

Modularising inductive families

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Workshop on Generic Programming
18 September 2011, Tokyo, Japan

Internalism Externalism

Internalism

Constraints internalised in datatypes

```
data Vec (A : Set) : Nat → Set where
```

```
  [] : Vec A 0
```

```
  _::__ : (x : A) →  
          {n : Nat} (xs : Vec A n) →  
          Vec A (suc n)
```

```
x :: y :: z :: [] : Vec A 3
```

Internalist clarity

Constraints cleanly expressed and managed

`zipWith3` :

`(f : A → B → C → D) →`

`Vec A n → Vec B n → Vec C n → Vec D n`

`zipWith3 f [] [] []`

`= []`

`zipWith3 f (x :: xs) (y :: ys) (z :: zs)`

`= f x y z :: zipWith3 f xs ys zs`

Internalist libraries

dreadful reusability/composability

`insert : Nat → List Nat → List Nat`

`vinsert :`

`Nat → Vec Nat n → Vec Nat (suc n)`

`sinsert :`

`(x : Nat) → SList b → SList (b ∩ x)`

`data SList : Nat → Set` where

`nil : ∀ {b} → SList b`

`cons : (x : Nat) → ∀ {b} → b ≤ x →
 (xs : SList x) → SList b`

Externalism

Predicates imposed on existing datatypes

$(xs : \text{List Nat}) \times \text{Length } n \text{ } xs$

data Length : Nat \rightarrow List A \rightarrow Set where

nil : Length 0 []

cons : $\forall \{x \ n \ xs\} \rightarrow$

Length n xs \rightarrow

Length (suc n) (x :: xs)

Externalist composability

Easy to impose multiple constraints

$(xs : List\ Nat) \times \text{Length } n\ xs \times \text{Sorted } b\ xs$

data Sorted : Nat → List Nat → Set where

nil : $\forall \{b\} \rightarrow \text{Sorted } b\ []$

cons : $\forall \{x\ b\} \rightarrow b \leq x \rightarrow$
 $\forall \{xs\} \rightarrow \text{Sorted } x\ xs \rightarrow$
 $\text{Sorted } b\ (x :: xs)$

Externalist composability

Ideal for structuring libraries

`insert : Nat → List Nat → List Nat`

`insert-length :`

`Length n xs →`

`Length (suc n) (insert x xs)`

`insert-sorted :`

`Sorted b xs →`

`Sorted (b ∩ x) (insert x xs)`

Is it possible to import

externalist composability

into

internalist libraries ?

Or: Can we get sorted vectors and
insert on sorted vectors for free?

Constraints

Multiple constraints

Internalism

??

??

Externalism

Predicates

**Pointwise conjunction
of predicates**

Constraints

Multiple constraints

Internalism

Ornaments!

??

Conor McBride

Externalism

Predicates

Pointwise conjunction
of predicates

Ornaments

Information added to a datatype to get a fancier one

```
data List (A : Set) : Set where
```

```
  [] : List A
```

```
  _::__ : (x : A) →  
         (xs : List A) →  
         List A
```

Ornaments

Information added to a datatype to get a fancier one

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data List (A : Set) : Nat → Set where
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```
_::_ : (x : A) →  
      {n : Nat} (xs : List A n) →  
      List A (suc n)
```

Ornaments

Information added to a datatype to get a fancier one

data Vec (A : Set) : Nat → Set where

[] : Vec A 0

:: : (x : A) →
{n : Nat} (xs : Vec A n) →
Vec A (suc n)

Constraints

Multiple constraints

Internalism

Ornaments

??

induce

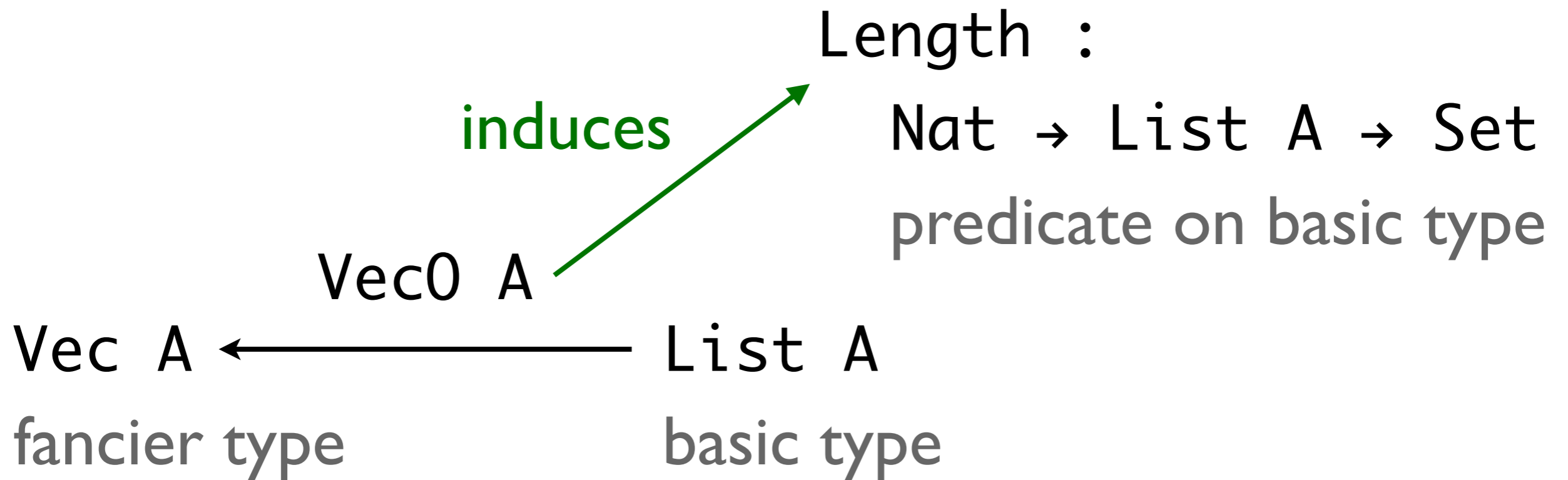


Externalism

Predicates

Pointwise conjunction
of predicates

Ornaments induce predicates



Ornaments induce predicates

and corresponding isomorphisms

internalist

externalist

$\text{Vec } A \ n \cong (\text{xs} : \text{List } A) \times \text{Length } n \ \text{xs}$

fancier type

basic type

induced predicate

$\text{SList } b \cong (\text{xs} : \text{List } \text{Nat}) \times \text{Sorted } b \ \text{xs}$

Function upgrade

with the help of the isomorphisms

$$\text{Vec Nat } n \cong (\text{xs} : \text{List Nat}) \times \text{Length } n \text{ xs}$$

$$\begin{array}{lcl} \text{vinsert} : \text{Nat} & \rightarrow & \\ \text{Vec Nat } n & \rightarrow & \text{Vec Nat (suc } n) \\ \quad \parallel & & \parallel \\ \text{xs} : \text{List Nat} & \mapsto & \text{insert } x \text{ xs} : \text{List Nat} \\ \text{l} : \text{Length } n \text{ xs} & \mapsto & \text{insert-length } \text{l} : \\ & & \text{Length (suc } n) \text{ (insert } x \text{ xs)} \end{array}$$

$\text{insert} : \text{Nat} \rightarrow \text{List Nat} \rightarrow \text{List Nat}$

$\text{insert-length} :$

$\text{Length } n \text{ xs} \rightarrow \text{Length (suc } n) \text{ (insert } x \text{ xs)}$

Function upgrade

with the help of the isomorphisms

$\text{Vec Nat } n \cong (\text{xs} : \text{List Nat}) \times \text{Length } n \text{ xs}$

$\text{vinsert} : \text{Nat} \rightarrow \text{Vec Nat } n \rightarrow \text{Vec Nat } (\text{suc } n)$

$\text{SList } b \cong (\text{xs} : \text{List Nat}) \times \text{Sorted } b \text{ xs}$

$\text{sininsert} : (\text{x} : \text{Nat}) \rightarrow \text{SList } b \rightarrow \text{SList } (b \sqcap \text{x})$

$\text{insert} : \text{Nat} \rightarrow \text{List Nat} \rightarrow \text{List Nat}$

$\text{insert-length} :$

$\text{Length } n \text{ xs} \rightarrow \text{Length } (\text{suc } n) (\text{insert } \text{x} \text{ xs})$

$\text{insert-sorted} :$

$\text{Sorted } b \text{ xs} \rightarrow \text{Sorted } (b \sqcap \text{x}) (\text{insert } \text{x} \text{ xs})$

Sorted vectors

data SList : Nat → Set where

nil : $\forall \{b\} \rightarrow$ SList b

cons : (x : Nat) → $\forall \{b\} \rightarrow b \leq x \rightarrow$
SList x → SList b

data Vec Nat : Nat → Set where

[] : Vec Nat 0

:: : Nat →

$\forall \{n\} \rightarrow$ Vec Nat n → Vec Nat (suc n)

Sorted vectors

= sorted lists + vectors!

data SVec : Nat → Nat → Set where

nil : $\forall \{b\} \rightarrow$ SVec b 0

cons : (x : Nat) → $\forall \{b\} \rightarrow b \leq x \rightarrow$

$\forall \{n\} \rightarrow$ SVec x n → SVec b (suc n)

Constraints

Multiple constraints

Internalism

Ornaments

Ornament fusion

induce



corresponds to



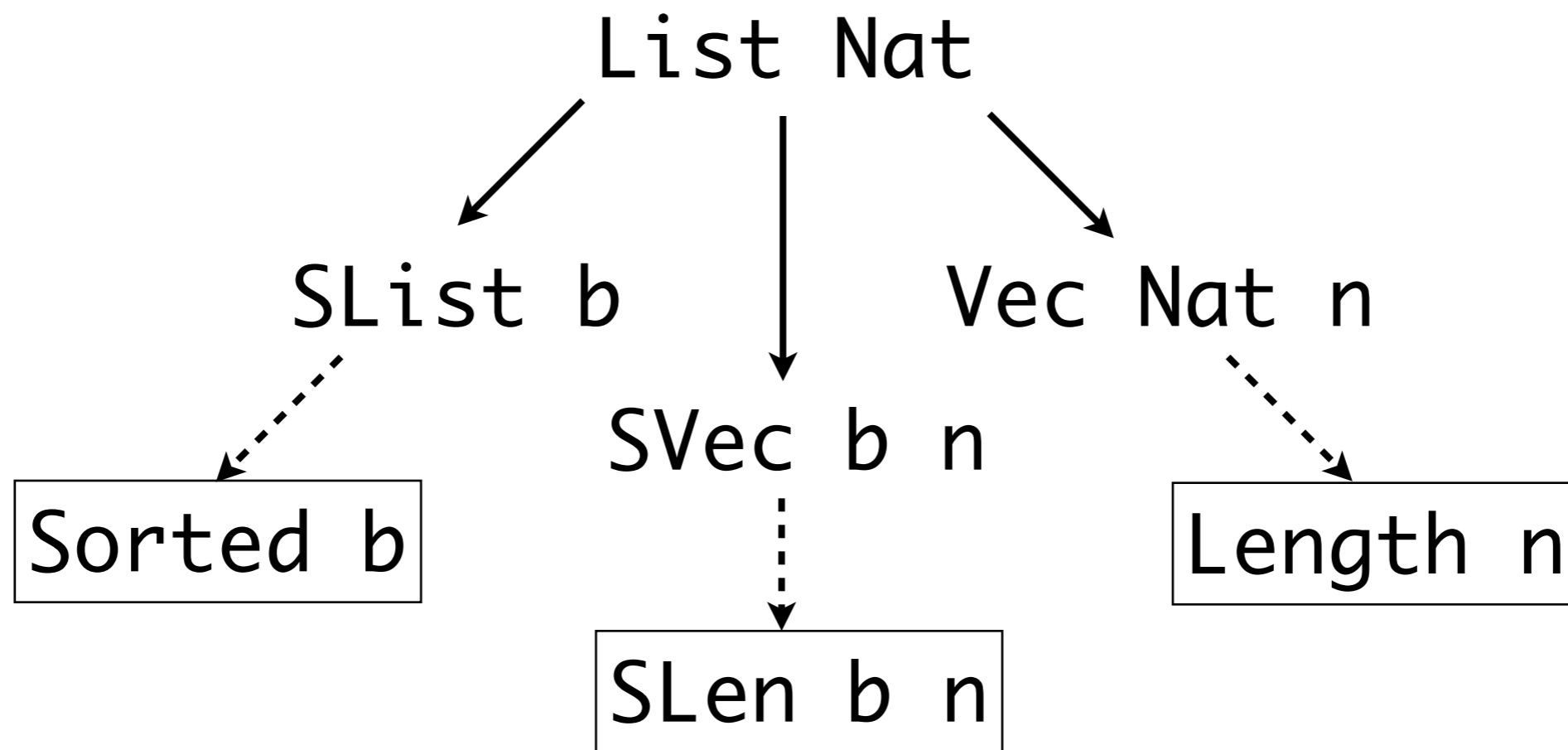
Externalism

Predicates

Pointwise conjunction
of predicates

Ornament fusion

corresponds to conjunction of induced predicates



$$\text{SLen } b \ n \ xs \cong \text{Sorted } b \ xs \times \text{Length } n \ xs$$

Ornament fusion

corresponds to conjunction of induced predicates

$SVec\ b\ n$

$\cong (xs : List\ Nat) \times SLen\ b\ n\ xs$

$\cong (xs : List\ Nat) \times Sorted\ b\ xs$
 $\quad \times Length\ n\ xs$

Function upgrade

with the help of the isomorphisms

$$\text{SVec } b \ n \cong (\text{xs} : \text{List Nat}) \times \text{Sorted } b \ \text{xs} \\ \times \text{Length } n \ \text{xs}$$

$$\text{svinsert} : (x : \text{Nat}) \rightarrow \\ \text{SVec } b \ n \rightarrow \text{SVec } (b \sqcap x) \ (\text{suc } n)$$

|||

|||

$$\text{xs} : \text{List Nat} \mapsto \text{insert } x \ \text{xs} : \text{List Nat}$$

$$s : \text{Sorted } b \ \text{xs} \mapsto \text{insert-sorted } s : \\ \text{Sorted } (b \sqcap x) \ (\text{insert } x \ \text{xs})$$

$$l : \text{Length } n \ \text{xs} \mapsto \text{insert-length } l : \\ \text{Length } (\text{suc } n) \ (\text{insert } x \ \text{xs})$$

Modular library structure

```
data List  
..., insert, ...
```

```
orn Vec inducing Length  
..., insert-length, ...
```

```
orn SList inducing Sorted  
..., insert-sorted, ...
```



Constraints

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corresponds to



Externalism

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Pointwise conjunction
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Thanks!

Descriptions

A “universe” datatype containing codes for datatypes

$\text{Vec} : \text{Set} \rightarrow \text{Desc Nat}$

$\text{Vec } A =$

$\sigma \text{ Bool } (\text{false} \mapsto \text{say zero}$

$\text{true} \mapsto \sigma A \quad \lambda x \rightarrow$

$\sigma \text{ Nat } \lambda n \rightarrow$

$\text{ask } n * \text{say } (\text{suc } n))$

$\mu : \text{Desc } I \rightarrow (I \rightarrow \text{Set})$

-- $\mu (\text{Vec } A) : \text{Nat} \rightarrow \text{Set}$

Ornaments

A richer universe of *relative* descriptions

`Vec0 : Set → Orn Nat ...`

`Vec0 A =`

`σ Bool (false ↦ say (ok zero))`

`true ↦ σ A λ x →`

`Δ Nat λ n →`

`ask (ok n) * say (ok (suc n))`)

`L_⊥ : Orn J ... → Desc J`

`-- L Vec0 A ⊥ ≈ Vec A`