

A causal Markov category with Kolmogorov products

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In [2] (Problem 6.7) the problem was posed of finding an interesting Markov category which is *causal* and has all (small) *Kolmogorov products*. Here we give an example where the deterministic subcategory is the category of Stone spaces (i.e. the dual of the category of Boolean algebras) and the kernels correspond to a restricted class of Kleisli arrows for the Radon monad.

We look at this from two perspectives. First via pro-completions and Stone spaces directly. Second via duality with Boolean algebras and effect algebras.

1 Perspective in terms of pro-completions

Let $\mathcal{C}$ be a Markov category, with deterministic wide subcategory $\det(\mathcal{C})$. A motivating example is $\mathcal{C} = \mathbf{FinKer}$, so that $\det(\mathcal{C}) = \mathbf{Fin}$, and $\mathbf{Pro}(\det(\mathcal{C}))$ is the category of Stone spaces.

We often write ‘ $X \rightarrow X'$ ’ for an object of $\mathbf{Pro}(\mathcal{C})$, as a shorthand for $X \cong \lim_{X \rightarrow X'} X'$, where the X' s are in $\mathcal{C}$.

Let $\mathbf{ProDet}(\mathcal{C})$ be the full subcategory of the profinite completion $\mathbf{Pro}(\mathcal{C})$ given by limits of cofiltered diagrams that factorize through $\det(\mathcal{C})$. Equivalently, $\mathbf{ProDet}(\mathcal{C})$ is the full image of

$$\mathbf{Pro}(\det(\mathcal{C})) \longrightarrow \mathbf{Pro}(\mathcal{C}),$$

where $\mathbf{Pro}(\det(\mathcal{C}))$ is cartesian monoidal. $\mathbf{ProDet}(\mathcal{C})$ is monoidal with

$$X \otimes Y = \lim_{\substack{X \rightarrow X' \\ Y \rightarrow Y' \\ \text{in } \mathbf{Pro}(\det(\mathcal{C}))}} X' \otimes Y',$$

i.e. the cartesian product $X \times Y$ in $\mathbf{Pro}(\det(\mathcal{C}))$. Every $X \in \mathbf{ProDet}(\mathcal{C})$ has a copy structure given by

$$X \longrightarrow X' \xrightarrow{\text{copy}} X' \otimes X'$$

for each $X \rightarrow X'$ in $\mathbf{Pro}(\det(\mathcal{C}))$.

Lemma 1. $\mathbf{ProDet}(\mathcal{C})$ is a Markov category with deterministic subcategory $\mathbf{Pro}(\det(\mathcal{C}))$.

Proof. The monoidal unit is the terminal object $1 \in \mathcal{C}$. Clearly every map in $\mathbf{Pro}(\det(\mathcal{C}))$ is deterministic. Let $f : X \rightarrow Y$ be deterministic in $\mathbf{ProDet}(\mathcal{C})$.

Then for every $Y \rightarrow Y'$ in $\text{Pro}(\det(\mathcal{C}))$, there is an $X \rightarrow X'$ in $\text{Pro}(\det(\mathcal{C}))$ such that

$$X \xrightarrow{f} Y \rightarrow Y' = X \rightarrow X' \xrightarrow{f'} Y'$$

for some $f' : X' \rightarrow Y'$ in $\mathcal{C}$, whence we have equations

$$\begin{aligned} X &\xrightarrow{f} Y \xrightarrow{\text{copy}} Y \otimes Y \rightarrow Y' \otimes Y' \\ &= X \xrightarrow{f} Y \rightarrow Y' \xrightarrow{\text{copy}} Y' \otimes Y' \\ &= X \rightarrow X' \xrightarrow{f'} Y' \xrightarrow{\text{copy}} Y' \otimes Y', \end{aligned}$$

and

$$\begin{aligned} X &\xrightarrow{\text{copy}} X \otimes X \xrightarrow{f \otimes f} Y \otimes Y \rightarrow Y' \otimes Y' \\ &= X \xrightarrow{\text{copy}} X \otimes X \rightarrow X' \otimes X' \xrightarrow{f' \otimes f'} Y' \otimes Y' \\ &= X \rightarrow X' \xrightarrow{\text{copy}} X' \otimes X' \xrightarrow{f' \otimes f'} Y' \otimes Y'. \end{aligned}$$

Thus for some factorization $X \rightarrow X'' \rightarrow X'$ in $\text{Pro}(\det(\mathcal{C}))$ we have

$$X'' \rightarrow X' \xrightarrow{f'} Y' \xrightarrow{\text{copy}} Y' \otimes Y' = X'' \rightarrow X' \xrightarrow{\text{copy}} X' \otimes X' \xrightarrow{f' \otimes f'} Y' \otimes Y',$$

whence we deduce that

$$X'' \rightarrow X' \xrightarrow{f'} Y'$$

is in $\det(\mathcal{C})$. Since $Y \rightarrow Y'$ was arbitrary, the map f is in $\text{Pro}(\det(\mathcal{C}))$. $\square$

Theorem 2. $\text{ProDet}(\mathcal{C})$ has all small Kolmogorov products.

Proof. Given $(X_i : i \in I)$, let X be their product in $\text{Pro}(\det(\mathcal{C}))$, as an object of $\text{ProDet}(\mathcal{C})$. Then X is an infinite tensor product of the X_i , since $\text{Pro}(\det(\mathcal{C})) \rightarrow \text{ProDet}(\mathcal{C})$ preserves cofiltered limits, essentially by construction. $\square$

In the case of $\text{ProDet}(\text{FinKer})$ there is also a faithful functor $\text{ProDet}(\text{FinKer}) \rightarrow \text{Kleisli}(R)$ into the Kleisli category of the Radon monad on compact Hausdorff spaces. On objects, this maps a cofiltered diagram to its limit in compact Hausdorff spaces: a Stone space. Recall that a morphism $f : \lim_i X_i \rightarrow \lim_j Y_j$ in $\text{ProDet}(\text{FinKer})$ is given by, for each j , a choice i and a morphism in $\text{FinKer}(X_i, Y_j)$ making suitable diagrams commute, and modulo choice of i . Thus for each deterministic point $x : 1 \rightarrow \lim_i X_i$, i.e. point of the Stone space, we have a morphism $p_j(x) \in \text{FinKer}(1, Y_j)$ for all j , i.e. a distribution $p_j(x)$ on Y_j for all j , and these are all suitably compatible; we then construct a function $p : \lim_i X_i \rightarrow R(\lim_j Y_j)$ between compact Hausdorff spaces, mapping x to the Kolmogorov extension of $p_j(x)$ on the Stone space $\lim_j Y_j$.

From [1, Example 11.35], the Markov category of all probability kernels between measurable spaces is causal. And so, from this faithful embedding, we see that $\text{ProDet}(\text{FinKer})$ is causal.

We revisit this in the Section 2, where the algebra brings out a different angle, and also suggests a direct proof of causality.

1.1 Aside: Additional basic results

Lemma 3. *If $\mathcal{C}$ is representable, then so is $\text{ProDet}(\mathcal{C})$.*

Proof. Suppose $\text{det}(\mathcal{C}) \rightarrow \mathcal{C}$ has a right adjoint $D : \mathcal{C} \rightarrow \text{det}(\mathcal{C})$. By 2-functoriality of Pro , this gives a right adjoint D^+ to $\text{Pro}(\text{det}(\mathcal{C})) \rightarrow \text{Pro}(\mathcal{C})$:

$$D^+X = \lim_{X \rightarrow X' \in \text{Pro}(\text{det}(\mathcal{C}))} DX'.$$

□

Lemma 4. *If $\mathcal{C}$ is a.s.-compatibly representable, then so is $\text{ProDet}(\mathcal{C})$.*

Proof. Let $p : T \rightarrow A$ and $f, g : A \rightarrow X$ be maps in $\text{ProDet}(\mathcal{C})$, and suppose that $f \equiv_{p\text{-a.s.}} g$, i.e.

$$T \xrightarrow{p} A \rightarrow A \otimes A \xrightarrow{\text{id} \otimes f} A \otimes X = T \xrightarrow{p} A \rightarrow A \otimes A \xrightarrow{\text{id} \otimes g} A \otimes X,$$

or equivalently,

$$\begin{aligned} T &\xrightarrow{p} A \rightarrow A \otimes A \xrightarrow{\text{id} \otimes f^\sharp} A \otimes D^+X \xrightarrow{\text{id} \otimes \text{samp}} A \otimes X \\ &= T \xrightarrow{p} A \rightarrow A \otimes A \xrightarrow{\text{id} \otimes g^\sharp} A \otimes D^+X \xrightarrow{\text{id} \otimes \text{samp}} A \otimes X. \end{aligned}$$

We are required to show that we can cancel the factor of $\text{id} \otimes \text{samp}$ from each side. It is straightforward to show that this is possible upon post-composing with projections $A \otimes X \rightarrow A' \otimes X'$. □

(Note: Lemmas 3 and 4 are not relevant to FinKer .)

Lemma 5. *Let $\mathcal{C} = \text{FinKer}$. Then for $X, Y \in \text{FinKer}$, $L \in \text{ProDet}(\text{FinKer})$, every kernel $p : X \rightsquigarrow Y \otimes L$ admits a conditional $k : X \otimes Y \rightsquigarrow L$.*

Proof. Without loss of generality, let $X = 1$. For $y \in Y$, the value of $p(*) (y, \top^Z)$ is independent of the deterministic projection $\pi : L \rightarrow Z \in \text{Pro}(\text{Fin})$ chosen. If it is zero, we define $k(*, y)$ arbitrarily. If it is not zero, say $0 < K \leq 1$, then for each $\pi : L \rightarrow Z$ we have a K -valued measure $p(*) (y, -)$ on Z forming a cone of kernels over the diagram for L . Normalizing by K , this is still a cone and hence we have defined a kernel $k : 1 \otimes Y \rightsquigarrow L$.

We must check that this really is the conditional, but this is straightforward. □

2 Dual perspective in terms of Boolean algebras and effect algebras

Boolean algebras and effect algebras Let BAlg denote the category of Boolean algebras, whose subcategory BAlg_{fin} of finite algebras is equivalent to Fin^{op} , the dual of the category of finite sets and all functions. We recall some generalities. The embedding

$$\text{BAlg} \hookrightarrow [\text{BAlg}_{\text{fin}}^{\text{op}}, \text{Set}] \simeq [\text{Fin}, \text{Set}]$$

is full and faithful, and the essential image is the closure of the representable functors under filtered colimits.

The category **EAlg** of *effect algebras* is traditionally presented in terms of partial algebras, but it is sometimes helpful to use the following viewpoint: the mapping

$$(E, \otimes, 0, 1) \mapsto \lambda n. \{(e_1, \dots, e_n) \in E^n : e_1 \otimes \dots \otimes e_n = 1\}$$

embeds **EAlg** as a full reflective subcategory of $[\mathbf{Fin}, \mathbf{Set}]$ containing **BAlg**

$$\mathbf{BAlg} \hookrightarrow \mathbf{EAlg} \hookrightarrow [\mathbf{Fin}, \mathbf{Set}].$$

We won't refer to the precise limit-preservation conditions that characterize **EAlg** this way [4, 5]. The cocartesian monoidal structure on finite Boolean algebras extends by Day convolution to a symmetric monoidal structure $\boxplus$ on $[\mathbf{Fin}, \mathbf{Set}]$, which restricts to the coproduct on **BAlg**. It's not obvious, but the reflection of the Day convolution into **EAlg** gives the classical effect algebra tensor $\otimes$ when restricted to **EAlg**. While the inclusion $\mathbf{BAlg} \hookrightarrow \mathbf{EAlg}$ does not preserve coproducts in general, it does preserve the initial Boolean algebra $\mathbb{2} = \{\perp, \top\}$. This makes **EAlg** an interesting semicocartesian category.

The interval effect monoid Let $\mathbb{I}$ denote the effect algebra whose underlying set is the real interval $[0, 1]$ with partial sum $a \oplus b$ defined whenever $a + b \leq 1$ in which case it is equal to $a + b$. As a functor $\mathbf{Fin} \rightarrow \mathbf{Set}$,

$$\mathbb{I}(n) = \{\phi \in [0, 1]^n : \sum_{i=1}^n \phi(i) = 1\}$$

and the action of a function $f : m \rightarrow n$ sends $\phi \in \mathbb{I}(m)$ to

$$\lambda i \in n. \sum_{f(j)=i} \phi(j).$$

As is well-known, $\mathbb{I}$ is a monoid with respect to $\otimes$, where the multiplication can be represented in terms of the Day convolution

$$\int^{a, b \in \mathbf{Fin}} \mathbf{Fin}(a \times b, n) \times \mathbb{I}(a) \times \mathbb{I}(b) \cong (\mathbb{I} \boxplus \mathbb{I})(n) \rightarrow \mathbb{I}(n)$$

by the dinatural transformation

$$\begin{aligned} \mathbf{Fin}(a \times b, n) \times \mathbb{I}(a) \times \mathbb{I}(b) &\rightarrow \mathbb{I}(n) \\ (f, \phi, \psi) &\mapsto \mathbb{I}(f)(\lambda(i, j) \in a \times b. \phi(i)\psi(j)) \end{aligned}$$

and the unit is just the unique map $\mathbb{2} \rightarrow \mathbb{I}$. (Recall that $\mathbb{I} \boxplus \mathbb{I}$ is not itself the effect tensor $\mathbb{I} \otimes \mathbb{I}$, rather it's an object of $[\mathbf{Fin}, \mathbf{Set}]$ whose reflection into **EAlg** is $\mathbb{I} \otimes \mathbb{I}$).

Measures on and kernels between Boolean algebras We will only work with probability measures. A measure (i.e. valuation) on a Boolean algebra A is a map $m : A \rightarrow \mathbb{I}$ in **EAlg**, i.e. a natural transformation. This is actually just the usual notion of finitely-additive measure (or modular valuation).

Definition 6. For $A, B \in \mathbf{BAlg}$, a *kernel* $k : A \rightsquigarrow B$ is a map $k : B \rightarrow \mathbb{I} \otimes A$.

Modules for the effect monoid $\mathbb{I}$ have been called “convex effect algebras” and have a representation theorem in terms of ordered vector spaces with a chosen strong unit [3]. For our purposes, the following special case is helpful.

Lemma 7. *Let $A \in \mathbf{BAlg}$. Then $\mathbb{I} \otimes A$ is isomorphic to the effect algebra $\mathbf{M}(A)$ whose elements are locally constant functions $f : S(A) \rightarrow [0, 1]$ from the Stone space $S(A)$ of A to the unit interval with the evident partial sum operation.*

Stone duality tells us that $\mathbf{BAlg}^{\text{op}}$ is equivalent to a full subcategory of $\mathbf{CHaus}$, i.e.

$$\mathbf{BAlg}(B, A) \cong \mathbf{CHaus}(S(A), S(B)).$$

It is fairly straightforward to check that measures on $A \in \mathbf{BAlg}$ are equivalent to Radon (probability) measures on $S(A)$. Thus we ‘extend’ Stone duality a little bit with the observation

$$\mathbf{EAlg}(B, \mathbb{I}) \cong \mathbf{CHaus}(1, R(S(B)))$$

where R is the Radon monad. In general, we only get an injection

$$\mathbf{EAlg}(B, \mathbb{I} \otimes A) \hookrightarrow \mathbf{CHaus}(S(A), R(S(B))).$$

The Radon kernels $k : S(A) \rightarrow R(S(B))$ that arise in this way are those with the property that, for every clopen set $C \subseteq S(B)$, the function $\lambda x \in S(A). k(x, C)$ is locally constant — in that case there is a (necessarily finite) partition of $S(A)$ into clopen sets such that the measure assigned to C is constant on each.

Example 8. A Radon kernel between Stone spaces that does not arise in this way is $b : 2^{\mathbb{N}} \rightsquigarrow 2$ where $b(\vec{x})$ is the Bernoulli distribution with bias $\sum_{i \in \mathbb{N}} 2^{-i-1} x_i$.

A Markov category of Boolean kernels

Definition 9. Let $\mathbf{BKer} = ((\mathbf{BAlg})_{\mathbb{I} \otimes (-)})^{\text{op}}$, the dual of the Kleisli category of the monad $\mathbb{I} \otimes (-)$ on $\mathbf{EAlg}$ restricted to the objects of $\mathbf{BAlg}$. Thus the objects of $\mathbf{BKer}$ are Boolean algebras and the morphism $A \rightarrow B$ are kernels, i.e. maps $B \rightarrow \mathbb{I} \otimes A$ in $\mathbf{EAlg}$.

Lemma 10. *$\mathbf{BKer}$ is a semicartesian monoidal category.*

Proof. The tensor of effect algebras straightforwardly induces a tensor of free $\mathbb{I}$ -modules. $\mathbf{BKer}$ is semicartesian because the monoidal unit is 2 , the initial object of $\mathbf{EAlg}$. $\square$

The copy-discard structure of $\mathbf{BKer}$ is given in the evident way. For a Boolean algebra A , the ‘codiscard’ map is the unique map

$$2 \rightarrow \mathbb{I} \otimes A$$

and the ‘cocompy’ map is

$$A \otimes A \rightarrow \mathbb{I} \otimes A$$

given the composing the codiagonal in $\mathbf{BAlg}$ with the unit of $\mathbb{I}$ in $\mathbf{EAlg}$.

$$A \otimes A \cong A +_{\mathbf{BAlg}} A \rightarrow A \cong 2 \otimes A \rightarrow \mathbb{I} \otimes A$$

Lemma 11. *The deterministic subcategory of $\mathbf{BKer}$ is precisely $\mathbf{BAlg}^{\text{op}} \hookrightarrow \mathbf{BKer}$.*

Proof. Let $f : B \rightarrow \mathbb{I} \otimes A$ represent a deterministic map $A \rightarrow B$ of $\mathbf{BKer}$. We want to show that, for all $b \in B$, $f(b)$ has the form $1 \otimes a$ for some $a \in A$. Write $f(b)$ in the form $\alpha_1 \otimes a_1 + \dots + \alpha_k \otimes a_k$ with the a_i mutually orthogonal and $0 \leq \alpha_i \leq 1$. Copyability implies

$$\alpha_1 \otimes a_1 + \dots + \alpha_k \otimes a_k = \alpha_1^2 \otimes a_1 + \dots + \alpha_k^2 \otimes a_k$$

whence $\alpha_i \in \{0, 1\}$, as required. $\square$

Causality From [1, Example 11.35], the Markov category of all probability kernels between measurable spaces is causal. Since $\mathbf{BKer}$ embeds faithfully into it with a morphism of Markov categories, $\mathbf{BKer}$ is also causal. Let's also argue directly.

Suppose we are given

$$f : B \rightarrow \mathbb{I} \otimes A \quad g : C \rightarrow \mathbb{I} \otimes B \quad h_i : D \rightarrow \mathbb{I} \otimes C$$

in $\mathbf{EAlg}$ for $i = 1, 2$ with $A, B, C, D \in \mathbf{BAlg}$ such that the two composites

$$(h_i \otimes 1_C) \circ \text{copy}_C \circ g \circ f : A \rightarrow D \otimes C$$

in $\mathbf{BKer}$ agree. We are required to show that the two composites

$$(((h_i \otimes 1_C) \circ \text{copy}_C \circ g) \otimes 1_B) \circ \text{copy}_B \circ f : A \rightarrow D \otimes C \otimes B$$

are equal. The latter are actually maps $D \otimes C \otimes B \rightarrow \mathbb{I} \otimes A$, hence it suffices to check equality on generators of the form $d \otimes c \otimes b$. By decomposing c , we can assume that c is ‘small’ enough that both $h_i(d) \in \mathbb{I} \otimes C$ are sums of the form $\sum_i \alpha_i x_i$ with, for each i , either $c \leq x_i$ or $c \wedge x_i = 0_C$. In terms of $\mathbf{M}(C)$, this says that both $h_i(d)$ are constant on the clopen subset of $S(C)$ corresponding to c . It follows that the maps $(h_i \otimes 1_C) \circ \text{copy}_C : C \rightarrow D \otimes C$ send, as functions $D \otimes C \rightarrow \mathbb{I} \times C$, the element $d \otimes c$ to $\xi_i \otimes c$ for some $\xi_i \in [0, 1]$. By hypothesis, either $g \circ f : A \rightarrow C$ sends, as a function $C \rightarrow \mathbb{I} \otimes A$, the element c to 0 or $\xi_1 = \xi_2$. In either case it is clear to see that the required equation holds.

Kolmogorov products Recall that $\mathbf{BAlg}$ is the closure under filtered colimits of the representable functors in $[\mathbf{Fin}, \mathbf{Set}]$, and that by construction the Day tensor preserves colimits in each argument. It follows that the inclusion $\mathbf{BAlg} \hookrightarrow \mathbf{EAlg}$ preserves filtered colimits, as does the (strict) monoidal functor $\mathbf{BAlg} \rightarrow \mathbf{BKer}^{\text{op}}$, and hence that $\mathbf{BKer}$ has all infinite tensor products of all small families, given by the coproduct in $\mathbf{BAlg}$ of the underlying objects. Obviously, the marginals are deterministic, i.e. $\mathbf{BAlg}$ has all Kolmogorov products.

References

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