

# A Library Approach to Information Flow Security in Haskell

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# Motivating Example

- Passwords

Dictionary attacks, offline attacks, ...



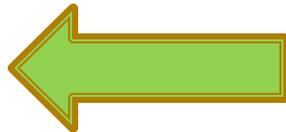
*/etc/passwd*



Universal Access



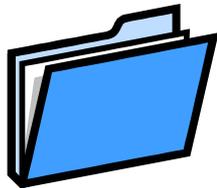
*/etc/shadow*



Root Access

# Motivating Example

- Passwords



*/etc/passwd*



*/etc/shadow*

Dictionary

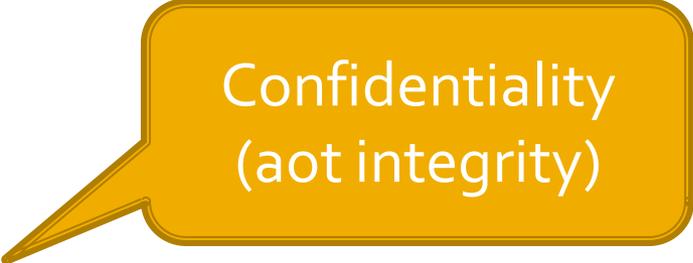
attack  
attack

Linux Shadow Password HOWTO:  
Adding shadow support to a C  
program

“Adding shadow support to a  
program is actually fairly  
straightforward. The only problem is  
that the program must be run by root  
in order for the the program to be  
able to **access** the */etc/shadow* file.”

# The Problem

- For the sake of
  - Intruders
  - People we let in (plug-ins)
  - Ourselves
- We want to restrict
  - Access to data
  - Where does data go?
  - Where is it used?

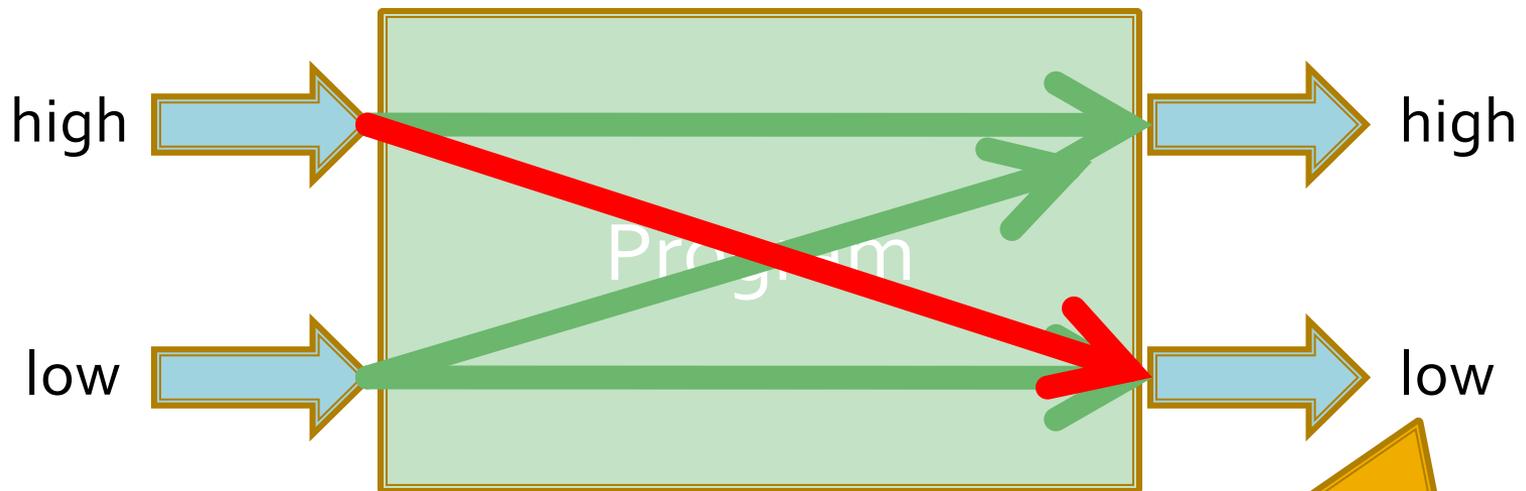


Confidentiality  
(not integrity)



“Information-flow  
security”

# The Model



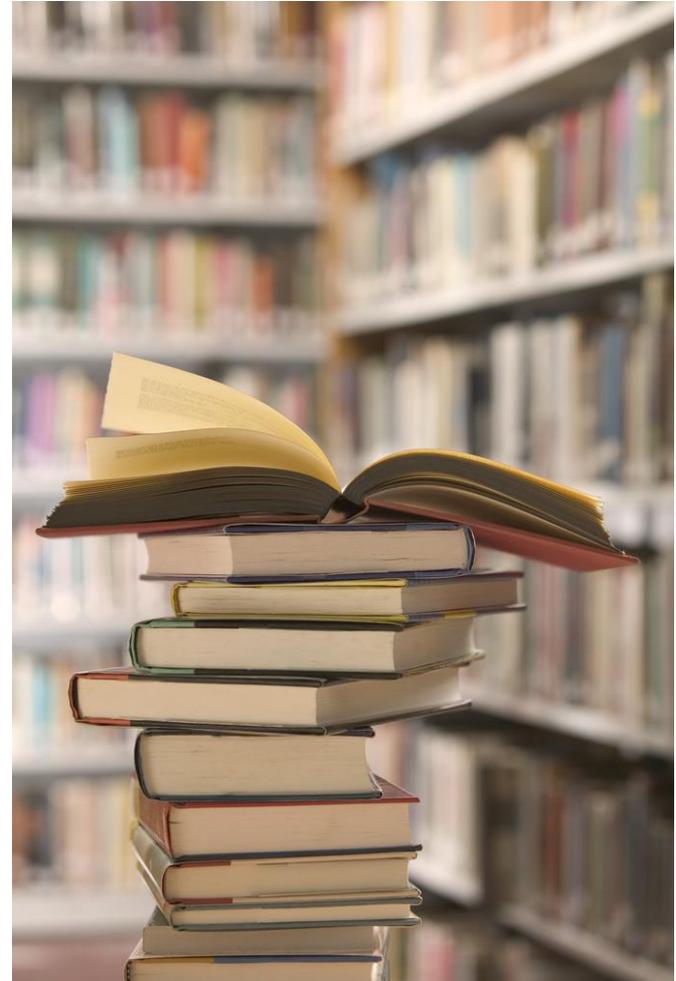
Non-interference: *Varying high inputs should not affect low inputs*

# "Attacker"

- Attacker
  - Not trusted
  - Intruder
  - Programmer
  - Yourself
- Everyone (including the attacker) can observe low security outputs

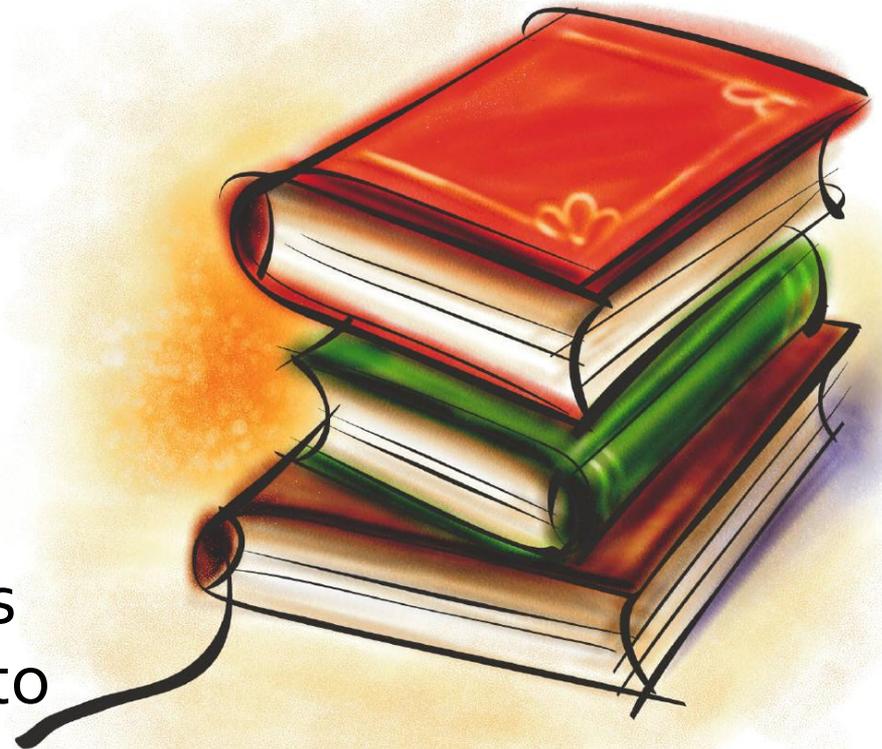
# Information-Flow Security

- Study for ~30 years
- Active research field
- Compilers
  - JIF (Java) 2001
    - Cornell University
  - FlowCaml (ML) 2002
    - INRIA (not actively developed)
- Impact on practice
  - Limited!



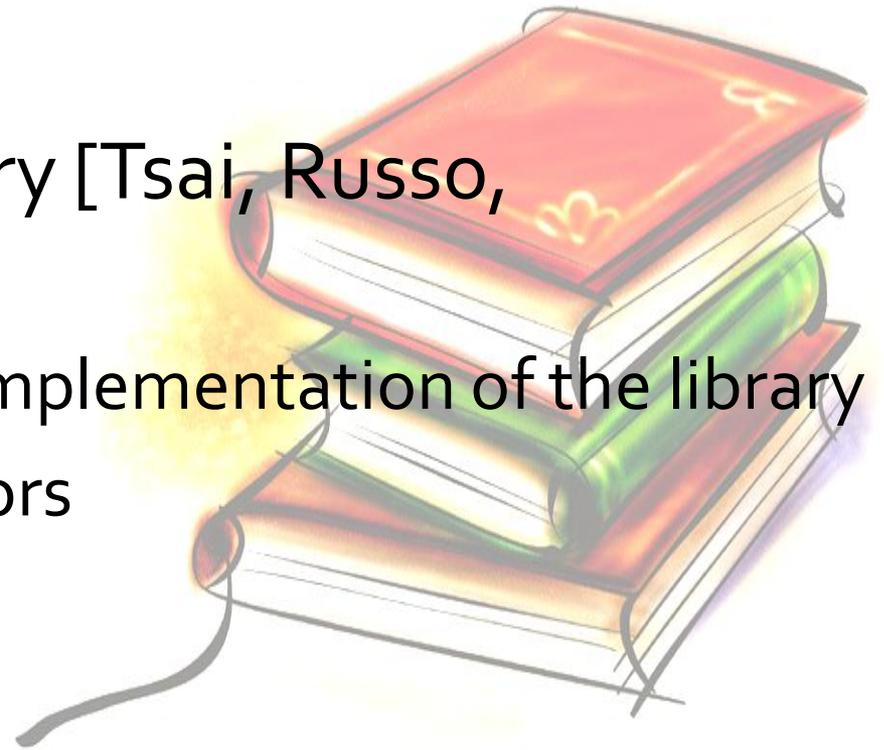
# Encoding IF in Haskell

- Possible to guarantee IF by a library
  - [Zdancewic & Li, 06]
  - Haskell
  - Arrows
- No need to write a compiler from scratch
- DSEL approach: Quick experimenting with ideas
- No restriction on the PL to use due to security



# Encoding IF in Haskell

- Limitations
  - No side effects
- Extension to the library [Tsai, Russo, Hughes'07]
  - Major changes in the implementation of the library
  - New arrows combinators
  - Lack of arrow notation
- Why arrows?
  - Zdancewic and Li mention that **monads** are **not suitable** for the design of the library

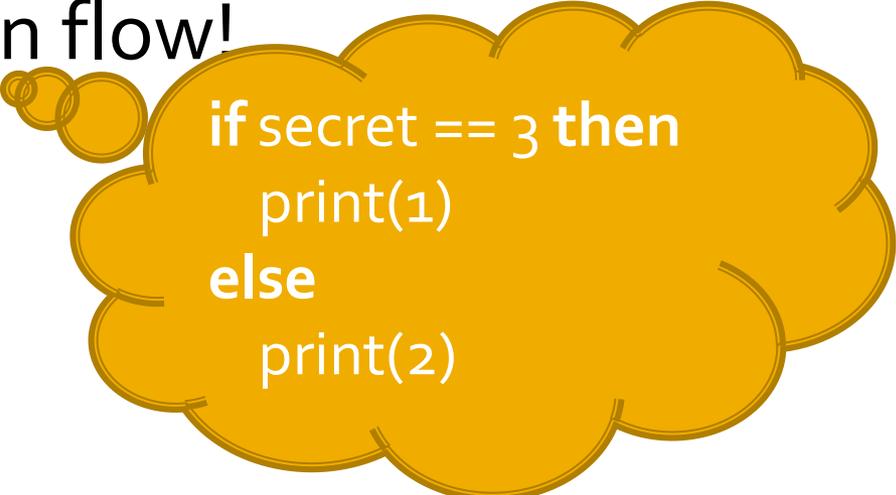


# Our Approach

- Light-weight
- Library-based
- Monad-based (not arrows)
- Restrict capabilities
  - Abstract types
  - Use of the module system
- Practical (?)

# Why Haskell?

- Pure language
  - No side effects
  - (Controlled side effects)
- Strong type system
  - Cannot "cheat"
- No implicit information flow!
  - Only explicit



```
if secret == 3 then  
  print(1)  
else  
  print(2)
```

# Example

```
f :: (Int {-secret-}, Char)
  -> (Int {-secret-}, Char)
```

```
f (n, c) = (n + 1, chr (ord c + 1))
```

YES

YES

```
f (n, c) = (n + ord c, 'a')
```

NO

```
f (n, c) = (n + ord c, chr n)
```

NO

```
f (n, c) | n > 0      = (42, c)
          | otherwise = (1, chr (ord c + 1))
```

# Simple Security API

```
type Sec a -- abstract  
sec  :: a -> Sec a  
open :: Sec a -> Key -> a
```



strict!

```
data Key = TheKey -- hidden
```

```
instance Functor Sec  
instance Monad Sec
```

# Non-Interference?

type A

type B

type C

type D

$f :: (\text{Sec } A, B) \rightarrow (\text{Sec } C, D)$

$f(a1, b) = (c, d) \Rightarrow f(a2, b) = (c', d)$

# Multiple Security Levels

```
type Sec s a -- abstract
```

```
sec  :: a -> Sec s a
```

```
open :: Sec s a -> s -> a
```

# Security Lattice

```
data H = H -- abstract
```

```
data L = L -- public
```

```
class Less low high where
```

```
  up :: Sec low a -> Sec high a
```

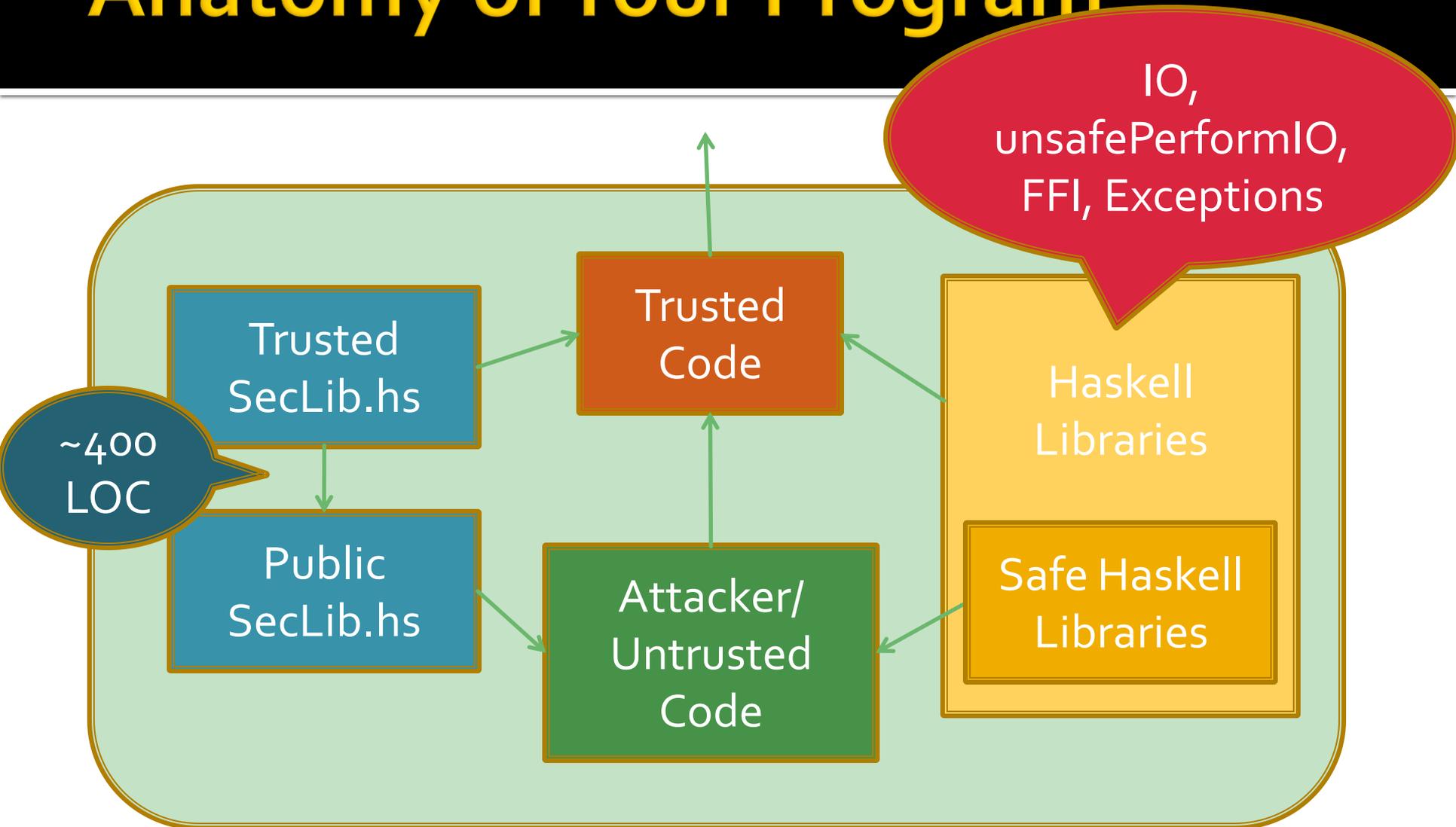
```
instance Less L H
```

```
instance Less L L
```

```
instance Less H H
```

Sec L a $\sim$ a
------------------

# Anatomy of Your Program



# How About IO?

- IO features
  - File IO
  - stdin/stdout
  - State references
  - Channels
  - ...
- This talk: Only File IO

# First Try

```
type File s -- abstract
```

```
readFileSec :: File s -> IO (Sec s String)
```

```
writeFileSec :: File s -> Sec s String -> IO ()
```

# High Control / Implicit Flow

- "Depending on a high value, write to file1 or file2"
- Leads to result types
  - IO (Sec H a)
  - Sec H (IO (Sec H a))
  - IO (Sec H (IO (Sec H a)))
  - ...
- Need a new type for "secure IO"

# SecIO

- \* Read from level  $s$  or lower
- \* Write to level  $s$  or higher
- \* Produce a value at level  $s$

```
type SecIO s a -- abstract
```

```
peek      :: Sec s a -> SecIO s a
```

```
readFileSec :: File s -> SecIO s String
```

```
writeFileSec :: File s -> String -> SecIO s ()
```

```
run      :: SecIO s a -> IO (Sec s a)
```

Side effects escape  
"Sec s"!

# SecIO

```
example :: Sec H Int -> SecIO s ()
example secret =
  do x <- peek secret
     if x == 42
       then writeFileSec file1 "foo"
       else writeFileSec file2 "bar"
```

# Anatomy of Your Program

shadow :: File H  
passwd :: File L

main = ... Untrusted.main shadow passwd ...

main :: File H -> File L -> IO (Sec H Answer)  
main shadow passwd = run (...)



# File Capabilities

```
type File m s -- abstract
```

```
data R
```

```
data W
```

```
readFileSec :: File R s -> SecIO s String
```

```
writeFileSec :: File W s -> String -> SecIO s ()
```

---

```
passwd :: File R L
```

```
shadow :: File R H
```

```
database :: File m H -- polymorphic
```

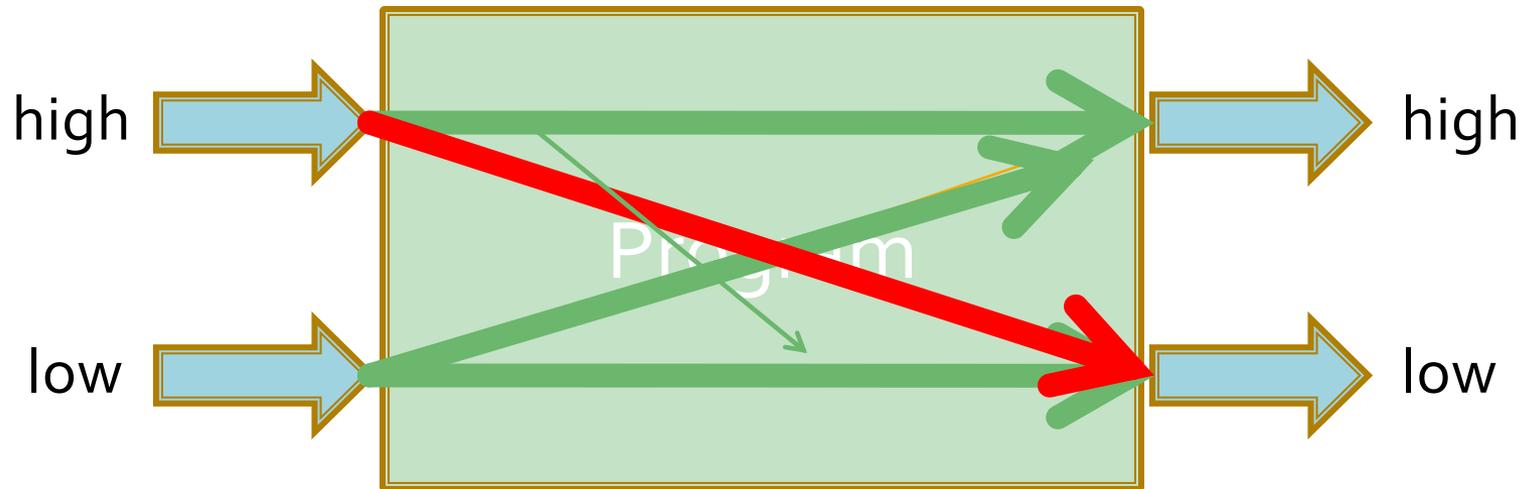
# Information Flow in Practice

- Login program
  - Get password from user input
  - Check if it is correct (compare with shadow)
  - Act accordingly
- It is necessary to leak information that depends on secrets!
  - `cypher inp == pwd`
  - Not non-interferent

# Declassification

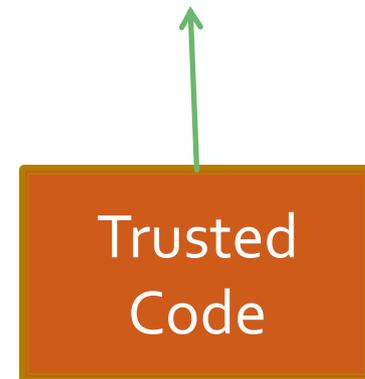
- Dimensions and principles of declassification [Sabelfeld and Sands, 06]
  - **What** information can be leaked?
  - **When** can information be leaked?
  - **Where** in the program is it safe to leak information?
  - **Who** can leak information?
- How to be certain that our programs leak what they are supposed to leak?

# The Adapted Model



# Introducing Declassification in the Library

- Our library should be able to handle different kind of *declassification policies*
- Policies are programs!
  - Trusted users of the library implement them
  - Controlled at run-time
- A module defines combinators for different declassification policies (**what, when, who**)



# Escape Hatches

- Declassification is performed by functions
  - Terminology: **escape hatches** [Sabelfeld and Myers, 2004]
- In our library:

```
type Hatch sH sL a b = Sec sH a -> Sec sL b
```

```
hatch :: (a -> b) -> Hatch sH sL a b -- hidden
```

- Example: checking password

monomorphic

```
check :: Hatch H L (String, Passwd) Bool  
check = hatch (\(inp, pwd) -> cypher inp == pwd)
```

# Escape Hatches

- We want to restrict capabilities of escape hatches

```
type Hatch sH sL a b =  
  Sec sH a -> IO (Maybe (Sec sL b))
```



internal state



may fail

# Declassification Combinators

```
-- restricting "what" (how often)
nTimes :: Int -> Hatch sH sL a b ->
                                         IO (Hatch sH sL a b)

-- example
check =
  nTimes 3
    (hatch (\(inp,pwd) -> cypher inp == pwd))
```

# Implementation

```
-- restricting "what" (how often)
nTimes :: Int -> Hatch sH sL a b ->
                                             IO (Hatch sH sL a b)
nTimes n hatch =
  do ref <- newIORef n
     return (\x ->
       do k <- readIORef ref
          if k >= 0
            then do writeIORef ref (k-1)
                   hatch x
            else do return Nothing)
```

# Declassification Combinators

```
-- restricting "when" (flow locks)
data Open    = Open (IO ())    -- hidden
data Close   = Close (IO ())   -- hidden

when :: Hatch sH sL a b ->
      IO (Hatch sH sL a b, Open, Close)
```

# Declassification Combinators

```
-- restricting "who" (flow locks)
data Authority s = Auth Open Close -- hidden

who :: Hatch sH sL a b ->
      IO (Hatch sH sL a b, Authority sH)

-- for use by attacker
certify :: s -> Authority s -> IO a -> IO a
```

# Declassification Combinators

- Powerful
- Expressive
- Theory of declassification is in its infancy
  - One dimension only
  - Weak results
- In practice, we want to combine things
  - Pragmatic

# Correctness Proof

- "Sec" -- obvious and trivial
- All other things
  - SecIO
  - Files
  - References
  - ...
- On top of Sec: also obvious
- With slight modification: small proof



To do



To do

# QuickChecking Correctness

- Modelled library + language as a Haskell datatype
  - Evaluate function
- Written a random generator
  - Respecting types
- Expressed non-interference as a QuickCheck property
  - Counter-examples for unsound versions of the library

# Conclusions

- Light-weight library (~400 LOC)
- Practical
  - Simple (Monads)
  - Features: files, stdio/stdout, references
  - Declassification
  - Examples: login system, bidding, banking system prototype,...
- Limitations
  - Timing leaks
  - Static security lattice