RSA encryption

Why do we need to keep things secret?

- Historically, secret messages were used in wars and battles
 - For example, the Enigma cipher in WW2
- Computers are nowadays used to store lots of information about us
 - The recent "lost CDs in the post" contained data that was not kept secret
- Also need to make sure that online purchases are secure

Terminology

- Cryptography: study of creating secret codes
- Plaintext: the original message
- Ciphertext: the message after encryption
- Encryption: the process of creating a ciphertext
- Decryption: reversing encryption
- Cipher: name often given to an encryption

- The RSA cipher uses two keys and uses lots of Maths to make it work
- The key used for encryption is called the public key and the key used for decryption is called the private key (which is kept secret)
- To make it hard to break, one of the keys needs to be around 600 digits (2048 bits)
- Named after Rivest, Shamir and Adleman

RSA recipe

- Pick two large primes p and q
- Work out $n = p \times q$
- Work out $k = (p-1) \times (q-1)$
- Pick a public key *e* which does not have any factors (except 1) in common with *k*
- Find a private key d that $d \times e = 1 \pmod{k}$
- To encrypt *m*, work out *m^e* (mod *n*)
- To decrypt *c*, work out *c*^d (mod *n*)

Primes and Factoring

- For RSA, we have to pick two prime numbers, let's say p and q.
- Remember a prime number has exactly two factors (so 2 and 3 are prime, but not 1 or 4)
- We have to work out the product n of the two primes so n = p × q
- To find p and q from n is hard we would have to check all the (prime) numbers up to p or q. This is what makes RSA secure.

RSA encryption uses modular arithmetic.

- We write 10 (mod 6) to mean "give me the remainder when 10 is divided by 6"
- So 10 (mod 6) = 4 (because $10 = 6 \times 1 + 4$) 75 (mod 9) = 3 (75 = 9 \times 8 + 3) 56 (mod 12) = 8 (56 = 12 \times 4 + 8) 89 (mod 10) = 9 (89 = 10 \times 8 + 9)

Keys

- We now pick a public key e
- We find a private key *d* such that:

$$d \times e = 1 \pmod{(p-1)(q-1)}$$

(There's a method called the extended Euclidean algorithm which can find this)

- To encrypt a number (message) M we do
 M^e (mod n)
- To decrypt, we do $C^d \pmod{n}$



- Pick p = 7, q = 13 so n = 7 × 13 = 91
- Pick e = 5 and then d = 29
 - Check 5 × 29 = 145 = (2 × 6 × 12) + 1
- To encrypt the number 10, we do

 $10^5 \pmod{91} = 10 \times 10 \times 10 \times 10 \times 10 \pmod{91} = 82$

To decrypt we do

 $82^{29} \pmod{91} = 10$

Working out powers

We needed to work out $82^{29} \pmod{91}$ – how can we do this quickly and easily?

We split up the power (29) as follows: $82^{29} = 82^1 \times 82^4 \times 82^8 \times 82^{16}$ and we work out these powers mod 91

Note that 82²⁹ is quite big with 56 digits.



$82^{2} = 6724 = 81 \pmod{91}$ $82^{4} = 82^{2} \times 82^{2} = 81 \times 81 = 9 \pmod{91}$ $82^{8} = 82^{4} \times 82^{4} = 9 \times 9 = 81 \pmod{91}$ $82^{16} = 82^{8} \times 82^{8} = 81 \times 81 = 9 \pmod{91}$

So, $82^{29} = 82 \times 9 \times 81 \times 9 \pmod{91}$ = 738 × 729 (mod 91) = 10 × 1 = 10 (mod 91)

Using Letters

So far we've seen how to encrypt numbers. To encrypt text we give letter a value A=1, B=2,...,Z=26 and Space=27and then combine the numbers using Base 28. Just as $234 = (2 \times 10^2) + (3 \times 10^1) + (4 \times 10^0)$ Then "BCD" = $(2 \times 28^2) + (3 \times 28^1) + (4 \times 28^0)$ = 1568 + 84 + 4 = 1656



- To make RSA fast we need to be able to find powers quickly – there are methods to help us do this
- The strength on RSA relies on the fact that it is generally hard to factor a very large number
- By using two keys, we can keep the private key secret
- There are many other "public key" methods