Introduction to Functional Programming using Haskell

Errata

April 7, 1999

Chapter 1

page 3, line 1 Replace \( \text{square} 14198724 \) by \( \text{square} 14197824 \).

page 9, line 14 Replace “hit the the interrupt key” by “hit the interrupt key”.

page 22, line 8 Insufficient space between names in \( \text{square square} \). [This unfortunates space compression occurs in various places throughout the text.]

page 23, line 14 “Look again the previous . . .” should read “Look again at the previous . . .”

page 25, line 22 Replace by “The link between the two is the requirement that the im-
plementation satisfies . . .”

Chapter 2

page 33, lines 1 - 8 The text is confused. There are two solutions:

Either replace it by “We can declare \( \text{Bool} \) to be an instance of \( \text{Ord} \) by writing

\[
\text{instance \textit{Ord}} \ \text{Bool} \ \text{where}
\]

\[
\text{False} \leq \text{False} \quad = \quad \text{True}
\]

\[
\text{False} \leq \text{True} \quad = \quad \text{True}
\]

\[
\text{True} \leq \text{False} \quad = \quad \text{False}
\]

\[
\text{True} \leq \text{True} \quad = \quad \text{True}
\]

The alternative definition, namely \( x \leq y = \text{not} \ \text{\( x \lor y \)} \), doesn’t quite work in the way expected (see Exercise 2.1.2). ”

Or replace lines 7 and 8 by “As an alternative definition we can write \( x < y = \text{not} \ \text{\( x \land y \)} \).

page 33, line 20 “Note that the two occurrences of . . .”

page 34 Delete Exercise 2.1.2.
In Exercise 2.1.9 replace the last sentence by “Show that these properties hold for the definition of (++) on the well-defined values of Bool.

“a different entity from the decimal number 7;...”

“but does not depend ...”

Systematically interchange the names toEnum and fromEnum in the whole of Section 2.3.

Interchange toEnum and fromEnum.

The type of plus should be

\[ \text{plus} :: (\alpha \to \beta, \gamma \to \delta) \to \text{Either} \ \alpha \gamma \to \text{Either} \ \beta \delta \]

Note that Haskell uses the name either rather than case.

Chapter 3

Replace \( h(\text{foldn} \ h \ b \ \text{Succ} \ n) \) by \( h(\text{foldn} \ h \ b \ n) \).

Replace by

\[ \text{Rat} \ x \ y = \text{Rat} \ u \ v = (x \times v) :: (y \times u) \]

Among possible representations we can choose one in which \(-5 \leq z < 5\) and \(\text{abs} \ y\) is as small as possible.

“since programs that avoid case analyses are clearer and simpler than those that do not, ...”

The definition of \( \text{done} \) should read

\[ \text{done} \ (m, n) = (m + 1 :: n) \]

In the definition of \( y_3 \) a division by 2 is missing:

\[ y_3 = (1.4167 + 2/1.4167)/2 = 1.4142157 \]
Chapter 4

page 103, line 14 First line in definition of \textit{init} should read: \textit{init} \([x] = []\).

page 112, line 10 “This equation is valid provided \(p\) and \(q\) are strict functions.”

page 114, line 4 Definition of \textit{pyth} should read

\[
\text{pyth} (x, y, z) = (x \times x + y \times y :: z \times z)
\]

page 116, line 14 Last line in definition of \textit{zip} should read:

\[
\text{zip} (x : xs) (y : ys) = (x, y) : \text{zip} \, xs \, ys
\]

page 116, line 18 “the scalar product of two vectors \(x\) and \(y\) of size \(n\) is defined by . . .”

page 121, line 22 The function \textit{zip} can be defined as an instance of \textit{foldr}: we have \textit{zip} = \textit{foldr} \(f\) \(e\) where

\[
e \, ys = []
\]

\[
f \, x \, g \, [] = []
\]

\[
f \, x \, g \, (y : ys) = (x, y) : g \, ys
\]

page 124, line 12 “the first is clearer, while the second is more efficient.”

page 124, line -6 The type of \textit{scanl} should read:

\[
\text{scanl} :: (\beta \to \alpha \to \beta) \to \beta \to [\alpha] \to [\beta]
\]

page 125, line -1 Replace \(a\) by \(e\) in equations involving \textit{scanr}.

page 126, lines 9,10 Replace \(a\) by \(e\) in equations involving \textit{scanr}.

page 127 Exercise 4.5.9 should read “What list does \textit{scanl} (\(/\)) \, 1 \, [1..n] produce?”

page 125 In Exercise 4.5.11 the type of \textit{convert} should be \textit{Liste} \(\alpha \to [\alpha]\).

page 129, line -4 “Both sides simplify to \(x \oplus y\).”

page 129, line -1 Replace \(z; xs\) by \(z : xs\).

page 130, line -7 Missing period at end of paragraph.

page 137 In Exercise 4.6.10 the law should read

\[
\text{foldl}1 (\oplus) \cdot \text{scanl} (\otimes) \, e = \text{fst} \cdot \text{foldl} (\odot) \, (e, e)
\]
Chapter 5

page 146, line 9 Replace \textit{assign} \textit{xs} by \textit{assign}.

page 146, line 12 The definition of \textit{mktriple} should read

\begin{align*}
\text{mktriple}(xn, xm) xr &= (xn, xm, xr)
\end{align*}

page 148, line 18 The last line of the definition of \textit{sortby} should read:

\begin{align*}
\text{sortby } f (x : y : xs) &= \text{mergeby } f (\text{cross (sortby } f, \text{sortby } f) (\text{divide}(x : y : xs)))
\end{align*}

page 163, line 12 The three = signs on the right-hand side of the definition of \textit{leap} should be \texttt{==} signs.

page 165, lines 12,15 The types of \textit{stackWith} and \textit{spreadWith} should be

\begin{align*}
\text{stackWith} &:: \text{Height} \rightarrow [\text{Picture}] \rightarrow \text{Picture} \\
\text{spreadWith} &:: \text{Width} \rightarrow [\text{Picture}] \rightarrow \text{Picture}
\end{align*}

page 167, line 16 Replace \textit{entries} \((d, s)\) with just \textit{entries}.

page 167, line 23 In the definition of \textit{dnames} the conversion to type \textit{Picture} is omitted, so prefix the right-hand side with \textit{row}.

Chapter 6

page 185, lines -6 – -1 Replace definition of \textit{fork} by

\begin{align*}
\textit{fork} &:: \text{Atree } \alpha \rightarrow \text{Atree } \alpha \rightarrow \text{Atree } \alpha \\
\textit{fork } xt yt &= \text{Fork}(\text{lsize } xt) xt yt \\
\text{lsize} &:: \text{Atree } \alpha \rightarrow \text{Int} \\
\text{lsize} (\text{Leaf } x) &= 1 \\
\text{lsize} (\text{Fork } n xt yt) &= n + \text{lsize} yt
\end{align*}

page 186, line 6 The two occurrences of \textit{mkBtree} should be replaced by \textit{mkAtree}.

page 187 In Exercise 6.1.3 the definition of \textit{subtrees} should read

\begin{align*}
\text{subtrees} &:: \text{Btree } \alpha \rightarrow [\text{Btree } \alpha] \\
\text{subtrees} (\text{Leaf } x) &= [\text{Leaf } x] \\
\text{subtrees} (\text{Fork } xt yt) &= [\text{Fork } xt yt] + \text{subtrees } xt + \text{subtrees } yt
\end{align*}

page 188, line 9 The type of \textit{member} should read

\begin{align*}
\text{member} &:: \text{Ord } \alpha \Rightarrow \alpha \rightarrow \text{Stree } \alpha \rightarrow \text{Bool}
\end{align*}
page 188, line 15 Space compressed in \textit{member} \textit{x} \textit{xt}.

page 188, line 17 The type of \textit{height} should read

\[
\text{height :: Ord } \alpha \Rightarrow \textit{Stree} \alpha \rightarrow \textit{Int}
\]

page 190, line -10 The identity should read:

\[
x \mathbin{+} y = x \mathbin{+} [\text{head } y] \mathbin{+} \text{tail } y
\]

page 191, line -1 Exercise 6.2.4 should read “Prove that \textit{inordered} (\textit{insert} \textit{x} \textit{xt}) = \textit{True} for all finite binary search trees \textit{xt}.

page 193, line -11 The type of \textit{heapify} should read

\[
\text{heapify :: Ord } \alpha \Rightarrow \textit{Htree} \alpha \rightarrow \textit{Htree} \alpha
\]

page 193, line -6 The type of \textit{sift} should read

\[
\text{sift :: Ord } \alpha \Rightarrow \alpha \rightarrow \textit{Htree} \alpha \rightarrow \textit{Htree} \alpha \rightarrow \textit{Htree} \alpha
\]

page 196, lines 13 – 17 Omit the local definition of \textit{maxlist}, and insert the following sentence in the text: “Recall that \textit{maxlist} = \textit{foldl1} (\textit{max}).

page 197, line -5 The left-hand expression should read: \( f (g(x, y), z, h(t)) \).

page 200, line -6 The type of \textit{combine} should read: \textit{combine :: [[\alpha]]} \rightarrow [[\alpha]].

page 201, line 8 Same correction as above.

page 205, line -3 The type declaration of \textit{CodeTable} should read:

\[
\text{type CodeTable} = [(\text{Char}, [\text{Bit}], \text{Int})]
\]

page 206, line 19 Replace local definition by \textbf{where} \((ys, zs) = \textit{span} (z :: x) \textit{xs}\).

Chapter 7

page 231, line 8 The type of \textit{dfcat} should read:

\[
\text{dfcat :: [Rose } \alpha] \rightarrow [\alpha] \rightarrow [\alpha]
\]

page 236, line 8 The second line in the definition of \textit{fills} should read:

\[
\text{fills} (w : ws) = [us : vss \mid (us, vs) \leftarrow \textit{splits} (w : ws); vss \leftarrow \textit{fills} vs]
\]

page 236, line -1 Space compression in \textit{fill} \textit{vs}.
Chapter 8

page 256, line 16 “The second implementation therefore has a different efficiency from the first,…”

page 256, line 21 The right-hand side of the second axiom for back should be

\[ \text{join } x (\text{back } (\text{join } y x)) \]

page 272, lines 13,15 Remove closing parenthesis from right-hand expressions.

page 278, line 12 The type of fork should be

\[ \text{fork} \quad : \quad \alpha \rightarrow \text{Htree } \alpha \rightarrow \text{Htree } \alpha \rightarrow \text{Htree } \alpha \]

page 279, line 1 The type of delMin should be

\[ \text{delMin} \quad : \quad \text{Ord } \alpha \Rightarrow \text{Htree } \alpha \rightarrow \text{Htree } \alpha \]

page 279, line 3 The type of union should be

\[ \text{union} \quad : \quad \text{Ord } \alpha \Rightarrow \text{Htree } \alpha \rightarrow \text{Htree } \alpha \rightarrow \text{Htree } \alpha \]

page 280, line 6 The type of mkBag should be

\[ \text{mkBag} \quad : \quad \text{Ord } \alpha \Rightarrow [\alpha] \rightarrow \text{Htree } \alpha \]

page 280, line 8 The type of mkTwo should be

\[ \text{mkTwo} \quad : \quad \text{Ord } \alpha \Rightarrow \text{Int } \rightarrow [\alpha] \rightarrow (\text{Htree } \alpha, [\alpha]) \]

page 284, line 9 Replace right-hand side of otherwise branch by

\[ \text{Fork } n \ x t \ (\text{update } y t \ (k \ - \ m) \ x) \]

page 288, line 4 Replace nullys by null ys.

page 288, line 22 It should be pointed out that the definition of abstr is exactly the same as in the implementation of Section 8.1 since

\[ \text{ys } + \ \text{reverse } \text{xs} \quad = \quad \text{reverse } (\text{xs } + \ \text{reverse } \text{ys}) \]

page 289, line 14 Replace reverseys by reverse ys.

page 290, line -1 Replace last line by

\[ (3, 0, \text{rot } (\text{rot } [\ ] [1] [\]) [3, 2] [\], [\]) \]

page 291, line 2 Replace by

\[ (3, 3, \text{rot } (\text{rot } [\ ] [1] [\]) [3, 2] [\], [6, 5, 4]) \]

page 291, line 4 Replace by

\[ (7, 0, \text{rot } (\text{rot } (\text{rot } [\ ] [1] [\]) [3, 2] [\]) [7, 6, 5, 4] [\], [\]) \]

page 291, line 6 Replace by

\[ (7, 7, \text{rot } (\text{rot } (\text{rot } [\ ] [1] [\]) [3, 2] [\]) [7, 6, 5, 4] [\], [14, 13..8]) \]
Chapter 9

page 296, line 17 “the computer determines the first four elements . . .”

Chapter 10

page 330, line 13 Should add “where $C$ may be empty”.

page 342 In Exercise 10.2.2 the type definition should read

\[
\text{newtype } \text{Count } \alpha \ = \ \text{CNT} (\alpha, \text{Counter})
\]

Chapter 11

page 365, line -10 The definition should read

\[
p \text{ orelse } q \ = \ \text{MkP } f \\
\text{ where } f s \ = \ \text{if } \text{null } ps \ \text{then } \text{apply } q s \ \text{else } ps \\
\text{ where } ps = \text{apply } p s
\]

page 365, line -6 The operator \text{orelse} does not satisfy the distributive law of \text{plus}.

page 368, line -1 Replace $\triangleright$ by $\triangleright\triangleright$.

page 373 In Exercise 11.4.1 add “for deterministic parsers $p$ and $q$”.

Chapter 12

page 384, line -11 Type of \text{notuple} should read

\[
\text{notuple :: Parser } [\text{Expr}]
\]

Appendix

page 411, line -1 Replace by \text{head} $(x : xs) = x$. 
