

Problem Sheet 1

Instructions: The problem sheets are designed to increase your understanding of the material taught in the lectures, as well as to prepare you for the final exam. You should attempt to solve the problems on your own after reading the lecture notes and other posted material, where applicable. Problems marked with an asterisk are optional. Once you have given sufficient thought to a problem, if you are stuck, you are encouraged to discuss with others in the course and with the lecturer during office hours. You are *not permitted* to search for solutions online.

1 Learning Hyper-rectangles

The concept class of hyper-rectangles over \mathbb{R}^n is defined as follows:

$$C_n = \{[a_1, b_1] \times \cdots \times [a_n, b_n] \mid a_i, b_i \in \mathbb{R}, a_i < b_i\}$$

Generalise the algorithm discussed in class (for rectangles in \mathbb{R}^2) and show that it *efficiently* PAC learns the class of hyper-rectangles. Give bounds on the number of samples required to guarantee that the error is at most ϵ with probability at least $1 - \delta$.

Note: You may assume that the distribution D over \mathbb{R}^n can be expressed using a continuous density function that is defined over all of \mathbb{R}^n . (*Optional:* As an extra challenge, argue why the algorithm still works even when such an assumption regarding the distribution does not hold.)

2 PAC Learning : Confidence Parameter

Say that an algorithm L *perhaps learns* a concept class C using hypothesis class H , if for every n , for every concept $c \in C_n$, for every distribution D over X_n and for every $0 < \epsilon < 1/2$, L given access to $\text{EX}(c, D)$ and inputs ϵ and $\text{size}(c)$, runs in time polynomial in n , $\text{size}(c)$ and $1/\epsilon$, and outputs a polynomially evaluable hypothesis $h \in H_n$, that with probability at least $3/4$ satisfies $\text{err}(h) \leq \epsilon$. In other words, we've set $\delta = 1/4$ in the definition of *efficient* PAC learning. Show that if C is "perhaps learnable" using H , then C is also *efficiently* PAC learnable using H .

Hint: You will have to use the Chernoff-Hoeffding bound. Your proof should not rely on $\delta < 1/2$, in particular it should work for $\delta = 3/4$, even though we have required $0 < \delta < 1/2$ in the definition of PAC learning.

3 Hardness of Learning Boolean Threshold Functions

We will consider the question of learning boolean threshold functions. Let $X_n = \{0, 1\}^n$ and for $w \in \{0, 1\}^n$ and $k \in \mathbb{N}$, $f_{w,k} : X_n \rightarrow \{0, 1\}$ is a boolean threshold function defined as follows:

$$f_{w,k}(x) = \begin{cases} 1 & \text{if } \sum_{i=1}^n w_i \cdot x_i \geq k \\ 0 & \text{otherwise} \end{cases}$$

Let $\text{TH}_n = \{f_{w,k} \mid w \in \{0, 1\}^n, 0 \leq k \leq n\}$ and $\text{TH} = \bigcup_{n \geq 1} \text{TH}_n$. Show that unless $\text{RP} = \text{NP}$, there is no efficient PAC-learning algorithm for learning $\overline{\text{TH}}$, if the output hypothesis is also required to be in TH , *i.e.*, a *proper* PAC-learning algorithm.

Hint: You should reduce from Zero-One Integer Programming (ZIP) which is known to be NP-complete. An instance of ZIP consists of an $s \times n$ matrix A with entries in $\{0, 1\}$, a vector $\mathbf{b} \in \{0, 1\}^s$ and a pair (\mathbf{c}, B) (the objective) with $\mathbf{c} \in \{0, 1\}^n$ and $B \in \mathbb{Z}$. The decision problem is to determine whether there exists an assignment for the n variables z_1, \dots, z_n , each variable taking a value in the set $\{0, 1\}$, such that for each $1 \leq i \leq s$, $\sum_{j=1}^n A_{ij}z_j \leq b_i$ and $\sum_{j=1}^n c_jz_j \geq B$.

4 Learning Parity Functions

Let $X_n = \{0, 1\}^n$ be the instance space. A parity function over X_n is defined by some subset $S \subseteq \{1, \dots, n\}$, and takes the value 1 if and odd number of the input literals in the set $\{x_i \mid i \in S\}$ are 1 and 0 otherwise. For example, if $S = \{1, 3, 4\}$, then the function $f = x_1 \oplus x_3 \oplus x_4$ computes the parity on the subset $\{x_1, x_3, x_4\}$. Note that any such parity function can be represented by a bit string of length n , by indicating which indices are part of S . Let **PARITIES** denote the class of all parity functions. Show that the class **PARITIES** is efficiently PAC-learnable by describing an algorithm, analysing its running time and proving its correctness.

Hint: The parity operation can be viewed as addition modulo 2.

5 Output Hypothesis as a Turing Machine

Recall that in the definition of *efficient* PAC-learning, we require that the hypothesis output by the learning algorithm be evaluable in polynomial time. Suppose we relax this restriction, and let H be the class of all Turing machines (not necessarily polynomial time)—so the output of the learning algorithm can be any program. Let C_n be the class of all boolean circuits of size at most $p(n)$ for some fixed polynomial p and having n boolean inputs. Show that $C = \bigcup_{n \geq 1} C_n$ is “efficiently” PAC-learnable using H (under this modified definition). Argue that this solution shows that the relaxed definition trivialises the model of learning.