

LECTURE 16: VENDING MACHINE CASE STUDY

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- Next, we need a type *PROD*, representing all the products that the machine can sell.

[*PROD*]

- We can define the state space of the vending machine thus:

<i>VendingMachine</i>
$cost : PROD \rightarrow \mathbb{N}$
$stock : \text{bag } PROD$
$float : \text{bag } COIN$
$\text{dom } stock \subseteq \text{dom } cost$

1 Specification of a Vending Machine

- In this lecture, we will give a complete specification of a vending machine – the sort you buy cans of coke or cigarettes from.
- First, we need to introduce some types; the first one will be *COIN*, representing all the coins that are accepted by the machine.

$COIN == \{100, 50, 20, 10, 5, 2, 1\}$

- That is, there are coins in denominations of 100, 50, 20, 10, 5, 2, and 1 pence.
 - We will also need a type for system messages
- this is parachuted in:

[*REPORT*]

- The function *cost* return the cost of a product in pence. For example,

$cost(MarsBar) = 25$
 $cost(Penguin) = 15$

- The bag *stock* tells us how many items of each type are in stock. For example,

$stock = \{Penguin \mapsto 2\}$

means that there are just 2 penguins in the machine.

- The bag *float* records the coins that are currently in the machine; for example

$float = \{100 \mapsto 2, 50 \mapsto 8, 5 \mapsto 20\}$

means that there are $2 \times £1$ coins, $8 \times 50p$ coins and $20 \times 5p$ coins.

- QUESTION: Why are *stock* and *float* bags and not sets or sequences?
- The invariant $\text{dom } stock \subseteq \text{dom } cost$ says that everything in the machine (i.e. in stock) must have a cost associated with it.

Operations

Here are the operations we shall specify:

- initialising the machine;
- pricing goods;
- restocking;
- buying goods.

Pricing Goods

- This simply means changing the price of an item in stock, or pricing an item that is going to be stocked.
- The inputs are the item and a price.

$\begin{array}{l} \textit{Price} \\ \Delta \textit{VendingMachine} \\ \textit{item}? : \textit{PROD} \\ \textit{price}? : \mathbb{N} \\ \hline \textit{cost}' = \textit{cost} \oplus \{ \textit{item}? \mapsto \textit{price}? \} \\ \textit{stock}' = \textit{stock} \\ \textit{float}' = \textit{float} \end{array}$
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Initialisation

$\begin{array}{l} \textit{InitVendingMachine} \\ \Delta \textit{VendingMachine} \\ \hline \textit{cost}' = \{ \} \\ \textit{stock}' = \square \\ \textit{float}' = \square \end{array}$

- So initially, the machine does not know the cost of anything, contains nothing, and has no float.

Restocking

- The next operation to specify is that of restocking the machine with more goods.
- The only input is a new bag of products.
- The precondition $\text{dom } \textit{new}? \subseteq \text{dom } \textit{cost}$ is implied by the invariant of $\textit{VendingMachine}'$.

$\begin{array}{l} \textit{Restock} \\ \Delta \textit{VendingMachine} \\ \textit{new}? : \text{bag } \textit{PROD} \\ \hline \textit{stock}' = \textit{stock} \uplus \textit{new}? \\ \textit{float}' = \textit{float} \\ \textit{cost}' = \textit{cost} \end{array}$
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- (Note that \uplus is the 'bag union' operator.)

- We shall now make the operation robust. The *Restock* operation fails when an attempt is made to add goods which are not priced. We need a schema to identify this situation.

$$\frac{\text{GoodsNotPriced} \quad \frac{\frac{\exists \text{VendingMachine} \quad \text{new?} : \text{bag PROD} \quad \text{rep!} : \text{REPORT}}{\neg(\text{dom new?} \subseteq \text{dom cost})} \quad \text{rep!} = \text{'Some goods are not priced'}}{\text{rep!} = \text{'Some goods are not priced'}}$$

- This schema expands to ...

$$\frac{\text{RestockOp} \quad \frac{\frac{\Delta \text{VendingMachine} \quad \text{new?} : \text{bag PROD} \quad \text{rep!} : \text{REPORT}}{\text{cost}' = \text{cost} \quad \text{float}' = \text{float} \quad (\text{stock}' = \text{stock} \uplus \text{new?} \wedge \text{rep!} = \text{'Okay'})} \quad \vee \quad (\neg(\text{dom new?} \subseteq \text{dom cost}) \wedge \text{stock}' = \text{stock} \wedge \text{rep!} = \text{'Some goods are not priced'})}{\text{rep!} = \text{'Some goods are not priced'}}$$

- We need an operation to report success...

$$\frac{\text{Success} \quad \frac{\text{rep!} : \text{REPORT}}{\text{rep!} = \text{'Okay'}}}{\text{rep!} = \text{'Okay'}}$$

- Now, we simply use the schema calculus to specify a robust version of the Restock operation, called *RestockOp*:

$$\text{RestockOP} \equiv (\text{Restock} \wedge \text{Success}) \vee \text{GoodsNotPriced}$$

Buying

- The buying operation is a somewhat more complex operation ...
- The inputs are the chosen item and some money.
- We have to check that the item is in stock, and that the user has entered enough money to buy it.
- We may also have to figure out what change to give ...

- We assume that a function $sum : bag\ COIN \rightarrow \mathbb{N}$ is available, which takes a bag of coins and calculates how much is in the bag. For example, given a bag containing $7 \times 2p$, and $3 \times 5p$ coins,

$$\begin{aligned} sum\{2 \mapsto 7, 5 \mapsto 3\} &= (2 \times 7) + (5 \times 3) \\ &= 14 + 15 \\ &= 29pence \end{aligned}$$

- $in?$ represents the coins entered; $out!$ represents the change;
- $item?$ is the item dispensed to the user;
- the 1st condition says that the item must be in stock;
- the 2nd condition says that the amount of money entered must be greater than or equal to the cost of the item;
- the 3rd condition says that the change given must have been part of the float;
- the 4th condition says that the the money entered must equal the change given plus the cost of the item;
- the 5th condition says that the stock before must be equal to the stock after, to which is added the dispensed item;
- the 6th condition says that the float after, together with the change dispensed must equal the float before plus the amount entered (i.e. no money disappears)

- The basic *Buy* operation is as follows:

Buy

$\Delta VendingMachine$

$in?, out! : bag\ COIN$

$item? : PROD$

$item? \text{ in stock}$

$sum(in?) \geq cost(item?)$

$out! \sqsubseteq float$

$sum(in?) = sum(out!) + cost(item?)$

$stock' \uplus \{item? \mapsto 1\} = stock$

$float' \uplus out? = float \uplus in?$

$cost' = cost$