

Relational ornaments

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Ornamentation

```
data List (A : Set) : Set where
```

```
  List A ∃ nil
```

```
    | cons (x : A) (xs : List A)
```

```
data Vec (A : Set) : Nat → Set where
```

```
  Vec A zero      ∃ nil
```

```
  Vec A (suc n) ∃ cons (x : A) (xs : Vec A n)
```

Ornamentation

```
data List (A : Set) : Set where
```

```
List A ∃ nil
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```
| cons (x : A) (xs : List A)
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```
data Vec (A : Set) : Nat → Set where
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```
Vec A zero ∃ nil
```

```
| cons
```

```
Vec A (suc n) ∃ nil
```

```
| cons (x : A) (xs : Vec A n)
```

Descriptions

data Desc (I : Set) : Set₁

ListD : Set → Desc T

data μ (D : Desc I) : I → Set

List A = μ (ListD A) tt : Set

Ornaments

```
data Orn {I : Set} (J : Set)
  (e : J → I)
  (D : Desc I) : Set1
```

```
Vec0 : (A : Set) → Orn Nat ! (ListD A)
```

```
[_] : Orn J e D → Desc J
```

```
Vec A = μ [ Vec0 A ] : Nat → Set
```

```
forget : (O : Orn J e D) →
  μ [ O ] j → μ D (e j)
```

```
toList = forget (Vec0 A) : Vec A n → List A
```

Relational Ornaments

```
data Orn {I : Set} {J : Set}
  (e : J → I)
  (D : Desc I)
  (E : Desc J) : Set1
```

```
VecD : Set → Desc Nat
```

```
ListD-VecD : (A : Set) →
  Orn ! (ListD A) (VecD A)
```

```
forget : (O : Orn e D E) →  $\mu$  E j →  $\mu$  D (e j)
```

```
toList =
```

```
forget (ListD-VecD A) : Vec A n → List A
```

Theory of parallel composition

Ornaments induce families of isomorphisms ...

$$\begin{array}{c} \text{toList} \\ \curvearrowright \\ \text{Vec } A \ n \cong \Sigma(xs : \text{List } A) \ \text{Length } n \ xs \end{array}$$

... and they can be composed.

$$\text{OrdList } b \cong \Sigma(xs : \text{List } \text{Val}) \ \text{Ordered } b \ xs$$

$$\text{OrdVec } b \ n \cong \Sigma(xs : \text{List } \text{Val}) \\ \text{Ordered } b \ xs \times \text{Length } n \ xs$$

Ordered vectors

```
data OrdList : Val → Set where  
  OrdList b ∃ nil  
  | cons (x : Val) (l : b ≤ x)  
    (xs : OrdList x)
```



```
data Vec Val : Nat → Set where  
  Vec Val zero ∃ nil  
  Vec Val (suc n) ∃ cons (x : Val)  
    (xs : Vec Val n)
```

Ordered vectors

```
data OrdVec : Val → Nat → Set where
  OrdVec b zero    ∃ nil
  OrdVec b (suc n) ∃ cons (x : Val)
                        (l : b ≤ x)
                        (xs : OrdVec x n)
```

Orderedness predicate

```
data OrdList : Val → Set where
  OrdList b ∃ nil
  | cons (x : Val) (l : b ≤ x)
    (xs : OrdList x)
```



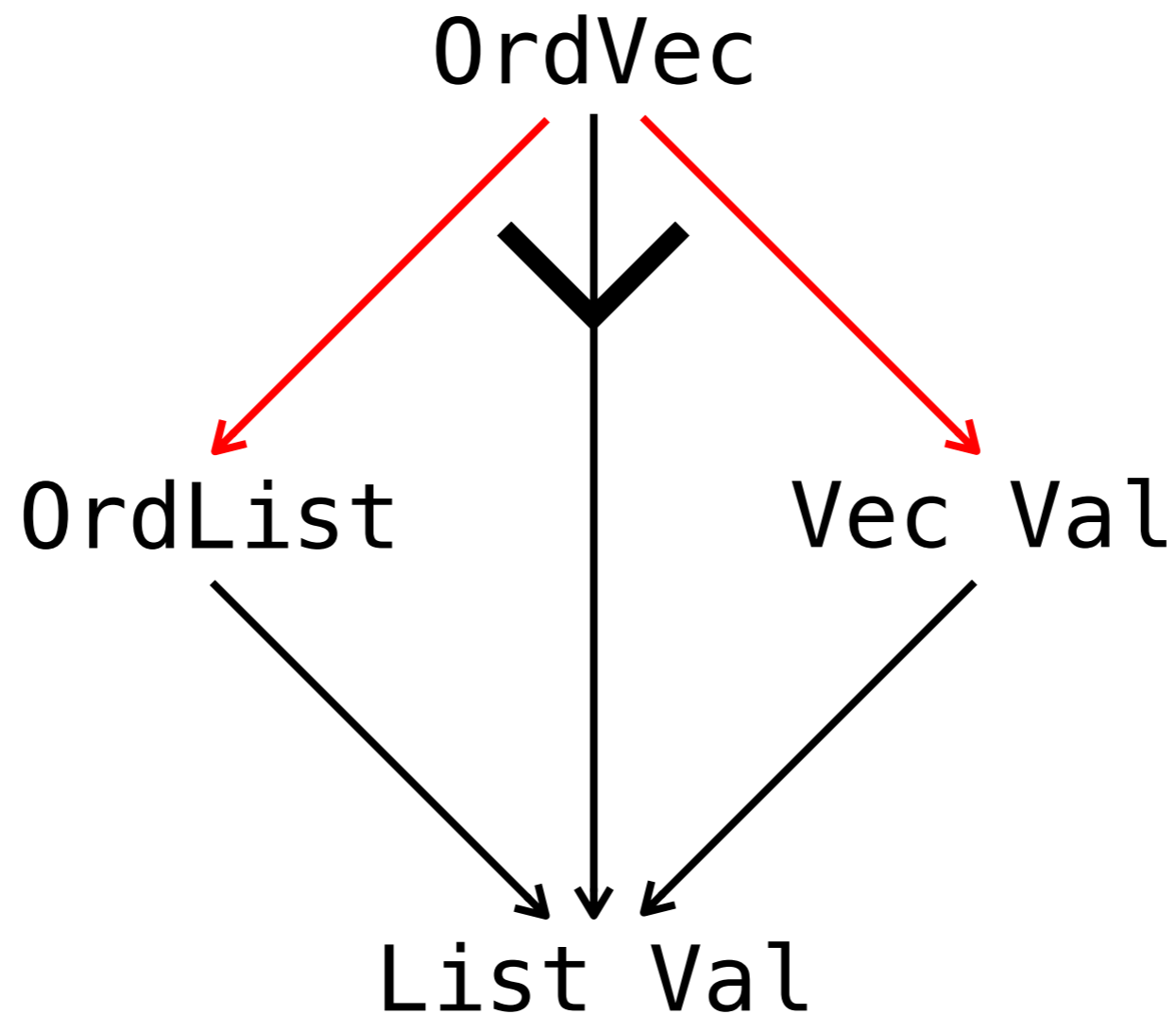
```
data ListS : List Val → Set where
  ListS nil ∃ nil
  ListS (cons x xs) ∃ cons (x : Val)
    (s : ListS xs)
```

Orderedness predicate

```
data Ordered : Val → List Val → Set where
  Ordered b nil           ∃ nil
  Ordered b (cons x xs)
    ∃ cons (x : Val) (l : b ≤ x)
          (ls : Ordered x xs)
```

Parallel composition

is a pullback in the category of descriptions and ornaments



sequential
composition

$\circ : \text{Orn } e \text{ } D \text{ } E \rightarrow \text{Orn } f \text{ } E \text{ } F \rightarrow \text{Orn } (e \circ f) \text{ } D \text{ } F$

From the pullback property

we can derive

$$\text{Vec } A \ n \cong \Sigma(xs : \text{List } A) \ \text{Length } n \ xs$$

$$\text{OrdList } b \cong \Sigma(xs : \text{List } \text{Val}) \ \text{Ordered } b \ xs$$

$$\begin{aligned} \text{OrdVec } b \ n &\cong \Sigma(xs : \text{List } \text{Val}) \\ &\quad \text{OrderedLength } b \ n \ xs \\ &\cong \Sigma(xs : \text{List } \text{Val}) \\ &\quad \text{Ordered } b \ xs \times \text{Length } n \ xs \end{aligned}$$

whose constructions are independent of choices of universe encodings.

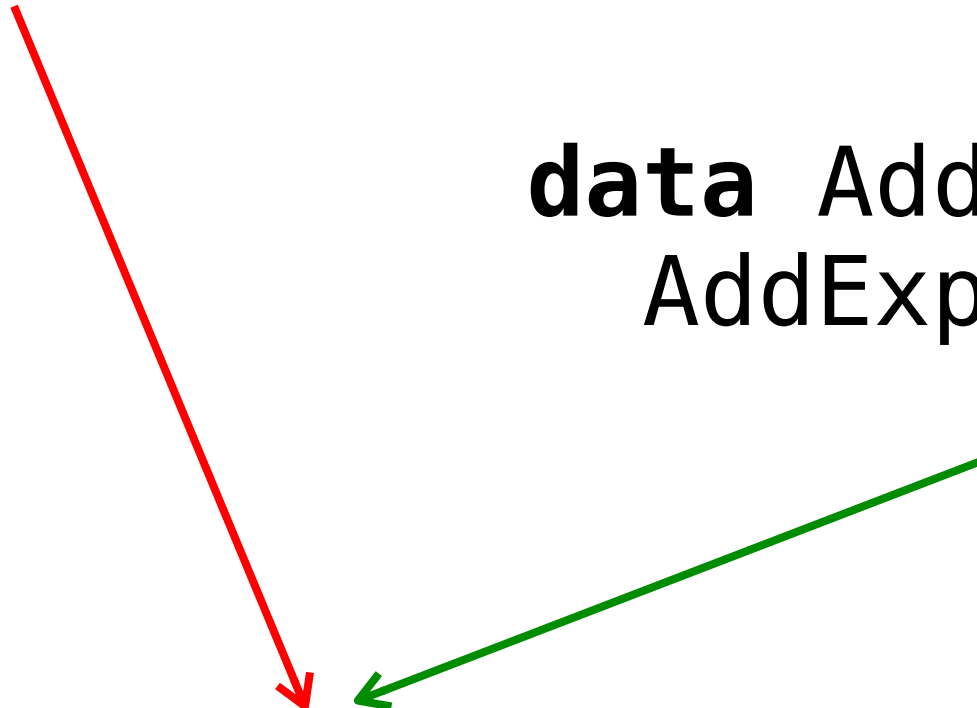
Datatypes à la carte

Coproduct of signatures

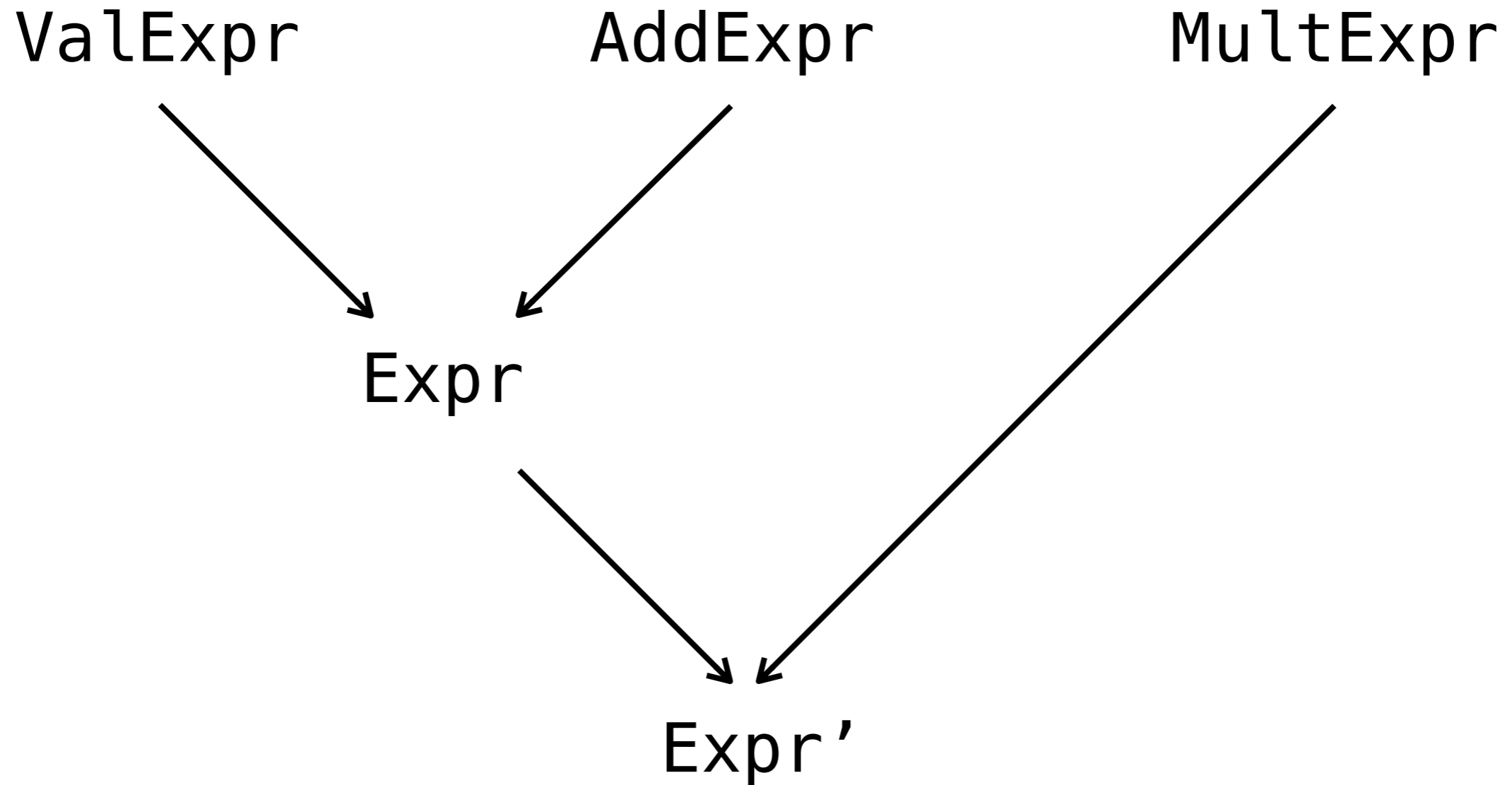
```
data ValExpr : Set where  
  ValExpr ∋ val (n : Nat)  
           | add
```

```
data AddExpr : Set where  
  AddExpr ∋ val  
           | add (l : AddExpr)  
              (r : AddExpr)
```

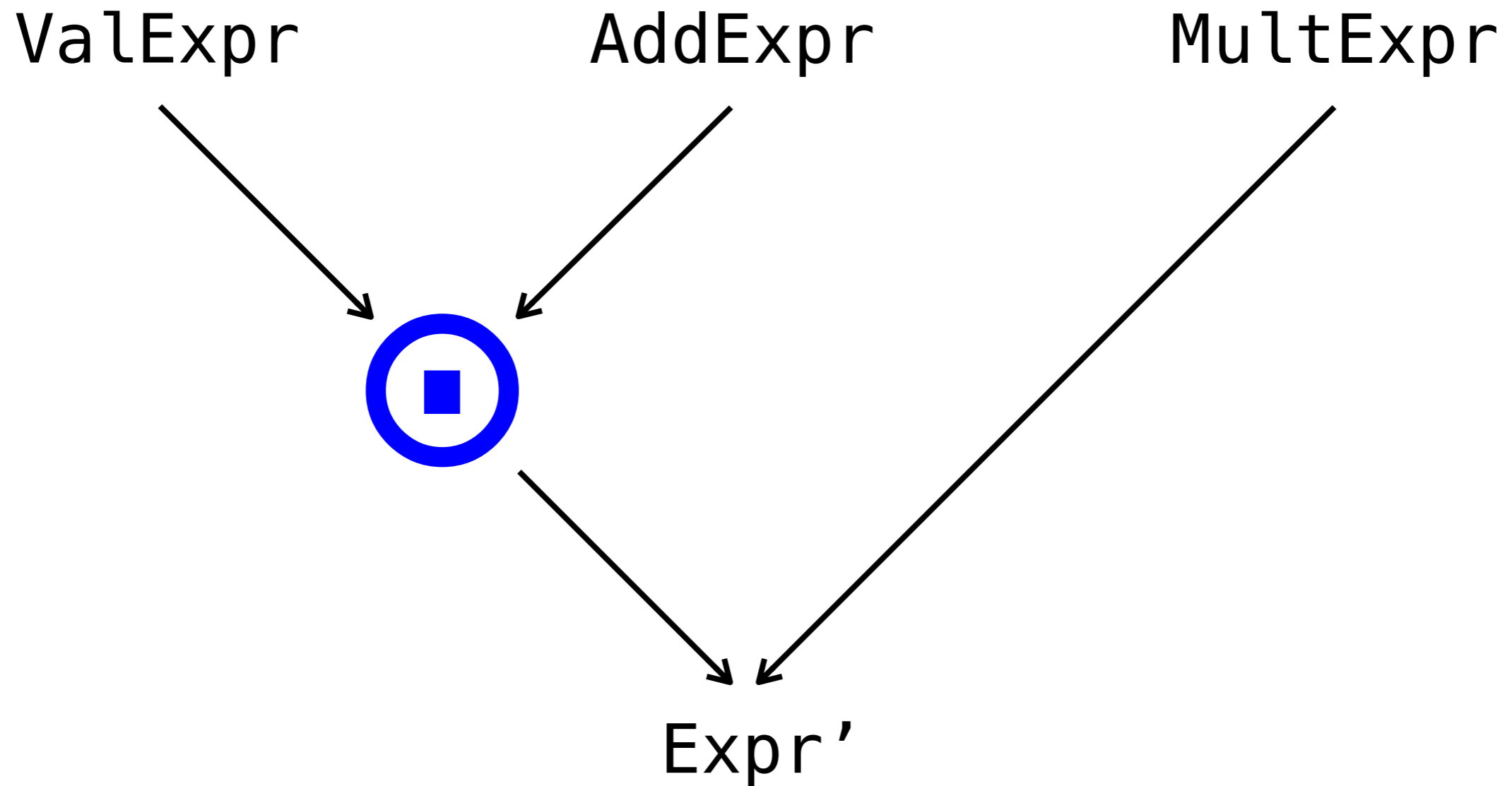
```
data Expr : Set where  
  Expr ∋ val (n : Nat)  
       | add (l : Expr) (r : Expr)
```



Even more signatures



Even more signatures



Smart constructors

AddExprD : Desc T

$_+_ : \{\{_ : \text{Orn ! D AddExprD}\}\} \rightarrow$
 $\mu D tt \rightarrow \mu D tt \rightarrow \mu D tt$

Summary

With relational ornaments, **creation of datatypes** is **decoupled** from **specification of ornamental relationship** between datatypes.

- **Ornaments can then arise between existing datatypes,**
 - leading to a nice theory of parallel composition.
- **New datatypes can arise at the less informative end,**
 - *so datatypes à la carte* may be subsumed in the theory of ornaments.