

Interpolation for guarded logics

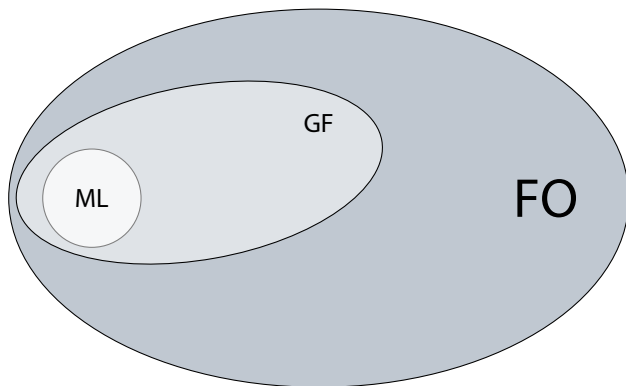
Michael Vanden Boom

University of Oxford

Highlights 2014
Paris, France

Joint work with
Michael Benedikt and Balder ten Cate

Some guarded logics

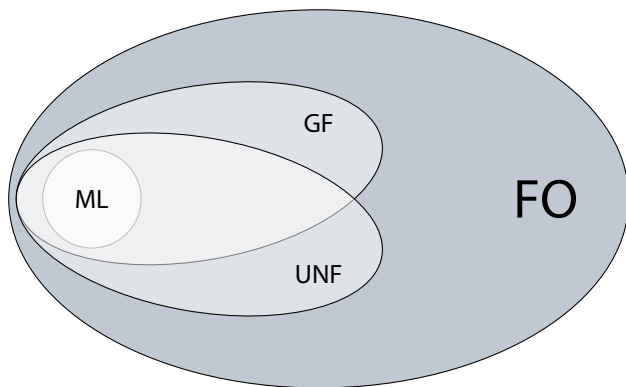


constrain
quantification

$$\begin{aligned} &\exists x(G(xy) \wedge \psi(xy)) \\ &\forall x(G(xy) \rightarrow \psi(xy)) \end{aligned}$$

[Andréka, van Benthem,
Németi '95-'98]

Some guarded logics



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$$\forall x(G(xy) \rightarrow \psi(xy))$$

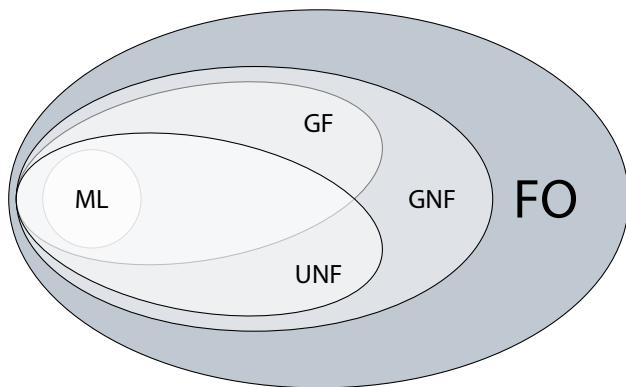
[Andréka, van Benthem,
Németi '95-'98]

constrain
negation

$$\exists x(\psi(xy))$$
$$\neg\psi(x)$$

[ten Cate, Segoufin '11]

Some guarded logics



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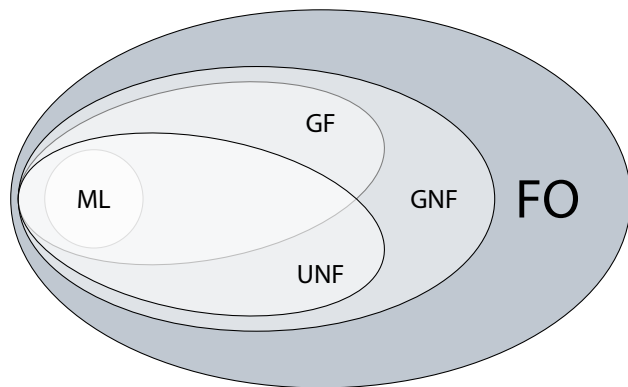
[Andréka, van Benthem,
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$$\exists x(\psi(xy))$$
$$G(xy) \wedge \neg \psi(xy)$$

[ten Cate, Segoufin '11]
[Bárány, ten Cate, Segoufin '11]

Some guarded logics



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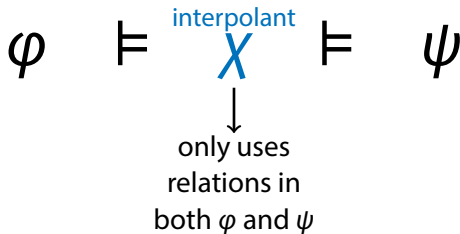
[ten Cate, Segoufin '11]

[Bárány, ten Cate, Segoufin '11]

These guarded logics are **decidable**, and **expressive** enough to capture many query languages and integrity constraints of interest in databases and knowledge representation.

$$\varphi \models \psi$$

Interpolation



Interpolation example

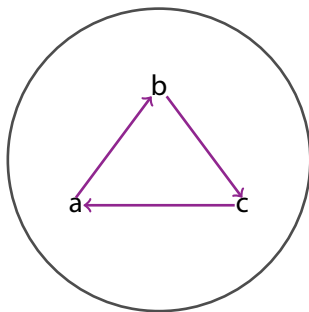
$$\exists xyz(T_{xyz} \wedge R_{xy} \wedge R_{yz} \wedge R_{zx}) \quad \models \quad \exists xy(R_{xy} \wedge ((S_x \wedge S_y) \vee (\neg S_x \wedge \neg S_y)))$$

“there is a T -guarded
3-cycle using R ”

Interpolation example

$$\exists xyz(Txyz \wedge Rxy \wedge Ryx \wedge Rzx) \quad \models \quad \exists xy(Rxy \wedge ((Sx \wedge Sy) \vee (\neg Sx \wedge \neg Sy)))$$

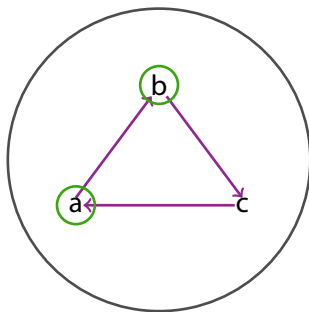
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Interpolation example

$$\exists xyz(T_{xyz} \wedge R_{xy} \wedge R_{yz} \wedge R_{zx}) \quad \models \quad \exists xy(R_{xy} \wedge ((S_x \wedge S_y) \vee (\neg S_x \wedge \neg S_y)))$$

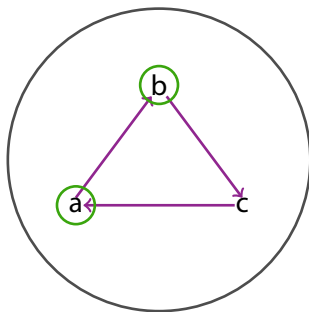
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Interpolation example

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“there is a T -guarded
3-cycle using R ”



interpolant $\chi := \exists xyz(R_{xy} \wedge R_{yz} \wedge R_{zx})$

“there is a 3-cycle using R ”

Why do we care?

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Why might someone care?

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- interpolation is a **benchmark** property of modal logic

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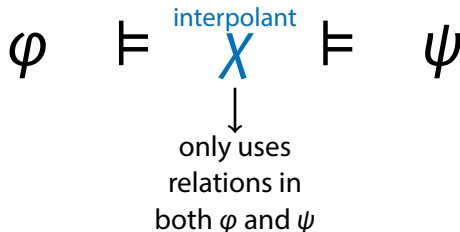
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- interpolation is a **benchmark** property of modal logic
- interpolation implies the **Beth definability** property
(implicit definability = explicit definability)
which indicates a good balance between syntax and semantics

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Why might someone care?

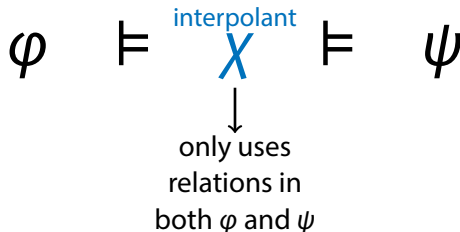
- interpolation is a **benchmark** property of modal logic
- interpolation implies the **Beth definability** property
(implicit definability = explicit definability)
which indicates a good balance between syntax and semantics
- for these guarded logics with connections to databases,
interpolation is related to **query rewriting** over views



Theorem (ten Cate, Segoufin '11; Barany, Benedikt, ten Cate '13)

Given GNF (respectively, UNF) formulas φ and ψ such that $\varphi \models \psi$, there is a GNF (respectively, UNF) interpolant χ .

Interpolation

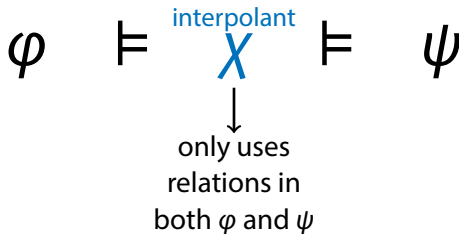


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Given GNF (respectively, UNF) formulas φ and ψ such that $\varphi \models \psi$, there is a GNF (respectively, UNF) interpolant χ .

No idea how to **compute** interpolants
(or other rewritings related to interpolation).

Interpolation



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Given GNF (respectively, UNF) formulas φ and ψ such that $\varphi \models \psi$, there is a GNF (respectively, UNF) interpolant χ .

Theorem (Benedikt, ten Cate, VB. '14)

Given GNF (respectively, UNF) formulas φ and ψ s.t. $\varphi \models \psi$, we can construct a GNF (respectively, UNF) interpolant χ of doubly exponential DAG-size.

Conclusion

	ML	GF	UNF	GNF
Interpolation?	✓	✗	✓	✓

adapted
mosaic method
from ML

[Benedikt,ten Cate,VB.'14]

Conclusion

	ML	GF	UNF	GNF	L_μ	GFP	UNFP	GNFP
Interpolation?	✓	✗	✓	✓	✓	✗	✓	✗

adapted
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[Benedikt,ten Cate,VB.'14]

Conclusion

	ML	GF	UNF	GNF	L_μ	GFP	UNFP	GNFP
Interpolation?	✓	✗	✓	✓	✓	✗	✓	✗
		adapted mosaic method from ML [Benedikt,ten Cate,VB.'14]				used automata for L_μ [Benedikt,ten Cate,VB. unpublished]		

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Interpolation?	✓	✗	✓	✓	✓	✗	✓	✗
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Open question

Is there a decidable extension of GNFP that has interpolation?