems, which does

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not require state-space discretization but only input set discretization, and that can be potentially more efficient (and thus scalable) than existing approaches. We elucidate the effectiveness of the proposed approach by synthesizing a schedule for the coordination of two traffic lights under some safety and fairness requirements for a road traffic model. Further we argue that this 5-dimensional linear stochastic control system cannot be studied with existing approaches based on state-space discretization due to the very large number of generated discrete states.

Keywords Stochastic control systems · Formal controller synthesis · Finite abstractions · Approximate bisimulation

1 Introduction

In the last decade many techniques have been developed providing controllers for control systems (both deterministic and, more recently, stochastic) in a *formal* and *automated* fashion against some complex *logical* specifications. Examples of such specifications include properties expressed as formulae in linear temporal logic (LTL) or as automata on infinite strings [3], and as such they are not tractable by classical techniques for control systems. A general scheme for providing such controllers is by leveraging *symbolic models* of original concrete systems. Symbolic models are discrete abstractions of the original systems in which each symbol represents an aggregate of continuous variables. When such symbolic models exist for the concrete systems, one can leverage the algorithmic machinery for automated synthesis of discrete models [1,14] to automatically synthesize discrete controllers which can be refined to hybrid controllers for the original systems.

The construction of symbolic models for *continuous-time* deterministic systems has been thoroughly investigated in the past few years. This includes results on the construction of approximately bisimilar symbolic models for incrementally stable control systems [15, 18], incrementally stable switched systems [7], and control systems with disturbances [19], non-uniform abstractions of nonlinear systems over a finite-time horizon [26], as well as the construction of sound abstractions based on the convexity of reachable sets [20], feedback refinement relations [21], robustness margins [13], and approximate alternating simulation relation [32]. Recently, there have been some results on the construction of symbolic models for *continuous-time* stochastic control system, without providing a quantitative relationship between abstract and concrete model [12], approximately bisimilar symbolic models for incrementally stable stochastic control systems [31], stochastic switched systems [29], and randomly switched stochastic systems [28], as well as abstractions for unstable stochastic control systems [30].

Note that all the techniques provided in [15,18,7,19,26,20,21,13,32,12,31,28, 30] are fundamentally based on the discretization of continuous states. Therefore, they suffer severely from the curse of dimensionality due to gridding those sets, which is especially irritating for models with high-dimensional state spaces. In this work we propose a novel approach for the construction of approximately bisimilar symbolic models for incrementally stable stochastic control systems not requiring any state set

discretization but only input set discretization. Therefore, it can be potentially more efficient than the proposed approaches in [31] when dealing with higher dimensional stochastic control systems. In particular, we provide a theoretical comparison with the approach in [31] and a simple criterion that helps choosing the most suitable among two approaches (in terms of the sizes of the symbolic models) for a given stochastic control system.

Another advantage of the technique proposed here is that it allows us to construct symbolic models with probabilistic output values, resulting in less conservative symbolic abstractions than those proposed in [31,28,30] that allow for deterministic output values exclusively. We then explain how the proposed symbolic models with probabilistic output values can be used for synthesizing hybrid controllers enforcing logic specifications. The proposed approaches in [29] also provide symbolic models with probabilistic output values and without any state set discretization. However, the results in [29] are for stochastic switched systems rather than stochastic control systems as in this work and they do not provide any intuition behind the control synthesis over symbolic models with probabilistic output values. The effectiveness of the proposed results is illustrated by synthesizing a schedule for the coordination of two traffic lights under some safety and fairness requirements for a model of road traffic which is a 5-dimensional linear stochastic control system. We also show that this example is not amenable to be dealt with the approaches proposed in [31]. Although the main proposed results in this work are for incrementally stable stochastic control systems, the similar results for incrementally stable deterministic control systems can be recovered in the same framework by simply setting the diffusion term to zero.

Alongside the relationship with and extension of [29,31], this paper provides a detailed and extended elaboration of the results first announced in [34], including the proofs of the main results, a detailed discussion on how to deal with probabilistic output values and a generalization of the corresponding result with no requirement on compactness, and finally discussing a new case study on road traffic control.

Similar to this contribution, the results in [23,25] also use sequences of symbols as symbolic states to provide finite abstractions of concrete systems. However, the results in [23,25] apply to discrete-time deterministic control systems, rather than continuous-time stochastic control systems. Furthermore, the results in [25] do not provide a constructive approach for the synthesis of finite abstractions, and the relationships established between abstractions and concrete systems do not necessarily preserve LTL properties. The results in [23] provide asynchronous ℓ -complete finite abstractions which are generally incomparable to the approximate bisimilar finite abstractions proposed in this work.

2 Stochastic Control Systems

2.1 Notation

The identity map over a set *B* is denoted by 1_B . The symbols \mathbb{N} , \mathbb{N}_0 , \mathbb{Z} , \mathbb{R} , \mathbb{R}^+ , and \mathbb{R}_0^+ denote the set of natural, nonnegative integer, integer, real, positive, and nonnegative real numbers, respectively. The symbols I_n , 0_n , and $0_{n \times m}$ denote the identity matrix,

the zero vector, and the zero matrix in $\mathbb{R}^{n \times n}$, \mathbb{R}^n , and $\mathbb{R}^{n \times m}$, respectively. Given a vector $x \in \mathbb{R}^n$, we denote by x_i the *i*-th element of *x*, and by ||x|| the infinity norm of *x*. Given a matrix $P = \{p_{ij}\} \in \mathbb{R}^{n \times n}$, we denote by $\operatorname{Tr}(P) = \sum_{i=1}^{n} p_{ii}$ the trace of *P*. Given a symmetric matrix *A*, we denote by $\lambda_{\min}(A)$ and $\lambda_{\max}(A)$ the minimum and maximum eigenvalues of *A*, respectively. The diagonal set $\Delta \subset \mathbb{R}^n \times \mathbb{R}^n$ is defined as: $\Delta = \{(x, x) \mid x \in \mathbb{R}^n\}$.

The closed ball centered at $u \in \mathbb{R}^m$ with radius λ is defined by $\mathscr{B}_{\lambda}(u) = \{v \in \mathbb{R}^m | ||u-v|| \leq \lambda\}$. A set $B \subseteq \mathbb{R}^m$ is called a *box* if $B = \prod_{i=1}^m [c_i, d_i]$, where $c_i, d_i \in \mathbb{R}$ with $c_i < d_i$ for each $i \in \{1, ..., m\}$. The *span* of a box B is defined as *span*(B) = min $\{|d_i - c_i| \mid i = 1, ..., m\}$. For a box $B \subseteq \mathbb{R}^m$ and $\mu \leq span(B)$, define the μ -approximation $[B]_{\mu} = [\mathbb{R}^m]_{\mu} \cap B$, where $[\mathbb{R}^m]_{\mu} = \{a \in \mathbb{R}^m \mid a_i = k_i \mu, k_i \in \mathbb{Z}, i = 1, ..., m\}$. Remark that $[B]_{\mu} \neq \emptyset$ for any $\mu \leq span(B)$. Geometrically, for any $\mu \in \mathbb{R}^+$ with $\mu \leq span(B)$ and $\lambda \geq \mu$, the collection of sets $\{\mathscr{B}_{\lambda}(p)\}_{p \in [B]_{\mu}}$ is a finite covering of B, i.e. $B \subseteq \bigcup_{p \in [B]_{\mu}} \mathscr{B}_{\lambda}(p)$. We extend the notions of *span* and *approximation* to finite unions of boxes as follows. Let $A = \bigcup_{j=1}^M A_j$, where each A_j is a box. Define *span*(A) = min $\{span(A_j) \mid j = 1, ..., M\}$, and for any $\mu \leq span(A)$, define $[A]_{\mu} = \bigcup_{j=1}^M [A_j]_{\mu}$.

Given a measurable function $f : \mathbb{R}_0^+ \to \mathbb{R}^n$, the (essential) supremum of f is denoted by $||f||_{\infty} := (ess)sup\{||f(t)||, t \ge 0\}$. A continuous function $\gamma : \mathbb{R}_0^+ \to \mathbb{R}_0^+$ is said to belong to class \mathscr{K} if it is strictly increasing and $\gamma(0) = 0$; γ is said to belong to class \mathscr{K} if $\gamma \in \mathscr{K}$ and $\gamma(r) \to \infty$ as $r \to \infty$. A continuous function $\beta : \mathbb{R}_0^+ \times \mathbb{R}_0^+ \to \mathbb{R}_0^+$ is said to belong to class $\mathscr{K} \mathscr{L}$ if, for each fixed *s*, the map $\beta(r,s)$ belongs to class $\mathscr{K} \mathscr{L}$ if, for each fixed nonzero *r*, the map $\beta(r,s)$ is decreasing with respect to *r* and, for each fixed nonzero *r*, the map $\beta(r,s)$ is decreasing with respect to *s* and $\beta(r,s) \to 0$ as $s \to \infty$. We identify a relation $R \subseteq A \times B$ with the map $R : A \to 2^B$ defined by $b \in R(a)$ iff $(a,b) \in R$. Given a relation $R \subseteq A \times B$, R^{-1} denotes the inverse relation defined by $R^{-1} = \{(b,a) \in B \times A : (a,b) \in R\}$. Given a finite sequence *S*, we denote by $\sigma := (S)^{\omega}$ the infinite sequence generated by repeating *S* infinitely, i.e. $\sigma := SSSSS \dots$

2.2 Stochastic control systems

Let $(\Omega, \mathscr{F}, \mathbb{P})$ be a probability space equipped with a filtration $\mathbb{F} = (\mathscr{F}_s)_{s\geq 0}$ satisfying the usual conditions of completeness and right continuity [10, p. 48]. Let $(W_s)_{s\geq 0}$ be a *p*-dimensional \mathbb{F} -adapted Brownian motion.

Definition 1 A stochastic control system Σ is a tuple $\Sigma = (\mathbb{R}^n, \bigcup, \mathcal{U}, f, \sigma)$, where \mathbb{R}^n is the state space, $\bigcup \subseteq \mathbb{R}^m$ is a bounded input set, and

- \mathscr{U} is a subset of the set of all measurable functions from \mathbb{R}^+_0 to U;
- $f: \mathbb{R}^n \times U \to \mathbb{R}^n$ satisfies the following Lipschitz assumption: there exist constants $L_x, L_u \in \mathbb{R}^+$ such that $||f(x, u) f(x', u')|| \le L_x ||x x'|| + L_u ||u u'||$ for all $x, x' \in \mathbb{R}^n$ and all $u, u' \in U$;
- $\sigma : \mathbb{R}^n \to \mathbb{R}^{n \times p}$ satisfies the following Lipschitz assumption: there exists a constant $Z \in \mathbb{R}^+$ such that: $\|\sigma(x) \sigma(x')\| \le Z \|x x'\|$ for all $x, x' \in \mathbb{R}^n$.

A continuous-time stochastic process $\xi : \Omega \times \mathbb{R}^+_0 \to \mathbb{R}^n$ is said to be a *solution* process of Σ if there exists $v \in \mathscr{U}$ satisfying the following stochastic differential equation (SDE) \mathbb{P} -almost surely (\mathbb{P} -a.s.):

$$d\xi = f(\xi, v) dt + \sigma(\xi) dW_t, \tag{1}$$

where f is known as the *drift* and σ as the *diffusion*. We also denote by $\xi_{a\nu}(t)$ the value of the solution process at time $t \in \mathbb{R}_0^+$ under the input curve ν from initial condition $\xi_{a\nu}(0) = a \mathbb{P}$ -a.s., where a is a \mathscr{F}_0 -measurable random variable. Let us emphasize that the solution process is unambiguously determined, since the assumptions on f and σ ensure its existence and uniqueness [16, Theorem 5.2.1, p. 68].

Example 1 For a *linear* stochastic control system $\Sigma = (\mathbb{R}^n, \bigcup, \mathcal{U}, f, \sigma)$, one has f(x, u) := Ax + Bu, for some matrices $A \in \mathbb{R}^{n \times n}$ and $B \in \mathbb{R}^{n \times m}$; and $\sigma(x) := [\sigma_1 x \dots \sigma_p x]$ for some matrices $\sigma_i \in \mathbb{R}^{n \times n}$.

3 Incremental Stability

We recall a stability notion for stochastic control systems, introduced in [31], on which the main results presented in this work rely.

Definition 2 A stochastic control system Σ is incrementally input-to-state stable in the *q*th moment (δ -ISS-M_{*q*}), where $q \ge 1$, if there exist a function $\beta \in \mathscr{KL}$ and a function $\gamma \in \mathscr{K}_{\infty}$ such that for any $t \in \mathbb{R}_{0}^{+}$, any \mathbb{R}^{n} -valued \mathscr{F}_{0} -measurable random variables *a* and *a'*, and any $\upsilon, \upsilon' \in \mathscr{U}$, the following condition is satisfied:

$$\mathbb{E}\left[\left\|\xi_{a\upsilon}(t) - \xi_{a'\upsilon'}(t)\right\|^{q}\right] \le \beta \left(\mathbb{E}\left[\left\|a - a'\right\|^{q}\right], t\right) + \gamma\left(\left\|\upsilon - \upsilon'\right\|_{\infty}\right).$$
(2)

It can be easily verified that a δ -ISS-M_q stochastic control system Σ is δ -ISS [2] in the absence of any noise as in the following:

$$\|\xi_{a\upsilon}(t) - \xi_{a'\upsilon'}(t)\| \le \beta \left(\left\| a - a' \right\|, t \right) + \gamma \left(\left\| \upsilon - \upsilon' \right\|_{\infty} \right), \tag{3}$$

for $a, a' \in \mathbb{R}^n$, some $\beta \in \mathscr{KL}$, and some $\gamma \in \mathscr{K}_{\infty}$.

Remark 1 For linear stochastic control systems Σ , as in Example 1, δ -ISS-M₂ is the same as asymptotic stability in the mean square sense (cf. [33, Lemma 5.1]). However, for nonlinear stochastic control systems, δ -ISS-M_q is a much stronger property than the usual stability in the *q*th moment [9].

Similar to the characterization of δ -ISS in terms of the existence of so-called δ -ISS Lyapunov functions in [2], one can describe δ -ISS-M_q in terms of the existence of so-called δ -ISS-M_q Lyapunov functions, as shown in [31] and defined next.

Definition 3 Consider a stochastic control system Σ and a continuous function V: $\mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}^+_0$ that is twice continuously differentiable on $\{\mathbb{R}^n \times \mathbb{R}^n\} \setminus \Delta$. The function V is called a δ -ISS-M_q Lyapunov function for Σ , where $q \ge 1$, if there exist functions $\underline{\alpha}, \overline{\alpha}, \rho \in \mathcal{K}_{\infty}$, and a constant $\kappa \in \mathbb{R}^+$, such that (i) $\underline{\alpha}$ (resp. $\overline{\alpha}$) is a convex (resp. concave) function;

(ii) for any $x, x' \in \mathbb{R}^n$, $\underline{\alpha}(\|x - x'\|^q) \le V(x, x') \le \overline{\alpha}(\|x - x'\|^q)$;

(iii) for any $x, x' \in \mathbb{R}^n$, $x \neq x'$, and for any $u, u' \in U$,

$$\begin{aligned} \mathscr{L}^{u,u'}V(x,x') &:= \left[\partial_x V \ \partial_{x'}V\right] \begin{bmatrix} f(x,u) \\ f(x',u') \end{bmatrix} + \\ \frac{1}{2} \operatorname{Tr}\left(\begin{bmatrix} \boldsymbol{\sigma}(x) \\ \boldsymbol{\sigma}(x') \end{bmatrix} \begin{bmatrix} \boldsymbol{\sigma}^T(x) \ \boldsymbol{\sigma}^T(x') \end{bmatrix} \begin{bmatrix} \partial_{x,x} V \ \partial_{x',x'}V \\ \partial_{x',x'}V \ \partial_{x',x'}V \end{bmatrix} \right) &\leq -\kappa V(x,x') + \rho(\|u-u'\|), \end{aligned}$$

where $\mathscr{L}^{u,u'}$ is the infinitesimal generator associated to the process $V(\xi, \xi')$ and ξ and ξ' are solution processes of the SDE (1) [16, Section 7.3]. The symbols ∂_x and $\partial_{x,x'}$ denote first- and second-order partial derivatives with respect to x and (x,x'), respectively.

Condition (*ii*) in the above definition implies that the growth rate of functions $\overline{\alpha}$ and $\underline{\alpha}$ is linear. Remark that this condition does not limit the behavior of $\overline{\alpha}$ and $\underline{\alpha}$ to linear ones on a compact subset of \mathbb{R}^n . Observe that condition (*i*) is not required in the context of deterministic control systems for the corresponding δ -ISS Lyapunov functions [2]. The following theorem, borrowed from [31], describes δ -ISS-M_q in terms of the existence of δ -ISS-M_q Lyapunov functions.

Theorem 1 A stochastic control system Σ is δ -ISS- M_q if it admits a δ -ISS- M_q Lyapunov function.

One can resort to available optimization software tools, such as SOSTOOLS [17], to search for appropriate δ -ISS-M_q Lyapunov functions for systems Σ of polynomial type (polynomial drift and diffusion). We refer the interested readers to the results in [31] for the discussion of special instances where these functions can be easily computed.

As an example, for linear stochastic control systems Σ , as in Example 1, one can search for δ -ISS-M_q Lyapunov functions of the form $V(x,x') := \frac{1}{q}((x-x')^T P(x-x'))^{\frac{q}{2}}$, for some positive definite $P \in \mathbb{R}^{n \times n}$ and $q \in \{1,2\}$, by solving the following linear matrix inequality (LMI):

$$A^{T}P + PA + \sum_{i=1}^{p} \sigma_{i}^{T} P \sigma_{i} \preceq -\widehat{\kappa}P,$$
(4)

for some positive constant $\hat{\kappa}$. We refer the interested readers to [31, Corollary 3.5] for more details.

3.1 Noisy and noise-free trajectories

In order to introduce the symbolic models with deterministic output values in Subsection 5.2 (Theorems 2 and 3) for a stochastic control system, we need the following technical result, borrowed from [31]. The following result provides an upper bound on the distance (in the qth moment) between the solution process of Σ and the solution of a derived deterministic control system $\overline{\Sigma}$ obtained by disregarding the diffusion term σ in Σ . From now on, we denote by $\overline{\xi}_{xv}$ the solution of $\overline{\Sigma} = (\mathbb{R}^n, \bigcup, \mathcal{U}, f, 0_{n \times p})^1$, starting from the deterministic initial condition $x = \overline{\xi}_{xv}(0)$ and under the input curve v, which satisfies the ordinary differential equation (ODE) $\dot{\xi}_{xv} = f(\overline{\xi}_{xv}, v)$.

Lemma 1 Consider a stochastic control system Σ such that $f(0_n, 0_m) = 0_n$ and $\sigma(0_n) = 0_{n \times p}$. Suppose that $q \ge 2$ and that there exists a δ -ISS- M_q Lyapunov function V for Σ such that its Hessian is a positive semidefinite matrix in $\mathbb{R}^{2n \times 2n}$ and $\partial_{x,x}V(x,x') \le P$, for any $x, x' \in \mathbb{R}^n$, and some positive semidefinite matrix $P \in \mathbb{R}^{n \times n}$. Then for any $x \in \mathbb{R}^n$ and any $v \in \mathcal{U}$, we have

$$\mathbb{E}\left[\left\|\xi_{x\upsilon}(t) - \overline{\xi}_{x\upsilon}(t)\right\|^{q}\right] \le h_{x}(t),\tag{5}$$

where

$$h_{x}(t) = \underline{\alpha}^{-1}\left(\frac{1}{2}\left\|\sqrt{P}\right\|^{2}n\min\{n,p\}Z^{2}e^{-\kappa t}\int_{0}^{t}\left(\beta\left(\|x\|^{q},s\right)+\gamma\left(\sup_{u\in\mathsf{U}}\{\|u\|\}\right)\right)^{\frac{2}{q}}\mathsf{d}s\right),$$

and Z is the Lipschitz constant, introduced in Definition 1, and β is the \mathcal{KL} function appearing in (2).

It can be readily seen that the nonnegative-valued function h_x tends to zero as $t \to 0, t \to +\infty$, or as $Z \to 0$, and is identically zero if the diffusion term is identically zero (i.e. Z = 0, which is the case for $\overline{\Sigma}$). The interested readers are referred to [31] providing (possibly less conservative) results in line with that of Lemma 1 for (linear) stochastic control systems Σ admitting a specific type of δ -ISS-M_q Lyapunov functions.

4 Systems and Approximate Equivalence Relations

4.1 Systems

We employ the abstract and general notion of "*system*," as introduced in [24], to describe both stochastic control systems and their symbolic models.

Definition 4 A system *S* is a tuple $S = (X, X_0, U, \longrightarrow, Y, H)$, where *X* is a (possibly infinite) set of states, $X_0 \subseteq X$ is a (possibly infinite) set of initial states, *U* is a (possibly infinite) set of inputs, $\longrightarrow \subseteq X \times U \times X$ is a transition relation, *Y* is a set of outputs, and $H : X \to Y$ is an output map. A transition $(x, u, x') \in \longrightarrow$ is also denoted by $x \xrightarrow{u} x'$. For a transition $x \xrightarrow{u} x'$, state x' is called a *u*-successor, or simply a successor, of state *x*. We denote by **Post**_{*u*}(*x*) the set of all *u*-successors of a state *x*. A system *S* is said to be

- *metric*, if the output set Y is equipped with a metric $\mathbf{d}: Y \times Y \to \mathbb{R}_0^+$;

¹ Here, we have abused notation by identifying $0_{n \times p}$ with the map $\sigma : x \to 0_{n \times p} \ \forall x \in \mathbb{R}^n$.

- *finite* (or *symbolic*), if *X* and *U* are finite sets;
- *deterministic*, if for any state $x \in X$ and any input $u \in U$, $|\mathbf{Post}_u(x)| \le 1$. \Box

For technical reasons, we assume that for any $x \in X$, there exists some *u*-successor of *x*, for some $u \in U$ — note that this is always the case for the considered systems in this paper.

Given a system $S = (X, X_0, U, \longrightarrow, Y, H)$ and any initial state $x_0 \in X_0$, a finite state run generated from x_0 is a finite sequence of transitions:

$$x_0 \xrightarrow{u_0} x_1 \xrightarrow{u_1} \cdots \xrightarrow{u_{n-2}} x_{n-1} \xrightarrow{u_{n-1}} x_n, \tag{6}$$

such that $x_i \xrightarrow{u_i} x_{i+1}$ for all $0 \le i < n$. A finite state run can be directly extended to an infinite state run as well. A finite output run is a sequence $\{y_0, y_1, \ldots, y_n\}$ such that there exists a finite state run of the form (6) with $y_i = H(x_i)$, for $i = 0, \ldots, n$. A finite output run can also be directly extended to an infinite output run as well.

4.2 Relations among systems

We recall the notion of approximate (bi)simulation relation, introduced in [6], which is crucial when analyzing or synthesizing controllers for deterministic systems.

Definition 5 Consider two metric systems $S_a = (X_a, X_{a0}, U_a, \xrightarrow{a}, Y_a, H_a)$ and $S_b = (X_b, X_{b0}, U_b, \xrightarrow{b}, Y_b, H_b)$ with the same output sets $Y_a = Y_b$ and metric **d**. For $\varepsilon \in \mathbb{R}^+_0$, a relation $R \subseteq X_a \times X_b$ is said to be an ε -approximate simulation relation from S_a to S_b if, for all $(x_a, x_b) \in R$, the following two conditions are satisfied:

(i) $\mathbf{d}(H_a(x_a), H_b(x_b)) \leq \varepsilon$; (ii) for any $x_a \xrightarrow{u_a} x'_a \in \underbrace{\longrightarrow}_a$ there exists $x_b \xrightarrow{u_b} x'_b \in \underbrace{\longrightarrow}_b$ such that $(x'_a, x'_b) \in R$.

A relation $R \subseteq X_a \times X_b$ is called an ε -approximate bisimulation relation between S_a and S_b if R is an ε -approximate simulation relation from S_a to S_b and R^{-1} is an ε -approximate simulation relation from S_b to S_a .

System S_a is ε -approximately simulated by S_b , or $S_b \varepsilon$ -approximately simulates S_a , denoted by $S_a \preceq_{\mathscr{S}}^{\varepsilon} S_b$, if there exists an ε -approximate simulation relation R from S_a to S_b such that:

- for every $x_{a0} \in X_{a0}$, there exists $x_{b0} \in X_{b0}$ such that $(x_{a0}, x_{b0}) \in R$.

System S_a is ε -approximately bisimilar to S_b , denoted by $S_a \cong_{\mathscr{S}}^{\varepsilon} S_b$, if there exists an ε -approximate bisimulation relation R between S_a and S_b such that:

- for every $x_{a0} \in X_{a0}$, there exists $x_{b0} \in X_{b0}$ such that $(x_{a0}, x_{b0}) \in R$;
- for every $x_{b0} \in X_{b0}$, there exists $x_{a0} \in X_{a0}$ such that $(x_{a0}, x_{b0}) \in R$.

5 Symbolic Models for Stochastic Control Systems

5.1 Describing stochastic control systems as metric systems

In order to show the main results of the paper, we use the notion of system introduced in Definition 4 to abstractly represent a stochastic control system. More precisely, given a stochastic control system Σ , we define an associated metric system $S(\Sigma) = (X, X_0, U, \longrightarrow, Y, H)$, where:

- X is the set of all Rⁿ-valued random variables defined on the probability space (Ω, F, P);
- X_0 is a subset of the set of \mathbb{R}^n -valued random variables that are measurable over \mathscr{F}_0 ;
- $U = \mathscr{U};$
- $x \xrightarrow{\upsilon} x'$ if x and x' are measurable in \mathscr{F}_t and $\mathscr{F}_{t+\tau}$, respectively, for some $t \in \mathbb{R}^+_0$ and $\tau \in \mathbb{R}^+$, and there exists a solution process $\xi : \Omega \times \mathbb{R}^+_0 \to \mathbb{R}^n$ of Σ satisfying $\xi(t) = x$ and $\xi_{x\upsilon}(\tau) = x'$ P-a.s.;
- Y = X;
- $H = 1_X.$

We assume that the output set *Y* is equipped with the *q*th moment metric $\mathbf{d}(y, y') = (\mathbb{E}[||y - y'||^q])^{\frac{1}{q}}$, for any $y, y' \in Y$ and some $q \ge 1$. One can readily observe that the set of states and inputs of $S(\Sigma)$ are uncountable and that $S(\Sigma)$ is a deterministic system in the sense of Definition 4, since (cf. Subsection 2.2) the solution process of Σ is uniquely determined. For the case of deterministic control system $\overline{\Sigma}$, one obtains $S(\overline{\Sigma}) = (X, X_0, U, \longrightarrow, Y, H)$, where $X = \mathbb{R}^n, X_0$ is a subset of $\mathbb{R}^n, U = \mathscr{U}, x \xrightarrow{v} x'$ iff $x' = \overline{\xi}_{xv}(\tau)$ for some $\tau \in \mathbb{R}^+, Y = X, H = 1_X$, and the metric on the output set reduces to the natural infinity one: $\mathbf{d}(y, y') = ||y - y'||$, for any $y, y' \in Y$.

Since the concrete system $S(\Sigma)$ is uncountably infinite, it does not allow for a straightforward symbolic controller synthesis with the discrete techniques in the literature [1,14]. We are thus interested in finding a finite abstract system that is (bi)similar to the concrete one $S(\Sigma)$. In order to discuss approximate (bi)simulation relations between two metric systems, they have to share the output space (cf. Definition 5). System $S(\Sigma)$ inherits a classical trace-based semantics [3] (cf. definition of output run after (6)), however the outputs of $S(\Sigma)$ (and necessarily those of any approximately (bi)similar one) are random variables. This fact is especially important due to the metric **d** that the output set is endowed with: for any deterministic point one can always find a non-degenerate random variable that is as close as desired to the original point in the metric **d**.

To further elaborate the discussion in the previous paragraph, let us consider the following example. Let $A \subset \mathbb{R}^n$ be a set (of deterministic points). Consider a safety problem, formulated as the satisfaction of the LTL formula² $\Box \varphi_A$, where φ_A is a label (or proposition) characterising the set *A*. Suppose that over the abstract system we are able to synthesize a control strategy that makes an output run of the abstraction satisfy

 $^{^2}$ We refer the interested readers to [3, Subsection 5.1.2] for the formal trace-based semantic of LTL formulae.

 $\Box \varphi_A$, which means that the output run of the abstraction is always inside the set A. Although the run would in general be consisting of random variables $y: \Omega \to \mathbb{R}^n$, the fact that $y \in A$ means that, with a slight abuse of notation, $y : \hat{\Omega} \to \{y\}$ where $\hat{\Omega} \subset \Omega$ and $|\hat{\Omega}| = 1$. Hence, y has a Dirac probability distribution centered at y, that is $y \in Y$ is a degenerate random variable that can be identified with a point in $A \subset \mathbb{R}^n \subset Y$. Since any deterministic point can be regarded as a random variable with a Dirac probability distribution centered at that point, \mathbb{R}^n can be embedded in Y, which we denote as $\mathbb{R}^n \subset Y$ with a slight abuse of notation. As a result, satisfying $\Box \varphi_A$ precisely means that the output run of the abstraction indeed stays in the set $A \subset \mathbb{R}^n$ forever. On the other hand, suppose that the original system is ε -approximate bisimilar to the abstraction. If we want to interpret the result $\Box \varphi_A$ obtained over the abstraction, we can guarantee that the corresponding output run of the original system satisfies $\Box \varphi_{A_{e}}$, that is any output y of the run of the original system is within ε d-distance from the set A: $\mathbf{d}(y,A) = \inf_{a \in A} \mathbf{d}(y,a) \leq \varepsilon$. Although the original set $A \subset Y$ is a subset of $\mathbb{R}^n \subset Y$, its ε -inflation $A_{\varepsilon} = \{y \in Y : \mathbf{d}(y,A) \leq \varepsilon\}$ is not a subset of \mathbb{R}^n anymore and hence contains non-degenerate random variables. In particular, $A_{\varepsilon} \neq \{y \in \mathbb{R}^n :$ $\inf_{a \in A} \|y - a\| \le \varepsilon$ and is in fact bigger than the latter set of deterministic points. As a result, although satisfying $\Box \varphi_{A_F}$ does not necessarily mean that a trajectory of Σ always stays within some deterministic set, it means that the associated random variables always belong to A_{ε} and, hence, are close to the deterministic set A with respect to the *q*th moment metric.

We are now able to provide two versions of finite abstractions: one whose outputs are always deterministic points – that is degenerate random variables, elements of $\mathbb{R}^n \subset Y$, and one whose outputs can be non-degenerate random variables. Recall, however, that in both cases the output set is still the whole *Y* and the semantics is the same as for the original system $S(\Sigma)$.

5.2 Main results

This subsection contains the main contributions of the paper. We show that for any δ -ISS- M_q (resp. δ -ISS) stochastic control system Σ (resp. deterministic control system $\overline{\Sigma}$), and for any precision level $\varepsilon \in \mathbb{R}^+$, we can construct a finite system that is ε -approximate bisimilar to Σ (resp. $\overline{\Sigma}$) without any state set discretization. The results in this subsection rely on additional assumptions on the model Σ that are described next. We restrict our attention to stochastic control systems Σ with input sets U that are assumed to be finite unions of boxes (cf. Subsection 2.1). We further restrict our attention to sampled-data stochastic control systems, where input curves belong to set \mathscr{U}_{τ} , which contains exclusively curves that are constant over intervals of length $\tau \in \mathbb{R}^+$, i.e.

$$\mathscr{U}_{\tau} = \Big\{ \upsilon \in \mathscr{U} \mid \upsilon(t) = \upsilon((k-1)\tau), t \in [(k-1)\tau, k\tau[, k \in \mathbb{N}] \Big\}.$$

Let us denote by $S_{\tau}(\Sigma)$ a sub-system of $S(\Sigma)$ obtained by selecting those transitions of $S(\Sigma)$ corresponding to solution processes of duration τ and to control inputs in \mathcal{U}_{τ} . This can be seen as the sampled version of Σ . More precisely, given a stochastic control system Σ and the corresponding metric system $S(\Sigma)$, we define a new associated metric system

$$S_{\tau}(\Sigma) = \left(X_{\tau}, X_{\tau 0}, U_{\tau}, \xrightarrow{\tau}, Y_{\tau}, H_{\tau}\right),$$

where $X_{\tau} = X$, $X_{\tau 0} = X_0$, $U_{\tau} = \mathscr{U}_{\tau}$, $Y_{\tau} = Y$, $H_{\tau} = H$, and $x_{\tau} \xrightarrow{v_{\tau}} x'_{\tau} \in \xrightarrow{\tau}$ if x_{τ} and x'_{τ} are measurable, respectively, in $\mathscr{F}_{k\tau}$ and $\mathscr{F}_{(k+1)\tau}$ for some $k \in \mathbb{N}_0$, and there exists a solution process $\xi : \Omega \times \mathbb{R}^+_0 \to \mathbb{R}^n$ of Σ satisfying $\xi(k\tau) = x_{\tau}$ and $\xi_{x_{\tau}v_{\tau}}(\tau) = x'_{\tau}$ \mathbb{P} -a.s..

Similarly, one can define $S_{\tau}(\overline{\Sigma})$ as the time discretization of $\overline{\Sigma}$. Observe that a finite state run $x_0 \xrightarrow{v_0} x_1 \xrightarrow{v_1} x_1 \xrightarrow{v_{N-1}} x_N$ of $S_{\tau}(\Sigma)$, where $v_{i-1} \in \mathscr{U}_{\tau}$ and $x_i = \xi_{x_{i-1}v_{i-1}}(\tau)$ P-a.s. for i = 1, ..., N, captures the solution process of Σ at times $t = 0, \tau, ..., N\tau$, started from the initial condition x_0 and resulting from a control input v obtained by the concatenation of the input curves v_{i-1} (i.e. $v(t) = v_{i-1}(t)$ for any $t \in [(i-1)\tau, i\tau[)$, for i = 1, ..., N.

Now, we have all the ingredients to introduce two fully symbolic systems for the concrete model Σ . Consider a stochastic control system Σ and a tuple $q = (\tau, \mu, N, x_s)$ of parameters, where τ is the sampling time, μ is the input set quantization, $N \in \mathbb{N}$ is a *temporal horizon*, and $x_s \in \mathbb{R}^n$ is a *source state*. Given Σ and q, let us introduce the following two symbolic systems:

$$\begin{split} S_{\mathsf{q}}(\Sigma) &= (X_{\mathsf{q}}, X_{\mathsf{q}0}, U_{\mathsf{q}}, \xrightarrow{\mathsf{q}}, Y_{\mathsf{q}}, H_{\mathsf{q}}), \\ \overline{S}_{\mathsf{q}}(\Sigma) &= (X_{\mathsf{q}}, X_{\mathsf{q}0}, U_{\mathsf{q}}, \xrightarrow{\mathsf{q}}, Y_{\mathsf{q}}, \overline{H}_{\mathsf{q}}), \end{split}$$

where

-
$$X_{q} = \{ (u_{1}, \dots, u_{N}) \in \overbrace{[U]_{\mu} \times \dots \times [U]_{\mu}}^{N \text{ times}} \};$$

- $X_{q0} = X_{q};$
- $U_{q} = [U]_{\mu};$
- $x_{q} \xrightarrow{u_{q}} x'_{q}$, where $x_{q} = (u_{1}, u_{2}, \dots, u_{N})$, if and only if $x'_{q} = (u_{2}, \dots, u_{N}, u_{q});$

- Y_q is the set of all \mathbb{R}^n -valued random variables defined on the probability space $(\Omega, \mathscr{F}, \mathbb{P});$

-
$$H_{q}(x_{q}) = \xi_{x_{s}x_{q}}(N\tau) \left(\overline{H}_{q}(x_{q}) = \overline{\xi}_{x_{s}x_{q}}(N\tau)\right).$$

Note that the transition relation in $S_q(\Sigma)$ admits a compact representation in the form of a *shift operator*. We have abused notation by identifying $u_q \in [U]_{\mu}$ with the constant input curve with domain $[0, \tau[$ and value u_q , and by identifying $x_q \in [U]_{\mu}^N$ with the concatenation of N control inputs $u_i \in [U]_{\mu}$ (i.e. $x_q(t) = u_i$ for any $t \in [(i - 1)\tau, i\tau[)$ for i = 1, ..., N. Notice that the proposed abstraction $S_q(\Sigma)$ (resp. $\overline{S}_q(\Sigma)$) is a deterministic system in the sense of Definition 4. We point out that H_q and \overline{H}_q are mappings from a deterministic point x_q to the random variable $\xi_{x_s x_q}(N\tau)$ and to the one with a Dirac probability distribution centered at $\overline{\xi}_{x_s x_q}(N\tau)$, respectively. Finally,



Fig. 1 Example of abstraction $S_q(\Sigma)$ with N = 3 and $U_q = \{0, 1\}$. The lower part of the states are their output values. Initial states are targets of sourceless arrows.

in the case of a deterministic control system $\overline{\Sigma}$, one obtains the symbolic system $\overline{S}_q(\overline{\Sigma}) = (X_q, X_{q0}, U_q, \xrightarrow{q}, Y_q, \overline{H}_q)$, where $X_q, X_{q0}, U_q, \xrightarrow{q}$, and \overline{H}_q are the same as before, but the output set reduces to $Y_q = \mathbb{R}^n$.

The main idea behind the definitions of symbolic models $S_q(\Sigma)$ and $\overline{S}_q(\Sigma)$ hinges on the δ -ISS-M_q property. Given an input $v \in \mathcal{U}$, all solution processes of Σ under the input v forget the mismatch between their initial conditions and converge to each other with respect to the *q*th moment metric. Therefore, the longer the applied inputs are, the less relevant is the mismatch between initial conditions. Then, the fundamental idea of the introduced abstractions consists in taking the *N* consequitive applied inputs as the state of the symbolic model.

The control synthesis over $\overline{S}_q(\Sigma)$ (resp. $\overline{S}_q(\overline{\Sigma})$) is simple as the outputs are deterministic points, whereas for $S_q(\Sigma)$ it is perhaps less intuitive. Hence, we discuss it in more details later in Subsection 5.3.

Example 2 An example of an abstraction $S_q(\Sigma)$ with N = 3 and $U_q = \{0, 1\}$ is depicted in Figure 1, where the initial states are shown as targets of sourceless arrows. Regardless of the size of the state set and of its dimension, $S_q(\Sigma)$ only has eight possible states, namely:

$$X_{q} = \{(0,0,0), (0,0,1), (0,1,0), (0,1,1), (1,0,0), (1,0,1), (1,1,0), (1,1,1)\}$$

In order to obtain some of the main results of this work, we raise an assumption on the δ -ISS-M_{*q*} Lyapunov function *V* which we will work with, as follows:

$$|V(x,y) - V(x,z)| \le \widehat{\gamma}(||y - z||), \tag{7}$$

for any $x, y, z \in \mathbb{R}^n$, and for some concave function $\widehat{\gamma} \in \mathscr{K}_{\infty}$. The function $\widehat{\gamma}$ in (7) can be easily computed as long as one is interested to work in a compact subset of \mathbb{R}^n . Particularly, for all $x, y, z \in D$, where $D \subset \mathbb{R}^n$ is compact, one can apply the mean value theorem to the function $y \to V(x, y)$ to get

$$|V(x,y) - V(x,z)| \le \widehat{\gamma}(||y-z||), \text{ where } \widehat{\gamma}(r) = \left(\max_{x,y\in \mathsf{D}\setminus\Delta}\left\|\frac{\partial V(x,y)}{\partial y}\right\|\right)r.$$

In particular, for the δ -ISS-M₁ Lyapunov function $V(x,x') := \sqrt{(x-x')^T P(x-x')}$, for some positive definite matrix $P \in \mathbb{R}^{n \times n}$ and for all $x, x' \in \mathbb{R}^n$, one obtains $\widehat{\gamma}(r) = \sqrt{\lambda_{\max}(P)r}$ due to the triangle inequality, which satisfies (7) globally on \mathbb{R}^n . Remark that for deterministic control systems, the concavity assumption of $\widehat{\gamma}$ is not required.

Before providing the main results of the paper, we need the following technical results.

Lemma 2 Consider a stochastic control system Σ , admitting a δ -ISS- M_q Lyapunov function V, and consider its corresponding symbolic model $\overline{S}_q(\Sigma)$. We have that

$$\eta \leq \left(\underline{\alpha}^{-1}\left(\mathrm{e}^{-\kappa N\tau}\max_{u_{\mathsf{q}}\in U_{\mathsf{q}}}V\left(\overline{\xi}_{x_{s}u_{\mathsf{q}}}(\tau), x_{s}\right)\right)\right)^{1/q},\tag{8}$$

where

$$\eta := \max_{\substack{u_{\mathsf{q}} \in U_{\mathsf{q}}, x_{\mathsf{q}} \in \mathbf{X}_{\mathsf{q}} \\ x'_{\mathsf{q}} \in \mathbf{Post}_{u_{\mathsf{q}}}(x_{\mathsf{q}})}} \left\| \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}\left(x'_{\mathsf{q}}\right) \right\|.$$
(9)

The proof of Lemma 2 is provided in the Appendix. The next lemma provides similar result as the one in Lemma 2, but without explicitly using any Lyapunov function.

Lemma 3 Consider a δ -ISS- M_q stochastic control system Σ and its corresponding symbolic model $\overline{S}_q(\Sigma)$. We have:

$$\eta \leq \left(\beta \left(\max_{u_{\mathsf{q}} \in U_{\mathsf{q}}} \left\| \overline{\xi}_{x_{s}u_{\mathsf{q}}}(\tau) - x_{s} \right\|^{q}, N\tau\right)\right)^{1/q},\tag{10}$$

where η is given in (9) and β is the \mathscr{KL} function appearing in (2).

The proof of Lemma 3 is provided in the Appendix. The next two lemmas provide similar results as Lemmas 2 and 3, but by using the symbolic model $S_q(\Sigma)$ rather than $\overline{S}_q(\Sigma)$.

Lemma 4 Consider a stochastic control system Σ , admitting a δ -ISS- M_q Lyapunov function V, and consider its corresponding symbolic model $S_q(\Sigma)$. One has:

$$\widehat{\eta} \leq \left(\underline{\alpha}^{-1} \left(e^{-\kappa N \tau} \max_{u_{\mathsf{q}} \in U_{\mathsf{q}}} \mathbb{E} \left[V \left(\xi_{x_{s}u_{\mathsf{q}}}(\tau), x_{s} \right) \right] \right) \right)^{1/q}, \tag{11}$$

where

$$\widehat{\eta} := \max_{\substack{u_{\mathsf{q}} \in U_{\mathsf{q}}, x_{\mathsf{q}} \in X_{\mathsf{q}} \\ x'_{\mathsf{q}} \in \mathsf{Post}_{u_{\mathsf{q}}}(x_{\mathsf{q}})}} \mathbb{E}\left[\left\| \xi_{H_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - H_{\mathsf{q}}\left(x'_{\mathsf{q}}\right) \right\| \right].$$
(12)

Proof The proof is similar to the one of Lemma 2 and can be shown by using convexity of $\underline{\alpha}$ and Jensen inequality [16].

Lemma 5 Consider a δ -ISS- M_q stochastic control system Σ and its corresponding symbolic model $S_q(\Sigma)$. We have:

$$\widehat{\eta} \le \left(\beta \left(\max_{u_{\mathsf{q}} \in U_{\mathsf{q}}} \mathbb{E}\left[\left\|\xi_{x_{s}u_{\mathsf{q}}}(\tau) - x_{s}\right\|^{q}\right], N\tau\right)\right)^{1/q},\tag{13}$$

where $\widehat{\eta}$ is given in (12) and β is the \mathscr{KL} function appearing in (2).

Proof The proof is similar to the one of Lemma 3 and can be shown by using Jensen inequality [16]. \Box

Remark 2 It can be readily verified that by choosing *N* sufficiently large, η and $\hat{\eta}$ can be made arbitrarily small. One can as well try to reduce the upper bound for η (in (8) for example) by selecting the source state x_s as follows:

$$x_s = \arg\min_{x \in \mathbb{R}^n} \max_{u_n \in U_n} V(\xi_{xu_q}(\tau), x).$$
(14)

We can now present the first main result of the paper, which relates the existence of a δ -ISS-M_q Lyapunov function to the construction of an approximately bisimilar symbolic model.

Theorem 2 Consider a stochastic control system Σ with $f(0_n, 0_m) = 0_n$ and $\sigma(0_n) = 0_{n \times p}$, admitting a δ -ISS- M_q Lyapunov function V, of the form of the one explained in Lemma 1, such that (7) holds for some concave $\widehat{\gamma} \in \mathscr{K}_{\infty}$. Let η be given by (9). For any $\varepsilon \in \mathbb{R}^+$ and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$e^{-\kappa\tau}\underline{\alpha}(\varepsilon^{q}) + \frac{1}{\mathsf{e}\kappa}\rho(\mu) + \widehat{\gamma}\Big((h_{x_{s}}((N+1)\tau))^{\frac{1}{q}} + \eta\Big) \leq \underline{\alpha}(\varepsilon^{q}), \quad (15)$$

the relation (cf. Definition 5)

$$R = \left\{ (x_{\tau}, x_{\mathsf{q}}) \in X_{\tau} \times X_{\mathsf{q}} \mid \mathbb{E} \left[V \left(x_{\tau}, \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}) \right) \right] \leq \underline{\alpha} \left(\varepsilon^{q} \right) \right\}$$

is an ε -approximate bisimulation relation between $\overline{S}_q(\Sigma)$ and $S_\tau(\Sigma)$.

The proof can be found in the Appendix. By choosing *N* sufficiently large and using the results in Lemmas 1 and 2, one can enforce $h_{x_s}((N+1)\tau)$ and η in (15) to be sufficiently small. Hence, for a given precision $\varepsilon \in \mathbb{R}_0^+$, there always exists a sufficiently small value of μ and a large value of *N*, such that the condition in (15) is satisfied. A result similar as that in Theorem 2 can be recovered for a δ -ISS deterministic control system $\overline{\Sigma}$, as provided in the following corollary.

Corollary 1 Consider a deterministic control system $\overline{\Sigma}$ admitting a δ -ISS Lyapunov function V such that (7) holds for some $\widehat{\gamma} \in \mathscr{K}_{\infty}$. Let η be given by (9). For any $\varepsilon \in \mathbb{R}^+$ and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$e^{-\kappa\tau}\underline{\alpha}(\varepsilon) + \frac{1}{e\kappa}\rho(\mu) + \widehat{\gamma}(\eta) \le \underline{\alpha}(\varepsilon), \qquad (16)$$

the relation

$$R = \left\{ (x_{\tau}, x_{\mathsf{q}}) \in X_{\tau} \times X_{\mathsf{q}} \mid V \left(x_{\tau}, \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}) \right) \leq \underline{\alpha}(\varepsilon) \right\}$$

is an ε *-approximate bisimulation relation between* $\overline{S}_{\mathsf{q}}(\overline{\Sigma})$ *and* $S_{\tau}(\overline{\Sigma})$.

The proof is similar to the one of Theorem 2. In order to mitigate the conservativeness that might rise from using Lyapunov functions, the next theorem provides a result that is similar to the one of Theorem 2, which is however not obtained by explicit use of δ -ISS-M_q Lyapunov functions, but by using functions β and γ as in (2).

Theorem 3 Consider a δ -ISS- M_q stochastic control system Σ , satisfying the result of Lemma 1. Let η be given by (9). For any $\varepsilon \in \mathbb{R}^+$, and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$(\beta \left(\varepsilon^{q}, \tau\right) + \gamma(\mu))^{\frac{1}{q}} + (h_{x_{s}}((N+1)\tau))^{\frac{1}{q}} + \eta \leq \varepsilon,$$
(17)

the relation

$$R = \left\{ (x_{\tau}, x_{\mathsf{q}}) \in X_{\tau} \times X_{\mathsf{q}} \mid \left(\mathbb{E} \left[\left\| x_{\tau} - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}) \right\|^{q} \right] \right)^{\frac{1}{q}} \le \varepsilon \right\}$$

is an ε -approximate bisimulation relation between $\overline{S}_{q}(\Sigma)$ and $S_{\tau}(\Sigma)$.

The proof can be found in the Appendix. By choosing N sufficiently large and using the results in Lemmas 1 and 3, one can force $h_{x_s}((N+1)\tau)$ and η in (17) to be sufficiently small. Hence, for a given precision ε , there always exist a sufficiently large value of τ and N and a small enough value of μ such that the condition in (17) is satisfied. However, unlike the result in Theorem 2, here for a given fixed sampling time τ , one may not find any values of N and μ satisfying (17) because the quantity $(\beta (\varepsilon^q, \tau))^{\frac{1}{q}}$ may be larger than ε . The symbolic model $\overline{S}_q(\Sigma)$, computed using the parameter q provided in Theorem 3 (whenever existing), is likely to have fewer states than the one computed using the parameter q provided in Theorem 2 – a similar fact has been experienced in the first example in [31]. A result similar to the one in Theorem 3 can be fully recovered for a δ -ISS deterministic control system $\overline{\Sigma}$, as provided in the following corollary.

Corollary 2 Consider a δ -ISS deterministic control system $\overline{\Sigma}$. Let η be given by (9). For any $\varepsilon \in \mathbb{R}^+$, and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq span(U)$ and³

$$\beta(\varepsilon,\tau) + \gamma(\mu) + \eta \le \varepsilon, \tag{18}$$

the relation

$$R = \left\{ (x_{\tau}, x_{\mathsf{q}}) \in X_{\tau} \times X_{\mathsf{q}} \mid \left\| x_{\tau} - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}) \right\| \leq \varepsilon \right\}$$

is an ε -approximate bisimulation relation between $\overline{S}_{q}(\overline{\Sigma})$ and $S_{\tau}(\overline{\Sigma})$.

The proof is similar to the one of Theorem 3. The next two theorems provide results that are similar to those of Theorems 2 and 3, but by using the symbolic model $S_q(\Sigma)$ rather than $\overline{S}_q(\Sigma)$.

³ Here, β and γ are the \mathscr{KL} and \mathscr{K}_{∞} functions, respectively, appearing in (3).

Theorem 4 Consider a stochastic control system Σ , admitting a δ -ISS- M_q Lyapunov function V such that (7) holds for some concave $\widehat{\gamma} \in \mathscr{K}_{\infty}$. Let $\widehat{\eta}$ be given by (12). For any $\varepsilon \in \mathbb{R}^+$ and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$e^{-\kappa\tau}\underline{\alpha}(\varepsilon^{q}) + \frac{1}{e\kappa}\rho(\mu) + \widehat{\gamma}(\widehat{\eta}) \leq \underline{\alpha}(\varepsilon^{q}),$$
(19)

the relation

$$R = \{(x_{\tau}, x_{\mathsf{q}}) \in X_{\tau} \times X_{\mathsf{q}} \mid \mathbb{E}\left[V(x_{\tau}, H_{\mathsf{q}}(x_{\mathsf{q}}))\right] \leq \underline{\alpha}\left(\varepsilon^{q}\right)\}$$

is an ε -approximate bisimulation relation between $S_q(\Sigma)$ and $S_\tau(\Sigma)$.

The proof is similar to the one of Theorem 2.

Theorem 5 Consider a δ -ISS- M_q stochastic control system Σ . Let $\hat{\eta}$ be given by (12). For any $\varepsilon \in \mathbb{R}^+$, and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq span(U)$ and

$$(\boldsymbol{\beta}\left(\boldsymbol{\varepsilon}^{q},\tau\right)+\boldsymbol{\gamma}(\boldsymbol{\mu}))^{\frac{1}{q}}+\widehat{\boldsymbol{\eta}}\leq\boldsymbol{\varepsilon},\tag{20}$$

the relation

$$R = \left\{ (x_{\tau}, x_{\mathsf{q}}) \in X_{\tau} \times X_{\mathsf{q}} \mid (\mathbb{E} \left[\|x_{\tau} - H_{\mathsf{q}}(x_{\mathsf{q}})\|^{q} \right])^{\frac{1}{q}} \le \varepsilon \right\}$$

is an ε -approximate bisimulation relation between $S_q(\Sigma)$ and $S_{\tau}(\Sigma)$.

The proof is similar to the one of Theorem 3.

Remark 3 The symbolic model $S_q(\Sigma)$, proposed in Theorem 4 (resp. Theorem 5), has fewer (or at most equal number of) states than the symbolic model $\overline{S}_q(\Sigma)$, proposed in Theorem 2 (resp. Theorem 3) while having the same precision. However, the states in $S_q(\Sigma)$ have probabilistic output values, rather than deterministic ones, which is likely to make control synthesis procedures more involved (cf. Subsection 5.3).

Remark 4 Although we assume that the set U is infinite, Theorems 2, 3, 4, and 5 and Corollaries 1 and 2 still hold when the set U is finite, with the following modifications. First, the systems Σ and $\overline{\Sigma}$ are required to satisfy properties (2) and (3), respectively, for v = v'. Second, take $U_q = U$ in the definitions of $\overline{S}_q(\Sigma)$ (resp. $S_q(\Sigma)$) and $\overline{S}_q(\overline{\Sigma})$. Finally, in the conditions (15), (16), (17), (18), (19), and (20) set $\mu = 0$.

Finally, we establish the results on the existence of symbolic models $\overline{S}_q(\Sigma)$ (resp. $S_q(\Sigma)$) and $\overline{S}_q(\overline{\Sigma})$ such that $\overline{S}_q(\Sigma) \cong^{\varepsilon}_{\mathscr{S}} S_{\tau}(\Sigma)$ (resp. $S_q(\Sigma) \cong^{\varepsilon}_{\mathscr{S}} S_{\tau}(\Sigma)$) and $\overline{S}_q(\overline{\Sigma}) \cong^{\varepsilon}_{\mathscr{S}} S_{\tau}(\overline{\Sigma})$.

Theorem 6 Consider the results in Theorem 2. If we select

$$X_{\tau 0} = \left\{ x \in \mathbb{R}^n | \left\| x - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0}) \right\| \le \left(\overline{\alpha}^{-1} \left(\underline{\alpha} \left(\varepsilon^q \right) \right) \right)^{\frac{1}{q}}, \exists x_{\mathsf{q}0} \in X_{\mathsf{q}0} \right\},$$

then we have $\overline{S}_{q}(\Sigma) \cong_{\mathscr{S}}^{\varepsilon} S_{\tau}(\Sigma)$.

Proof We start by proving that $S_{\tau}(\Sigma) \preceq_{\mathscr{S}}^{\varepsilon} \overline{S}_{q}(\Sigma)$. For every $x_{\tau 0} \in X_{\tau 0}$ there always exists $x_{q0} \in X_{q0}$ such that $||x_{\tau 0} - \overline{H}_{q}(x_{q0})|| \leq (\overline{\alpha}^{-1}(\underline{\alpha}(\varepsilon^{q})))^{\frac{1}{q}}$. Then,

$$\mathbb{E}\left[V\left(x_{\tau 0}, \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0})\right)\right] = V\left(x_{\tau 0}, \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0})\right) \leq \overline{\alpha}(\|x_{\tau 0} - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0})\|^{q}) \leq \underline{\alpha}\left(\varepsilon^{q}\right),$$

since $\overline{\alpha}$ is a \mathscr{K}_{∞} function. Hence, $(x_{\tau 0}, x_{q 0}) \in R$ implying that $S_{\tau}(\Sigma) \preceq_{\mathscr{S}}^{\varepsilon} \overline{S}_{q}(\Sigma)$. In a similar way, we can show that $\overline{S}_{q}(\Sigma) \preceq_{\mathscr{S}}^{\varepsilon} S_{\tau}(\Sigma)$ which completes the proof. \Box

The next theorem provides a similar result in line with the one of previous theorem, but by using a different relation.

Theorem 7 Consider the results in Theorem 3. If we select

$$X_{\tau 0} = \left\{ x \in \mathbb{R}^n \mid \left\| x - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0}) \right\| \le \varepsilon, \ \exists x_{\mathsf{q}0} \in X_{\mathsf{q}0} \right\},\$$

then we have $\overline{S}_{q}(\Sigma) \cong_{\mathscr{S}}^{\varepsilon} S_{\tau}(\Sigma)$.

Proof We start by proving that $S_{\tau}(\Sigma) \preceq_{\mathscr{S}}^{\varepsilon} \overline{S}_{q}(\Sigma)$. For every $x_{\tau 0} \in X_{\tau 0}$ there always exists $x_{q0} \in X_{q0}$ such that $||x_{\tau 0} - \overline{H}_{q}(x_{q0})|| \leq \varepsilon$ and $(\mathbb{E}[||x_{\tau 0} - \overline{H}_{q}(x_{q0})||^{q}])^{\frac{1}{q}} \leq \varepsilon$. Hence, $(x_{\tau 0}, x_{q0}) \in R$ implying that $S_{\tau}(\Sigma) \preceq_{\mathscr{S}}^{\varepsilon} \overline{S}_{q}(\Sigma)$. In a similar way, we can show that $\overline{S}_{q}(\Sigma) \preceq_{\mathscr{S}}^{\varepsilon} S_{\tau}(\Sigma)$ which completes the proof.

The next two corollaries provide similar results as the ones of Theorems 6 and 7, but for deterministic control systems $\overline{\Sigma}$.

Corollary 3 Consider the results in Corollary 1. If we select

$$X_{\tau 0} = \left\{ x \in \mathbb{R}^n | \left\| x - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0}) \right\| \le \left(\overline{\alpha}^{-1}\left(\underline{\alpha}\left(\varepsilon\right)\right) \right), \exists x_{\mathsf{q}0} \in X_{\mathsf{q}0} \right\},$$

then we have $\overline{S}_{\mathsf{q}}(\overline{\Sigma}) \cong_{\mathscr{S}}^{\varepsilon} S_{\tau}(\overline{\Sigma}).$

The proof is similar to the one of Theorem 6.

Corollary 4 Consider the results in Corollary 2. If we select

$$X_{\tau 0} = \left\{ x \in \mathbb{R}^n \mid \left\| x - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}0}) \right\| \le \varepsilon, \ \exists x_{\mathsf{q}0} \in X_{\mathsf{q}0} \right\},\$$

then we have $\overline{S}_{q}(\overline{\Sigma}) \cong^{\varepsilon}_{\mathscr{G}} S_{\tau}(\overline{\Sigma})$.

The proof is similar to the one of Theorem 7. The next two theorems provide similar results as the ones of Theorems 6 and 7, but by using the symbolic model $S_q(\Sigma)$.

Theorem 8 Consider the results in Theorem 4. Let \mathscr{A} denote the set of all \mathbb{R}^n -valued random variables, measurable over \mathscr{F}_0 . If we select

$$X_{\tau 0} = \left\{ a \in \mathscr{A} \mid (\mathbb{E} \left[\|a - H_{\mathsf{q}}(x_{\mathsf{q}0})\|^{q} \right])^{\frac{1}{q}} \le \left(\overline{\alpha}^{-1} \left(\underline{\alpha} \left(\varepsilon^{q} \right) \right) \right)^{\frac{1}{q}}, \exists x_{\mathsf{q}0} \in X_{\mathsf{q}0} \right\},$$

then we have $S_{\mathsf{q}}(\Sigma) \cong_{\mathscr{S}}^{\varepsilon} S_{\tau}(\Sigma).$

The proof is similar to the one of Theorem 6.

Theorem 9 Consider the results in Theorem 5. Let \mathscr{A} denote the set of all \mathbb{R}^n -valued random variables, measurable over \mathscr{F}_0 . If we select

$$X_{\tau 0} = \left\{ a \in \mathscr{A} \mid \left(\mathbb{E} \left[\| a - H_{\mathsf{q}}(x_{\mathsf{q}0}) \|^{q} \right] \right)^{\frac{1}{q}} \leq \varepsilon, \ \exists x_{\mathsf{q}0} \in X_{\mathsf{q}0} \right\},$$

then we have $S_q(\Sigma) \cong_{\mathscr{S}}^{\varepsilon} S_{\tau}(\Sigma)$.

The proof is similar to the one of Theorem 7.

5.3 Control synthesis over $S_q(\Sigma)$

Note that both $\overline{S}_q(\Sigma)$ and $S_q(\Sigma)$ are finite systems. The only difference is that the outputs of the former system are always deterministic points, whereas those of the latter can be non-degenerate random variables. Let us describe the control synthesis for these systems over quantitative specifications, and for example over the safety formula $\Box \varphi_A$, for $A \subset \mathbb{R}^n \subset Y$ (as already been used in Subsection 5.1). Clearly, since the original system $S_\tau(\Sigma)$ is stochastic in the sense that its outputs are non-degenerate random variables similarly to $S_q(\Sigma)$, it would be too conservative to require that it satisfies the formula exactly. Thus, we are rather interested in an input policy that makes $S_\tau(\Sigma)$ satisfy $\Box \varphi_{A_\varepsilon}$ with some $\varepsilon > 0$: recall from Subsection 5.1 that the latter LTL formula can be satisfied by non-degenerate random variables, in contrast to $\Box \varphi_A$. Let us recap how to use abstractions for this task, and let us start with $\overline{S}_q(\Sigma)$ belonging to a more familiar type of systems in which the outputs of states are deterministic.

We label a state x_q of $\overline{S}_q(\Sigma)$ with *A* if $\overline{H}_q(x_q) \in A$ and, say, with *B* otherwise. As a result, we obtain a transition system with labels over the states and can synthesize a control strategy by solving a safety game [24] that makes an output run of $\overline{S}_q(\Sigma)$ satisfy $\Box \varphi_A$. After that, we can exploit ε -approximate bisimilarity to guarantee that the refined input policy makes the corresponding output run of the original system $S_\tau(\Sigma)$ satisfy $\Box \varphi_{A_{\varepsilon}}$.

The main subtlety in the case of $S_q(\Sigma)$ is how to label its states. We cannot do this as for $\overline{S}_q(\Sigma)$, since $H_q(x_q)$ may never be an element of A for any $x_q \in X_q$: indeed, the latter is a set of deterministic points, whereas all the outputs of $S_q(\Sigma)$ can happen to be non-degenerate random variables. In order to cope with this issue, we propose to relax the original problem and at the same time to strengthen the quality of the abstraction. Namely, we can consider a relaxed problem $\Box \varphi_{A_\delta}$ over the abstraction $S_q(\Sigma)$, for some $\delta \in]0 \varepsilon[$, where the latter is now required to be $(\varepsilon - \delta)$ -approximate (rather than ε -approximate) bisimilar to the original system. Clearly $(A_\delta)_{\varepsilon-\delta} \subseteq A_{\varepsilon}$, so that whenever the control policy for $\Box \varphi_{A_\delta}$ is synthesized over $S_q(\Sigma)$, its refined version is guaranteed to enforce $\Box \varphi_{A_\varepsilon}$ over the original system. Thanks to the fact that A_δ contains non-degenerate random variables, we eliminate the conservativeness presented before in the sense that it is likely that there are now points $x_q \in X_q$ in $S_q(\Sigma)$ such that $H_q(x_q) \in A_\delta$. The only remaining question is how to check whether $H_q(x_q) \in A_\delta$. To answer this question, we check that the distance

$$\mathbf{d}\left(H_{\mathbf{q}}(x_{\mathbf{q}}),A\right) = \inf_{a \in A} \left(\mathbb{E} \|\xi_{x_{s}x_{\mathbf{q}}}(N\tau) - a\|^{q}\right)^{1/q}$$
(21)

is smaller than δ , which involves both computing the expectation over the solution of the SDE, and optimizing the value of this expectation. Clearly, such a computation in general cannot be done analytically, and the evaluation of the expectation itself is a highly non-trivial task unless the SDE has a very special form.

We propose a Monte Carlo approach to compute an approximation of the quantity in (21) by means of empirical expectations. Using such an approach, we can estimate $\mathbf{d}(H_q(x_q), A)$ only up to some precision, say θ . If the estimated distance is less than $\delta - \theta$, we are safe to label x_q with A, whereas all other states are labeled by B. Furthermore, since this result is based on a Monte Carlo method, it holds true only with a certain confidence level $1 - \pi$ where $\pi \in [0 \ 1]$. The benefit of our approach is that it is not only valid asymptotically (as the number of samples grows to infinity), but we are also able to provide a number of simulations that is sufficient to estimate $\mathbf{d}(H_q(x_q), A)$ with any given precision θ and with any given confidence $1 - \pi$. This can be considered as an extension of the well-known Hoeffding's inequality [8] to the case when one has to deal with an optimization problem. Note that like in other cases when using Monte-Carlo for distance estimation, there are two levels of error. First, in A_{ε} the ε still refers to the precision of symbolic model in the qth moment metric, but now also a level of (lack of) confidence $\pi \in [0 \ 1]$ is added on top of it.

Regardless of the specification of interest, the main task over $S_q(\Sigma)$ is always to compute some distance as in (21) for any set that appears in the specification, so the method below applies not only to the safety formula $\Box \varphi_A$, but also to more general formulae, which are left as object of the future research.

Suppose that *A* as in (21) is a compact subset of \mathbb{R}^n , and let A^r be the smallest subset of $[\mathbb{R}^n]_r$ such that $A \subseteq \bigcup_{p \in A^r} \mathscr{B}_{\frac{r}{2}}(p)$. Let *M* be the number of samples and let

$$\mathbf{d}_{M}^{r} := \min_{a \in A^{r}} \left(\frac{1}{M} \sum_{i=1}^{M} \left\| \xi_{\mathbf{x}_{s} \mathbf{x}_{q}}^{i}(N\tau) - a \right\|^{q} \right)^{\frac{1}{q}}$$

where the superscript i denotes the index of samples. Now, we have the following theorem.

Theorem 10 For any stochastic control system Σ one has $|\mathbf{d}(H_q(x_q), A) - \mathbf{d}_M^r| \le \theta$ with confidence of at least $1 - \pi$, given that $r < 2\theta$ and that

$$M \ge \frac{|A^r|b(a^*, 2q)}{\pi(\theta - r/2)^{2q}}$$

where $b(a,p) := (1 + |x_s - a|^p) e^{p(p+1) \max\{L_x, Z\}N\tau}$ and $a^* \in \arg \max_{a' \in A^r} ||x_s - a'||$.

The proof can be found in the Appendix. Let us make some comments on Theorem 10. First of all, no matter how many distances one has to evaluate, one can always use the same samples ξ^i and there is no need to generate new samples. Second, to the best of our knowledge, logarithmic bounds on M (as per [11]) are not available in this general case due to the fact that we deal with an unbounded state space. 5.4 Relationship with existing results in the literature

Note that given any precision $\varepsilon \in \mathbb{R}^+$ and sampling time τ , one can always use the results in Theorem 6 to construct a symbolic model $\overline{S}_q(\Sigma)$ that is ε -approximate bisimilar to $S_{\tau}(\Sigma)$ without any state set discretization, but only input set discretization. The results in Theorem 5.1 in [31] also provide symbolic models that are ε -approximate bisimilar to $S_{\tau}(\Sigma)$. However, the results in [31] require both state and input set discretization and cannot be applied for any sampling time τ if the precision ε is lower than the thresholds introduced in inequality (5.5) in [31]. Furthermore, while the results in [31] only provide symbolic models with states having deterministic output values, the ones in this work provide symbolic models with states having probabilistic output values as well, which can result in less conservative symbolic models (cf. Remark 3).

One can compare the results provided in Theorems 2 (corr. 6) and 3 (corr. 7) with the results provided in Theorems 5.1 and 5.3 in [31] in terms of the size of the generated symbolic models. It can be readily verified that the precisions of the symbolic models proposed here and the ones proposed in [31] are approximately the same as long as both use the same input set quantization parameter μ and the state space quantization parameter, called here ν , in [31] is equal to the parameter η in (9), i.e. $\nu \leq (\underline{\alpha}^{-1}(e^{-\kappa N\tau}\eta_0))^{1/q}$, where $\eta_0 = \max_{u_q \in U_q} V(\overline{\xi}_{x,u_q}(\tau), x_s)$. The reason their precisions are approximately (rather than exactly) the same is because we use $h_{x_s}(\sigma, (N+1)\tau)$ in conditions (15) and (17) in this paper rather than $h(\tau) = \sup_{x \in D} h_x(\tau)$ that is being used in conditions 5.4 and 5.14 in [31] for a compact set $D \subset \mathbb{R}^n$. By assuming that $h_{x_s}(\sigma, (N+1)\tau)^{\frac{1}{q}}$ and $h(\tau)^{\frac{1}{q}}$ are much smaller than η and ν , respectively, or $h_{x_s}(\sigma, (N+1)\tau) \approx h(\tau)$, one should expect to get the same precisions for the symbolic models provided here and those provided in [31].

The number of states of the proposed symbolic model in this paper is $|[U]_{\mu}|^N$. Assume that we are interested in the dynamics of Σ on a compact set $D \subset \mathbb{R}^n$. Since the set of states of the proposed symbolic model in [31] is $[D]_{\nu}$, its size is $|[D]_{\nu}| = \frac{K}{\nu^n}$, where *K* is a positive constant proportional to the volume of *D*. Hence, it is more convenient to use the proposed symbolic model here rather than the one proposed in [31] as long as:

$$|[\mathsf{U}]_{\mu}|^{N} \leq \frac{K}{\left(\underline{lpha}^{-1}\left(\mathsf{e}^{-\kappa N au}\eta_{0}
ight)
ight)^{n/q}}$$

Without loss of generality, one can assume that $\underline{\alpha}(r) = r$ for any $r \in \mathbb{R}_0^+$. Hence, for sufficiently large value of *N*, the size of the proposed symbolic model here is smaller than the one proposed in [31] as long as:

$$|[\mathsf{U}]_{\mu}|e^{\frac{-\kappa m}{q}} \le 1. \tag{22}$$

One can readily verify that if the state-space dimension of Σ is very large, i.e. n >> 1, or Σ is "strongly" δ -ISS-M_q, i.e. $\kappa >> 1$, inequality (22) is most likely satisfied.

Finally, remark that the framework proposed in this paper lets us to construct less conservative symbolic models with probabilistic output values while the proposed one in [31] only provides symbolic models with deterministic output values.

6 Example

We show the effectiveness of the results presented in this work by constructing a bisimilar symbolic model for the model of a road network. The road is divided in 5 cells of 250 meters with 2 entries and 2 ways out, as depicted schematically in Figure 2. The model is borrowed from [5], explained in details in [27], however it is now affected by noise and described in continuous time.



Fig. 2 Model of a road divided in 5 cells with 2 entries and 2 ways out.

The two entries are controlled by traffic lights, denoted by u_1 and u_2 , that enable (green light) or not (red light) the vehicles to pass. In this model the length of a cell is in kilometres (0.25 km), and the flow speed of the vehicles is 70 kilometers per hour (km/h). Moreover, during the sampling time interval τ , it is assumed that 6 vehicles pass the entry controlled by the light u_1 , 8 vehicles pass the entry controlled by the light u_2 , and one quarter of vehicles that leave cell 1 goes out on the first exit. We assume that both lights cannot be red at the same time. The model of Σ is described by:

$$d\xi = (A\xi + Bv) dt + \xi dW_t, \qquad (23)$$

where

$$A = 10^{4} \times \begin{bmatrix} -0.0541 & 0 & 0 & 0 & 0 \\ 0.3224 & -0.1370 & 0 & 0 & 0 \\ -0.7636 & 0.3224 & -0.0541 & 0 & 0 \\ 2.1122 & -0.7636 & 0.1260 & -0.0541 & 0 \\ -6.2132 & 2.1122 & -0.2205 & 0.1260 & -0.0541 \end{bmatrix},$$

$$B = 10^{4} \times \begin{bmatrix} 0.0696 & 0 \\ -0.2743 & 0 \\ 0.7075 & 0.0696 \\ -2.0081 & -0.0924 \\ 5.9802 & 0.1911 \end{bmatrix},$$

 $U = \{u_0, u_1, u_2\} = \{[6 \ 8]^T, [6 \ 0]^T, [0 \ 8]^T\}, \text{ and } \xi_i \text{ is the number of vehicles in cell } i \text{ of the road. We point out that } \mathscr{U}_{\tau} \text{ contains curves taking values in U. Since U is finite, as explained in Remark 4, } \mu = 0 \text{ is to be used in (15), (17), (19), and (20). Using LMI}$

(4), one can verify that the function $V(x,x') = (x - x')^T P(x - x')$, for any $x, x' \in \mathbb{R}^5$, where

$$P = 10^{4} \times \begin{bmatrix} 76763.4393 & -2101.1583 & 3790.9182 & -155.6576 & -125.9871 \\ -2101.1583 & 10676.9437 & 1237.3552 & -86.6855 & 100.5718 \\ 3790.9182 & 1237.3552 & 1823.02431 & 171.1549 & -71.1162 \\ -155.6576 & -86.6855 & 171.1549 & 229.2134 & -5.5649 \\ -125.9871 & 100.5718 & -71.1162 & -5.5649 & 33.3977 \end{bmatrix}$$

satisfies conditions (i)-(iii) in Definition 3 with q = 2, $\kappa = 300$, $\underline{\alpha}(r) = \frac{1}{2}\lambda_{\min}(P)r$, $\overline{\alpha}(r) = \frac{1}{2}\lambda_{\max}(P)r$, $\rho(r) = \frac{5||\underline{\beta}||^2 ||\underline{P}||}{2\kappa}r^2$, $\forall r \in \mathbb{R}^+_0$. Hence, Σ is δ -ISS-M₂, equipped with the δ -ISS-M₂ Lyapunov function *V*. Using the results of Theorem 1, provided in [31], one gets that functions $\beta(r,s) = \underline{\alpha}^{-1}(\overline{\alpha}(r)e^{-\kappa s})$ and $\gamma(r) = \underline{\alpha}^{-1}(\frac{1}{e\kappa}\rho(r))$ satisfy property (2) for Σ . We choose the source state as the one picked in [5], i.e. $x_s = [3.8570 \ 3.3750 \ 3.3750 \ 8.5177 \ 8.5177]^T$.

For a given precision $\varepsilon = 0.5$ and fixed sampling time $\tau = 0.00277$ h (10 sec), the parameter *N* for $\overline{S}_q(\Sigma)$, based on inequality (15) in Theorem 2, is obtained as 14. Therefore, the resulting cardinality of the set of states for $\overline{S}_q(\Sigma)$ is $|U|^{14} = 3^{14} =$ 4782969. Using the aforementioned parameters, one gets $\eta \leq 6.0776 \times 10^{-6}$, where η is given in (9). Remark that the results in Theorems 3 and 5 cannot be applied here because $(\beta(\varepsilon^q, \tau))^{\frac{1}{q}} > \varepsilon$. Using criterion (22), one has $|U|e^{\frac{-\kappa\tau n}{q}} = 0.37$, implying that the proposed approach in this paper is more efficient in terms of the size of the abstraction than the one proposed in [31]. We elaborate more on this at the end of the section.

Remark 5 By considering the deterministic control system $\overline{\Sigma}$ and using the results in Corollary 1 and the same parameters q as the ones in $\overline{S}_q(\Sigma)$, one obtains $\varepsilon = 0.01$ in (16). Therefore, as expected, $\overline{S}_q(\overline{\Sigma})$ (i.e. symbolic model for the deterministic control system $\overline{\Sigma}$) provides much smaller precision than $\overline{S}_q(\Sigma)$ (i.e. symbolic model for the stochastic control system Σ) while having the same size as $\overline{S}_q(\Sigma)$.

Now the objective, as inspired by the one suggested in [5], is to design a schedule for the coordination of traffic lights enforcing Σ to satisfy a *safety* and a *fairness* property. The safety part is to keep the density of traffic lower than 16 vehicles per cell which can be encoded via the LTL specification $\Box \varphi_W$, where $W = [0 \ 16]^5$ and φ_W is a label characterizing the set W. The fairness part requires to alternate the accesses between the two traffic lights and to allow only 3 identical consecutive modes of red light ensuring fairness between two traffic lights. We implemented the proposed techniques in this paper on top of the recently developed synthesis toolbox SCOTS [22] on an iMac with CPU 3.5GHz Intel Core i7. The CPU time used for computing the abstraction and synthesizing the controller have amounted to 68 and 5 seconds, respectively. Starting from the initial condition $x_0 = [1.417 \ 4.993 \ 10.962 \ 9.791 \ 14.734]^T$, we obtain a periodic schedule $v = (u_0u_0u_0u_2u_1u_0u_0u_2u_1u_2)^{\omega}$ keeping u_0 as much as possible in each period in order to maximize number of vehicles accessing the road. One can readily verify from the computed schedule v that the fairness property is satisfied. Figure 3 displays a few realizations of the closed-loop solution process ξ_{x_0v} . In Figure 3 bottom right, we show the average value (over 100000 experiments) of the distance (in the 2nd moment metric) in time of the solution process ξ_{x_0v} to the set *W*, namely $\|\xi_{x_0v}(t)\|_W$, where the point-to-set distance is defined as $\|x\|_W = \inf_{w \in W} \|x - w\|$. Observe that the empirical average distance is as expected lower than the precision $\varepsilon = 0.5$ due to the conservative nature of δ -ISS-M₂ function *V*.



Fig. 3 A few realizations of the closed-loop solution process ξ_{x_0v} (top panel and the first two figures from the left in the bottom panel) and the average values (over 100000 experiments) of the distance of the solution process ξ_{x_0v} to the set *W* (bottom right panel).

To compute exactly the size of the symbolic model, proposed in Theorem 5.1 in [31], we consider the dynamics of Σ over the subset $D = [0 \ 16]^5$ of \mathbb{R}^5 . Remark that Theorem 5.3 in [31] cannot be applied here because $(\beta(\varepsilon^q, \tau))^{\frac{1}{q}} > \varepsilon$. Using the same precision $\varepsilon = 0.5$ and sampling time $\tau = 0.00277$ as the ones here, and the inequalities (5.3) and (5.4) in [31], we obtain the state space quantization parameter as $v \le 6.0776 \times 10^{-6}$. Therefore, if one uses $v = 6.0776 \times 10^{-6}$, the cardinality of the state set of the symbolic model, provided by the results in Theorem 5.1 in [31], is equal to $(\frac{16}{v})^5 = 1.2645 \times 10^{32}$ which is practically not tractable and much higher than the one proposed here, amounting instead to 4782969 states.

7 Conclusions

In this paper we have proposed a symbolic abstraction technique for incrementally stable stochastic control systems (and corresponding deterministic model), which features only the discretization of the input set. The proposed approach is potentially more scalable than the one proposed in [31] for higher dimensional stochastic control systems.

Future work will concentrate on *efficient* implementations of the symbolic models proposed in this work on top of the recently developed synthesis toolbox SCOTS [22], as well as on more efficient controller synthesis techniques by using binary search trees data structure. The current implementation is not computationally very efficient for the closed-loop implementation of the synthesized symbolic controllers for deterministic control systems because one requires to loop over the symbolic states to find the ones with the outputs ε -close to the current measured state of the concrete system in each feedback iteration.

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Appendix

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Proof of Lemma 2: Let $x_q \in X_q$, where $x_q = (u_1, u_2, ..., u_N)$, and $u_q \in U_q$. Using the definition of $\overline{S}_q(\Sigma)$, one obtains $x'_q = (u_2, ..., u_N, u_q) \in \mathbf{Post}_{u_q}(x_q)$. Since *V* is a δ -ISS-M_{*q*} Lyapunov function for Σ , we have:

$$\underline{\alpha}\left(\left\|\overline{\xi}_{\overline{H}_{q}(x_{q})u_{q}}(\tau)-\overline{H}_{q}\left(x_{q}'\right)\right\|^{q}\right)\leq V(\overline{\xi}_{\overline{H}_{q}(x_{q})u_{q}}(\tau),\overline{H}_{q}\left(x_{q}'\right))$$

$$=V(\overline{\xi}_{\overline{\xi}_{x_{s}x_{q}}(N\tau)u_{q}}(\tau),\overline{\xi}_{x_{s}x_{q}'}(N\tau))=V(\overline{\xi}_{\overline{\xi}_{x_{s}u_{1}}(\tau)(u_{2},...,u_{N},u_{q})}(N\tau),\overline{\xi}_{x_{s}(u_{2},...,u_{N},u_{q})}(N\tau))$$

$$\leq e^{-\kappa N\tau}V(\overline{\xi}_{x_{s}u_{1}}(\tau),x_{s}).$$
(24)

We refer the interested readers to the proof of Theorem 1 in [31] to see how we derived the inequality (24). Hence, one gets

$$\|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}\left(x_{\mathsf{q}}'\right)\| \le (\underline{\alpha}^{-1}(\mathrm{e}^{-\kappa N\tau}V(\overline{\xi}_{x_{s}u_{1}}(\tau), x_{s})))^{1/q},$$
(25)

because of $\underline{\alpha} \in \mathscr{K}_{\infty}$. Since the inequality (25) holds for all $x_q \in X_q$ and $u_q \in U_q$, and $\underline{\alpha} \in \mathscr{K}_{\infty}$, inequality (8) holds.

Proof of Lemma 3: Let $x_q \in X_q$, where $x_q = (u_1, u_2, ..., u_N)$, and $u_q \in U_q$. Using the definition of $\overline{S}_q(\Sigma)$, one obtains $x'_q = (u_2, ..., u_N, u_q) \in \mathbf{Post}_{u_q}(x_q)$. Since Σ is δ -ISS-M_q and using inequality (2), we have:

$$\begin{split} \|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}\left(x_{\mathsf{q}}'\right)\|^{q} &= \|\overline{\xi}_{\overline{\xi}_{x_{s}x_{\mathsf{q}}}(N\tau)u_{\mathsf{q}}}(\tau) - \overline{\xi}_{x_{s}x_{\mathsf{q}}'}(N\tau)\|^{q} \\ &= \|\overline{\xi}_{\overline{\xi}_{x_{s}u_{1}}(\tau)(u_{2},\ldots,u_{N},u_{\mathsf{q}})}(N\tau) - \overline{\xi}_{x_{s}(u_{2},\ldots,u_{N},u_{\mathsf{q}})}(N\tau)\|^{q} \leq \beta(\|\overline{\xi}_{x_{s}u_{1}}(\tau) - x_{s}\|^{q}, N\tau). \end{split}$$

Hence, one gets

$$\|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\| \le (\beta(\|\overline{\xi}_{x_{s}u_{1}}(\tau) - x_{s}\|^{q}, N\tau))^{1/q}.$$
(26)

Since the inequality (26) holds for all $x_q \in X_q$ and all $u_q \in U_q$, and β is a \mathscr{K}_{∞} function with respect to its first argument when the second one is fixed, inequality (10) holds.

Proof of Theorem 2: We start by proving that *R* is an ε -approximate simulation relation from $S_{\tau}(\Sigma)$ to $\overline{S}_{q}(\Sigma)$. Consider any $(x_{\tau}, x_{q}) \in R$. Condition (i) in Definition 5 is satisfied because

$$\left(\mathbb{E}[\|x_{\tau} - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}})\|^{q}]\right)^{\frac{1}{q}} \le \left(\underline{\alpha}^{-1}(\mathbb{E}[V(x_{\tau}, \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}))])\right)^{\frac{1}{q}} \le \varepsilon.$$
(27)

We used the convexity assumption of $\underline{\alpha}$ and the Jensen inequality [16] to show the inequalities in (27). Let us now show that condition (ii) in Definition 5 holds. Consider any $v_{\tau} \in U_{\tau}$. Choose an input $u_{q} \in U_{q}$ satisfying

$$\|v_{\tau} - u_{\mathsf{q}}\|_{\infty} = \|v_{\tau}(0) - u_{\mathsf{q}}(0)\| \le \mu.$$
(28)

Note that the existence of such u_q is guaranteed by U being a finite union of boxes and by the inequality $\mu \leq span(U)$ which guarantees that $U \subseteq \bigcup_{p \in [U]_{\mu}} \mathscr{B}_{\mu}(p)$. Consider the transition $x_{\tau} \xrightarrow{v_{\tau}} x'_{\tau} = \xi_{x_{\tau}v_{\tau}}(\tau)$ P-a.s. in $S_{\tau}(\Sigma)$. Since V is a δ -ISS-M_q Lyapunov function for Σ and using inequality (28), we have (cf. equation (3.3) in [31])

$$\mathbb{E}[V(x_{\tau}',\xi_{\overline{H}_{q}(x_{q})u_{q}}(\tau))] \leq \mathbb{E}[V(x_{\tau},\overline{H}_{q}(x_{q}))]e^{-\kappa\tau} + \frac{1}{e\kappa}\rho(\|\upsilon_{\tau}-u_{q}\|_{\infty}) \leq \underline{\alpha}\left(\varepsilon^{q}\right)e^{-\kappa\tau} + \frac{1}{e\kappa}\rho(\mu).$$
(29)

Observe that existence of u_q , by the definition of $\overline{S}_q(\Sigma)$, implies the existence of $x_q \xrightarrow{u_q} x'_q$ in $\overline{S}_q(\Sigma)$. Using Lemma 1, the concavity of $\hat{\gamma}$, the Jensen inequality [16],

equation (9), the inequalities (7), (15), (29), and triangle inequality, we obtain

$$\begin{split} & \mathbb{E}[V(x'_{\tau},\overline{H}_{\mathsf{q}}(x'_{\mathsf{q}}))] = \mathbb{E}[V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)) + V(x'_{\tau},\overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})) - V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau))] \\ & = \mathbb{E}[V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau))] + \mathbb{E}[V(x'_{\tau},\overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})) - V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau))] \\ & \leq \underline{\alpha}(\varepsilon^{q})\mathsf{e}^{-\kappa\tau} + \frac{1}{\mathsf{e}\kappa}\rho(\mu) + \mathbb{E}[\widehat{\gamma}(\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})\|)] \\ & \leq \underline{\alpha}(\varepsilon^{q})\mathsf{e}^{-\kappa\tau} + \frac{1}{\mathsf{e}\kappa}\rho(\mu) \\ & \quad + \widehat{\gamma}(\mathbb{E}[\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) + \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})\|]) \\ & \leq \underline{\alpha}(\varepsilon^{q})\mathsf{e}^{-\kappa\tau} + \frac{1}{\mathsf{e}\kappa}\rho(\mu) \\ & \quad + \widehat{\gamma}(\mathbb{E}[\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|] + \|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})\|]) \\ & \leq \underline{\alpha}(\varepsilon^{q})\mathsf{e}^{-\kappa\tau} + \frac{1}{\mathsf{e}\kappa}\rho(\mu) + \widehat{\gamma}((h_{x_{s}}((N+1)\tau)))^{\frac{1}{q}} + \eta) \leq \underline{\alpha}(\varepsilon^{q}). \end{split}$$

Therefore, we conclude that $(x'_{\tau}, x'_{q}) \in R$ and that condition (ii) in Definition 5 holds.

Now we prove that R^{-1} is an ε -approximate simulation relation from $\overline{S}_q(\Sigma)$ to $S_\tau(\Sigma)$. Consider any $(x_\tau, x_q) \in R$ (or equivalently $(x_q, x_\tau) \in R^{-1}$). As showed in the first part of the proof, condition (i) in Definition 5 is satisfied. Let us now show that condition (ii) in Definition 5 holds. Consider any $u_q \in U_q$. Choose the input $v_\tau = u_q$ and consider $x'_\tau = \xi_{x_\tau v_\tau}(\tau) \mathbb{P}$ -a.s. in $S_\tau(\Sigma)$. Since V is a δ -ISS-M_q Lyapunov function for Σ , one obtains (cf. equation 3.3 in [31]):

$$\mathbb{E}[V(x_{\tau}',\xi_{\overline{H}_{q}(x_{q})u_{q}}(\tau))] \leq e^{-\kappa\tau} \mathbb{E}[V(x_{\tau},\overline{H}_{q}(x_{q}))] \leq e^{-\kappa\tau} \underline{\alpha}\left(\varepsilon^{q}\right).$$
(30)

Using Lemma 1, the definition of $\overline{S}_q(\Sigma)$, the concavity of $\hat{\gamma}$, the Jensen inequality [16], equation (9), the inequalities (7), (15), (30), and triangle inequality, we obtain

$$\begin{split} & \mathbb{E}[V(x'_{\tau},\overline{H}_{\mathsf{q}}(x'_{\mathsf{q}}))] = \mathbb{E}[V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)) + V(x'_{\tau},\overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})) - V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau))] \\ & = \mathbb{E}[V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau))] + \mathbb{E}[V(x'_{\tau},\overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})) - V(x'_{\tau},\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau))] \\ & \leq e^{-\kappa\tau}\underline{\alpha}(\varepsilon^{q}) + \mathbb{E}[\widehat{\gamma}(\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})\|)] \\ & \leq e^{-\kappa\tau}\underline{\alpha}(\varepsilon^{q}) + \widehat{\gamma}(\mathbb{E}[\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) + \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})\|]) \\ & \leq e^{-\kappa\tau}\underline{\alpha}(\varepsilon^{q}) + \widehat{\gamma}(\mathbb{E}[\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|] + \|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x'_{\mathsf{q}})\|)) \\ & \leq e^{-\kappa\tau}\underline{\alpha}(\varepsilon^{q}) + \widehat{\gamma}((h_{x_{s}}((N+1)\tau))^{\frac{1}{q}} + \eta)) \leq \underline{\alpha}(\varepsilon^{q}). \end{split}$$

Therefore, we conclude that $(x'_{\tau}, x'_{q}) \in R$ (or equivalently $(x'_{q}, x'_{\tau}) \in R^{-1}$) and condition (ii) in Definition 5 holds.

Proof of Theorem 3: We start by proving that *R* is an ε -approximate simulation relation from $S_{\tau}(\Sigma)$ to $\overline{S}_{q}(\Sigma)$. Consider any $(x_{\tau}, x_{q}) \in R$. Condition (i) in Definition 5 is satisfied by the definition of *R*. Let us now show that condition (ii) in Definition 5 holds. Consider any $v_{\tau} \in U_{\tau}$. Choose an input $u_{q} \in U_{q}$ satisfying

$$\|v_{\tau} - u_{\mathsf{q}}\|_{\infty} = \|v_{\tau}(0) - u_{\mathsf{q}}(0)\| \le \mu.$$
 (31)

Note that the existence of such u_q is guaranteed by U being a finite union of boxes and by the inequality $\mu \leq span(U)$ which guarantees that $U \subseteq \bigcup_{p \in [U]_{\mu}} \mathscr{B}_{\mu}(p)$. Consider the transition $x_{\tau} \xrightarrow{v_{\tau}} x'_{\tau} = \xi_{x_{\tau}v_{\tau}}(\tau)$ P-a.s. in $S_{\tau}(\Sigma)$. It follows from the δ -ISS-M_q assumption on Σ and (31) that:

$$\mathbb{E}[\|x_{\tau}' - \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|^{q}] \leq \beta(\mathbb{E}[\|x_{\tau} - \overline{H}_{\mathsf{q}}(x_{q})\|^{q}], \tau) + \gamma(\|\upsilon_{\tau} - u_{\mathsf{q}}\|_{\infty}) \leq \beta(\varepsilon^{q}, \tau) + \gamma(\mu).$$
(32)

Existence of u_q , by the definition of $\overline{S}_q(\Sigma)$, implies the existence of $x_q \xrightarrow{u_q} x'_q$ in $\overline{S}_q(\Sigma)$. Using equation (9), the inequalities (5), (17), (32), and triangle inequality, we obtain

$$\begin{split} (\mathbb{E}[\|x_{\tau}' - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\|^{q}])^{\frac{1}{q}} &= (\mathbb{E}[\|x_{\tau}' - \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) + \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) \\ &- \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) + \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\|^{q}])^{\frac{1}{q}} \\ &\leq (\mathbb{E}[\|x_{\tau}' - \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|^{q}])^{\frac{1}{q}} + (\mathbb{E}[\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|^{q}])^{\frac{1}{q}} \\ &+ (\mathbb{E}[\|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\|^{q}])^{\frac{1}{q}} \\ &\leq (\beta(\varepsilon^{q}, \tau) + \gamma(\mu))^{\frac{1}{q}} + (h_{x_{s}}((N+1)\tau))^{\frac{1}{q}} + \eta \leq \varepsilon. \end{split}$$

Therefore, we conclude that $(x'_{\tau}, x'_{q}) \in R$ and that condition (ii) in Definition 5 holds.

Now we prove that R^{-1} is an ε -approximate simulation relation from $\overline{S}_q(\Sigma)$ to $S_{\tau}(\Sigma)$. Consider any $(x_{\tau}, x_q) \in R$ (or equivalently $(x_q, x_{\tau}) \in R^{-1}$). Condition (i) in Definition 5 is satisfied by the definition of R. Let us now show that condition (ii) in Definition 5 holds. Consider any $u_q \in U_q$. Choose the input $v_{\tau} = u_q$ and consider $x'_{\tau} = \xi_{x_{\tau}v_{\tau}}(\tau)$ P-a.s. in $S_{\tau}(\Sigma)$. Since Σ is δ -ISS-M_q, one obtains:

$$\mathbb{E}[\|x_{\tau}' - \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|^{q}] \le \beta(\mathbb{E}[\|x_{\tau} - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}})\|^{q}], \tau) \le \beta(\varepsilon^{q}, \tau).$$
(33)

Using definition of $\overline{S}_q(\Sigma)$, equation (9), the inequalities (5), (17), (33), and the triangle inequality, we obtain

$$\begin{split} (\mathbb{E}[\|x_{\tau}' - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\|^{q}])^{\frac{1}{q}} &= (\mathbb{E}[\|x_{\tau}' - \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) + \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) \\ &- \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) + \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\|^{q}])^{\frac{1}{q}} \\ &\leq (\mathbb{E}[\|x_{\tau}' - \xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|^{q}])^{\frac{1}{q}} + (\mathbb{E}[\|\xi_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau)\|^{q}])^{\frac{1}{q}} \\ &+ (\mathbb{E}[\|\overline{\xi}_{\overline{H}_{\mathsf{q}}(x_{\mathsf{q}})u_{\mathsf{q}}}(\tau) - \overline{H}_{\mathsf{q}}(x_{\mathsf{q}}')\|^{q}])^{\frac{1}{q}} \\ &\leq (\beta(\varepsilon^{q}, \tau))^{\frac{1}{q}} + (h_{x_{s}}((N+1)\tau))^{\frac{1}{q}} + \eta \leq \varepsilon. \end{split}$$

Therefore, we conclude that $(x'_{\tau}, x'_{q}) \in R$ (or equivalently $(x'_{q}, x'_{\tau}) \in R^{-1}$) and condition (ii) in Definition 5 holds.

Proof of Theorem 10: Denote $\hat{\theta} := \theta - r/2 > 0$, and $\mathbf{d}_M(a) := \left(\frac{1}{M} \sum_{i=1}^M \|\xi_{x_s x_q}^i - a\|^q\right)^{\frac{1}{q}}$ for all $a \in \mathbb{R}^n$. It follows from [11, Theorem 4.5.4] that for all $p \ge 1$ and $a \in \mathbb{R}^n$

$$\mathbb{E}\left[\|\xi_{x_s x_{\mathsf{g}}}(N\tau) - a\|^p\right] \le b(a, p).$$

Since we do not assume that the set of continuous states is bounded, the distance can be any positive real number, and the usual method of applying Hoeffding's inequality does not work in this case. Instead we use Chernoff-type inequality (e.g. see above formula (1) in [4]), which implies that for any $a' \in A^r$:

$$\mathbb{P}\left(\left|\left(\mathbf{d}(H_{\mathsf{q}}(x_{\mathsf{q}}),a')\right)^{q}-(\mathbf{d}_{M}(a'))^{q}\right|\geq\hat{\theta}\right)\leq\frac{b(a',2q)}{M\hat{\theta}^{2}}.$$

Furthermore, since $x \mapsto x^q$ is Hölder continuous with power q,

$$\mathbb{P}\left(\left|\mathbf{d}(H_{\mathsf{q}}(x_{\mathsf{q}}),a')-\mathbf{d}_{M}(a')\right|\geq\hat{\theta}\right)\leq\frac{b(a',2q)}{M\hat{\theta}^{2q}}.$$

Thus, for the union of such events over $a' \in A^r$, we have

$$\mathbb{P}\left(\exists a' \in A^r \text{ s.t. } \left| \mathbf{d}(H_{\mathsf{q}}(x_{\mathsf{q}}), a') - \mathbf{d}_M(a') \right| \ge \hat{\theta} \right) \le \frac{|A^r|b(a^*, 2q)}{M\hat{\theta}^{2q}}, \tag{34}$$

due to the fact that the probability of a union is dominated by the sum of probabilities. Let $[\cdot]: A \to A^r$ be any surjective map such that $||a - [a]|| \le r/2$ for all $a \in A$, i.e. $[\cdot]$ chooses an r/2-close point in the grid A^r . Using this map, we can extrapolate the inequality (34) to the whole set *A* since

$$\begin{aligned} |\mathbf{d}(H_{q}(x_{q}), a) - \mathbf{d}_{M}([a])| &\leq |\mathbf{d}(H_{q}(x_{q}), a) - \mathbf{d}(H_{q}(x_{q}), [a])| + |\mathbf{d}(H_{q}(x_{q}), [a]) - \mathbf{d}_{M}([a])| \\ &\leq r/2 + |\mathbf{d}(H_{q}(x_{q}), [a]) - \mathbf{d}_{M}([a])|, \end{aligned}$$

where we used the fact that $|\mathbf{d}(H_q(x_q), a) - \mathbf{d}(H_q(x_q), [a])| \le ||a - [a]||$ by the triangle inequality. As a result, the following inequality holds:

$$\mathbb{P}\left(\exists a \in A \text{ s.t. } |\mathbf{d}(H_{q}(x_{q}), a) - \mathbf{d}_{M}([a])| \geq \theta\right)$$

$$\leq \mathbb{P}\left(\exists a' \in A' \text{ s.t. } |\mathbf{d}(H_{q}(x_{q}), a') - \mathbf{d}_{M}(a')| \geq \hat{\theta}\right).$$
(35)

On the other hand, since for any two functions $f, g : A \to \mathbb{R}$ it holds that

$$\left|\inf_{a\in A} f(a) - \inf_{a\in A} g(a)\right| \le \sup_{a\in A} |f(a) - g(a)|,$$

we obtain that

$$\mathbb{P}\left(|\mathbf{d}(H_{\mathsf{q}}(x_{\mathsf{q}}), A) - \mathbf{d}_{M}^{r}| \ge \theta\right) \le \mathbb{P}\left(\exists a \in A \text{ s.t. } |\mathbf{d}(H_{\mathsf{q}}(x_{\mathsf{q}}), a) - \mathbf{d}_{M}([a])| \ge \theta\right).$$

Combining the latter inequality with (34) and (35) yields:

$$\mathbb{P}\left(|\mathbf{d}(H_{\mathsf{q}}(x_{\mathsf{q}}), A) - \mathbf{d}_{M}^{r}| \ge \theta\right) \le \frac{|A^{r}|b(a^{*}, 2q)}{M\hat{\theta}^{2q}},$$

and in case *M* satisfies the assumption of the theorem, the right-hand side is bounded above by π as desired.