

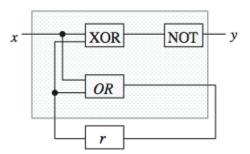


Formal Synthesis of Control Strategies for Dynamical Systems

Calin Belta

Tegan Family Distinguished Professor Mechanical Engineering, Systems Engineering, Electrical and Computer Engineering

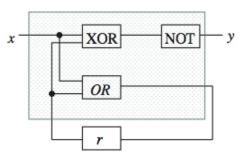
Boston University





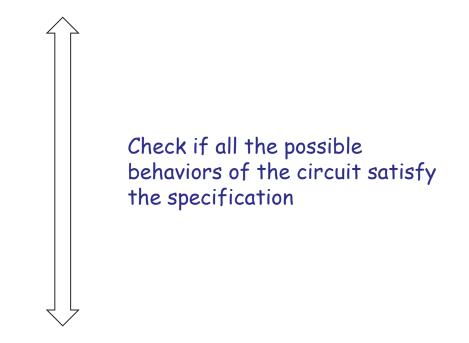
C. Bayer and J-P Katoen, Principles of Model Checking, MIT Press, 2008

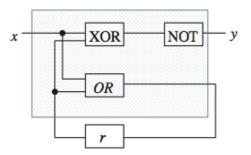
Specification: "If x is set infinitely often, then y is set infinitely often."



Process

Specification: "If x is set infinitely often, then y is set infinitely often."

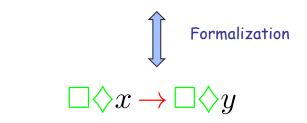




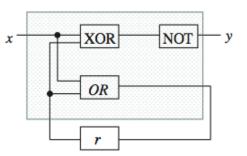
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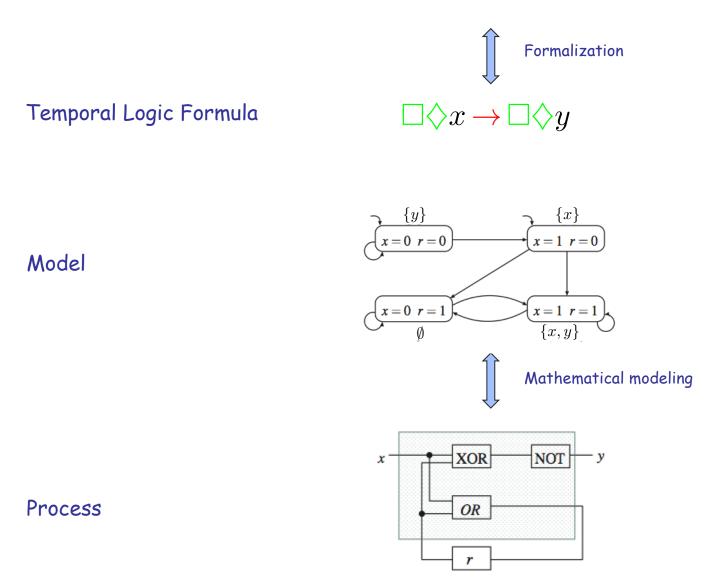


Temporal Logic Formula

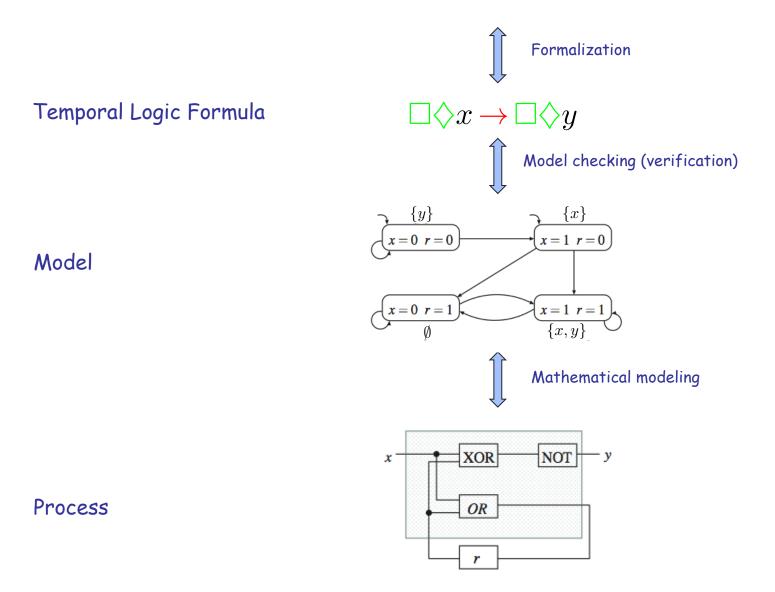


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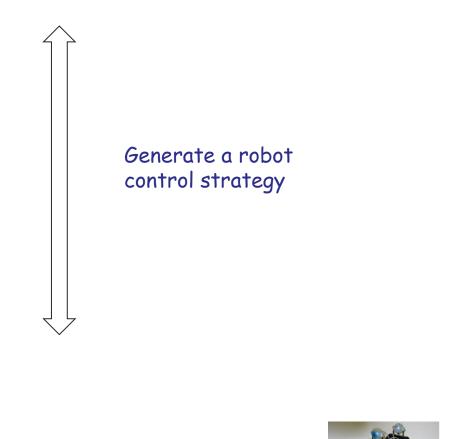


Specification: "drive from A to B."

• A

Process

Specification: "drive from A to B."

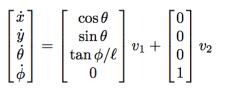


Process

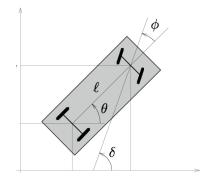
• A

Specification: "drive from A to B."

Model



οB



Mathematical modeling



Process

S. Sastry - Nonlinear Systems: analysis, stability, and control, Springer, 1999

Specification: "drive from A to B."

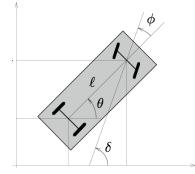
Formalization

Stabilization Problem: "make B an asymptotically stable equilibrium"

Model

$$\begin{bmatrix} \dot{x} \\ \dot{y} \\ \dot{\theta} \\ \dot{\phi} \end{bmatrix} = \begin{bmatrix} \cos \theta \\ \sin \theta \\ \tan \phi/\ell \\ 0 \end{bmatrix} v_1 + \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix} v_2$$

οB



Mathematical modeling



Process

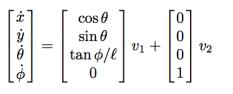
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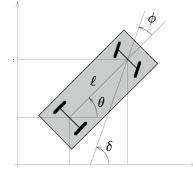
Formalization

Control

Stabilization Problem: "make B an asymptotically stable equilibrium"

Model



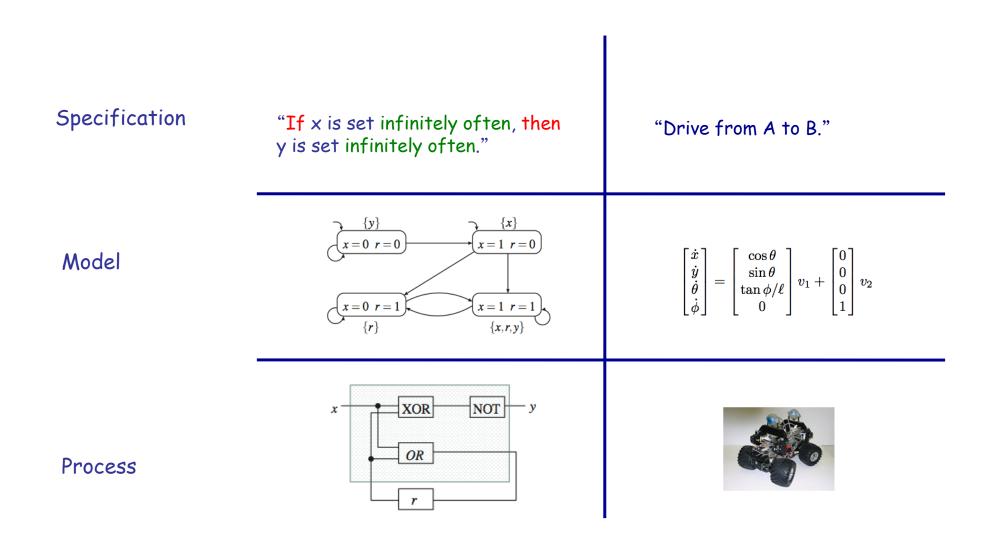


Mathematical modeling

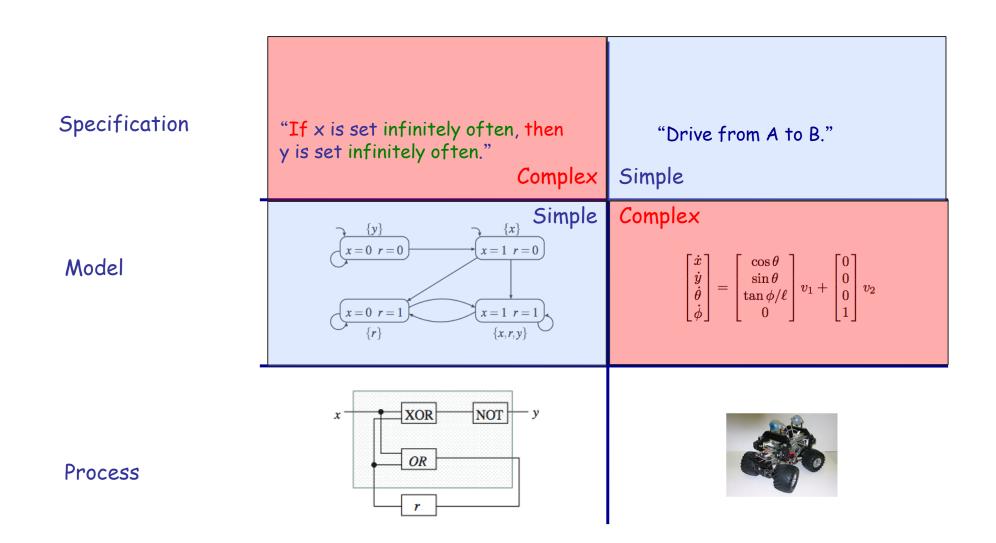


Process

Formal methods vs. dynamics

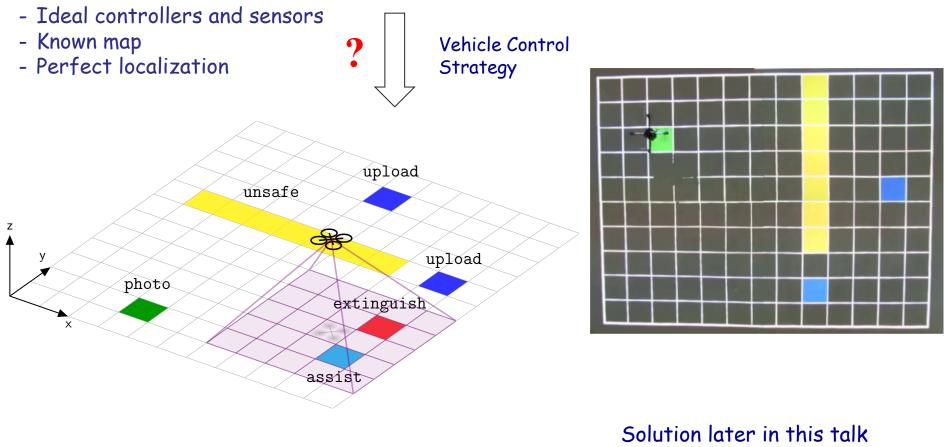


Formal methods vs. dynamics



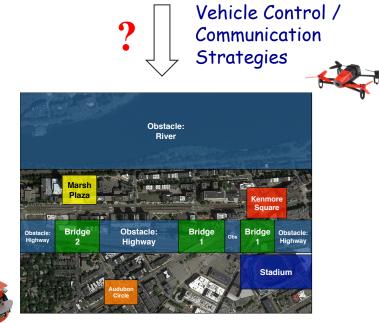
Need for formal methods in dynamical systems

Spec: **Off-line**: "Keep taking photos and upload current photo before taking another photo. **On-line**: Unsafe regions should always be avoided. If fires are detected, then they should be extinguished. If survivors are detected, then they should be provided medical assistance. If both fires and survivors are detected locally, priority should be given to the survivors."



Need for formal methods in dynamical systems

Spec: Maximize the probability of satisfying: "Always avoid all obstacles and Visit Marsh Plaza, Kenmore Square, Fenway Park, and Audubon Circle infinitely often and Bridge 2 should only be used for Northbound travel and Bridges 1 should only be used for Southbound travel. Uncertainty should always be below 0.9 m² and when crossing bridges it should be below 0.6 m²."

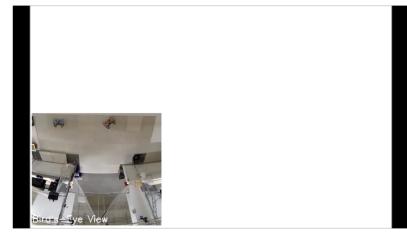




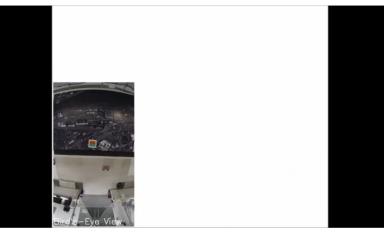
- Noisy controllers and sensors
- Unknown map
- Probabilistic localization

Example later in the talk

E. Cristofalo, et.al., *ISER* 2016. C. I. Vasile et.al., CDC 2016



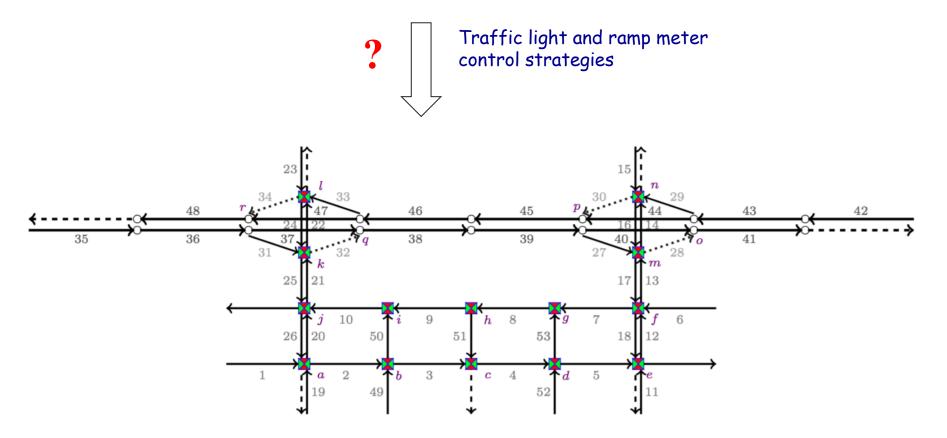
Map unknown environment



Localization and control

Need for formal methods in dynamical systems

- **Spec:** always the network is not congested
 - each queue at a junction will be actuated at least once every two minutes
 - whenever the number of vehicles on a link exceeds 40, within 3 min it should decreases below 20



Example later in the talk

Coogan, et.al.,, ACC 2015, IEEE TCNS 2016 Sadradini, Belta, ACC 2016, CDC 2016 Coogan, Arcak, Belta, ACC 2016, CSM 2017

Outline

TL verification and control for finite systems

Conservative TL control for small & simple dynamical systems

Conservative TL control for large & complex dynamical systems

Less conservative optimal TL control for small & simple dynamical systems

Less conservative TL control for large & (possibly) complex dynamical systems

Less conservative optimal TL control for large & simple dynamical systems



TL = Temporal Logic

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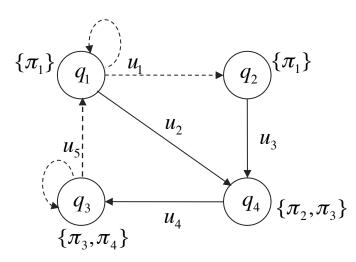
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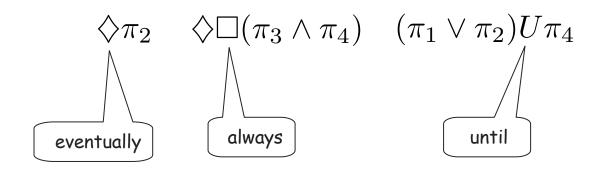
Finite system

(Fully-observable) nondeterministic (non-probabilistic) labeled transition systems with finitely many states, actions (controls), and observations (properties)



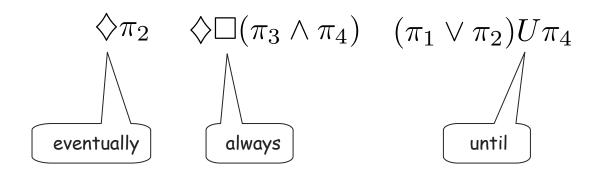
Linear Temporal Logic (LTL)

Syntax



Linear Temporal Logic (LTL)

Syntax

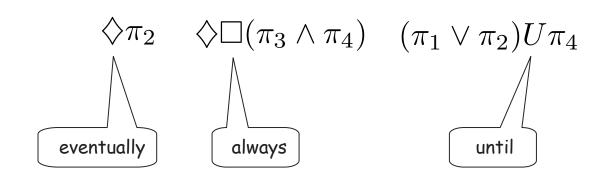


Semantics

Word:
$$\{\pi_1\}\{\pi_2,\pi_3\}\{\pi_3,\pi_4\}\{\pi_3,\pi_4\}\cdots$$

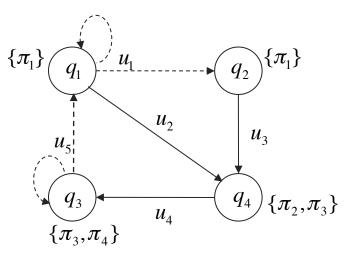
Linear Temporal Logic (LTL)

Syntax



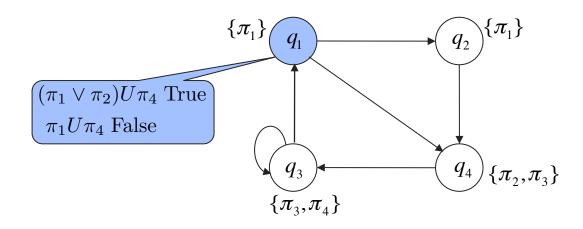
Semantics

Run (trajectory): $q_1, q_4, q_3, q_3, \dots$ Word: $\{\pi_1\}\{\pi_2, \pi_3\}\{\pi_3, \pi_4\}\{\pi_3, \pi_4\}\cdots$



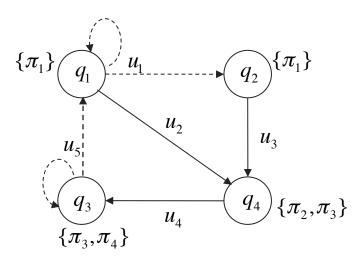
LTL verification (model checking)

Given a transition system and an LTL formula over its set of propositions, check if the language (i.e., all possible words) of the transition system starting from all initial states satisfies the formula.



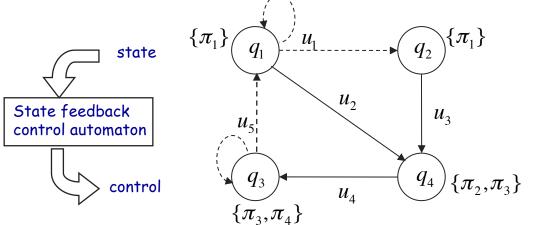
LTL control (synthesis)

Given a transition system and an LTL formula over its set of propositions, find a set of initial states and a control strategy for all initial states such that the produced language of the transition system satisfies the formula.



Did not receive much attention until recently!

LTL control (synthesis)





Particular cases:

- LTL without "eventually always": Buchi game
- LTL without "always" (syntactically co-safe LTL): the automaton is an FSA

Extensions

Optimal Temporal Logic Control for Finite Deterministic Systems Optimal Temporal Logic Control for Finite MDPs Temporal Logic Control for POMDPs Temporal Logic Control and Learning

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TL = Temporal Logic

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Conservative TL control for **large & complex** dynamical systems

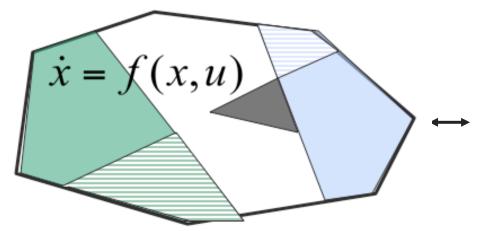
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Less conservative TL control for large & (possibly) complex dynamical systems

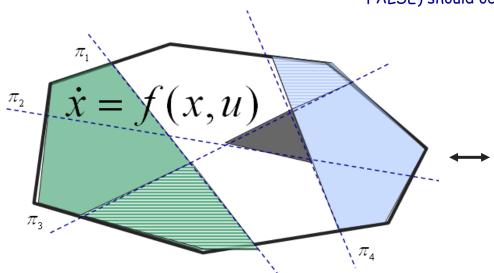
Less conservative optimal TL control for large & simple dynamical systems



TL = Temporal Logic

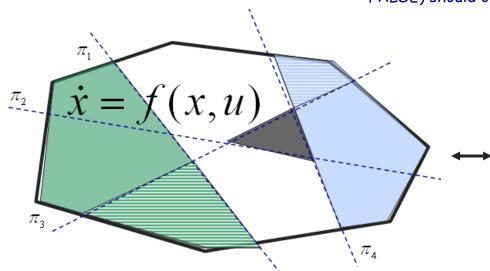


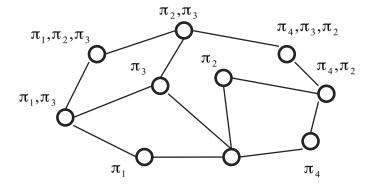
"(pi2 = TRUE and pi4 = FALSE and pi3 = FALSE) should never happen. Then pi4 = TRUE and then pi1 = TRUE should happen. After that, (pi3 = TRUE and pi4 = TRUE) and then (pi1 = TRUE and pi3 = FALSE) should occur infinitely often."



 $\Box \neg (\pi_2 \land \neg \pi_4 \land \neg \pi_3)) \land \\ \Diamond (\pi_4 \land \Diamond (\pi_1 \land \Diamond ((\pi_3 \land \pi_4) \land \Diamond (\pi_1 \land \neg \pi_3))))))$

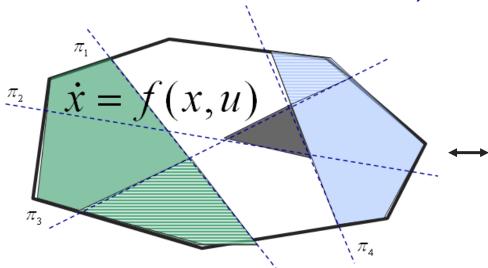
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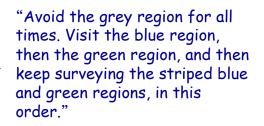




 $\Box \neg (\pi_2 \land \neg \pi_4 \land \neg \pi_3)) \land \\ \Diamond (\pi_4 \land \Diamond (\pi_1 \land \Diamond (\pi_1 \land (\pi_3 \land \pi_4) \land \Diamond (\pi_1 \land \neg \pi_3)))))$

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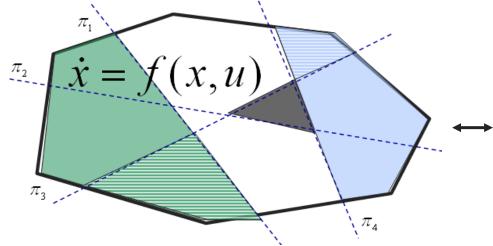


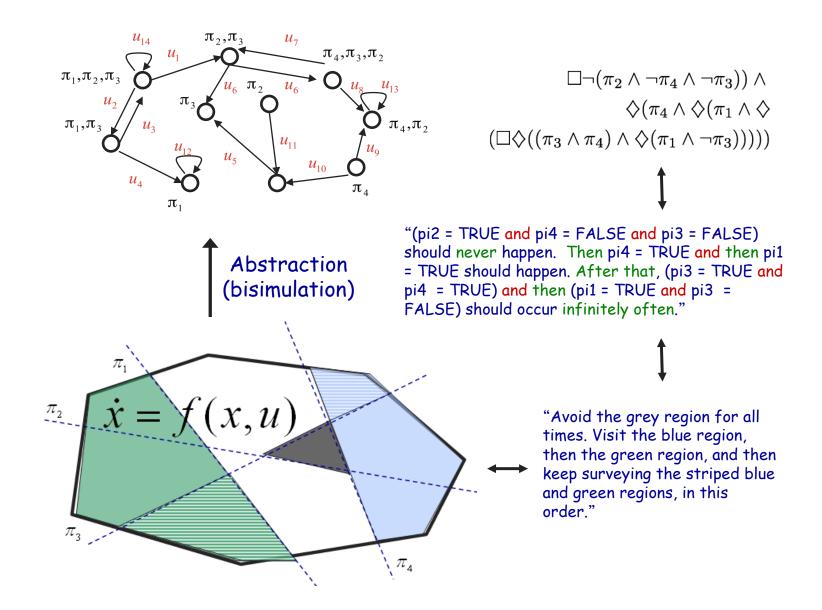


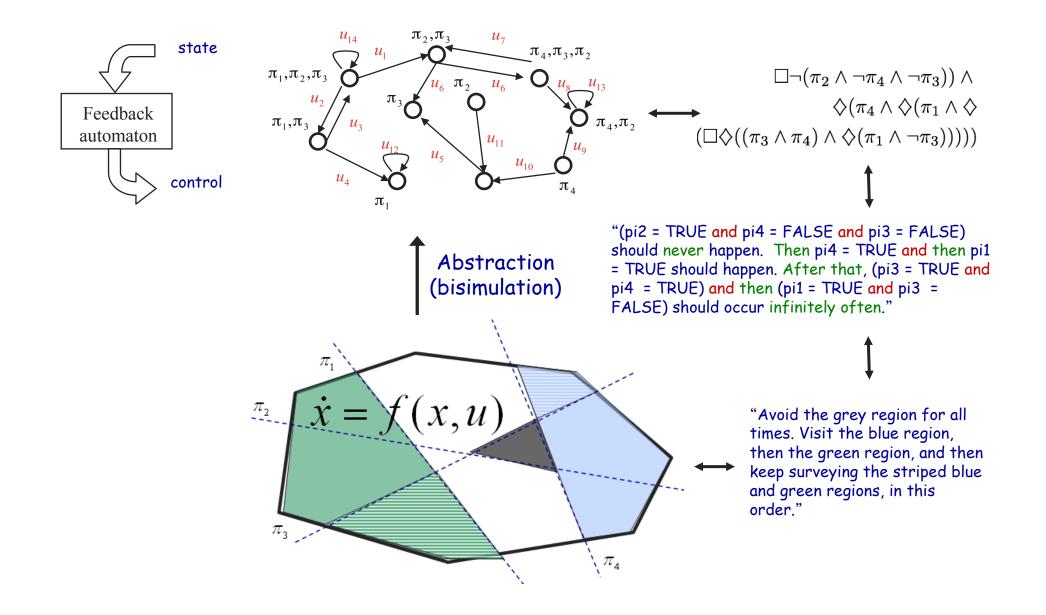
 π_{1},π_{2},π_{3} π_{1},π_{3},π_{2} π_{1},π_{3} π_{1},π_{3} π_{1},π_{2} π_{1},π_{2} π_{1},π_{2} π_{1},π_{2} π_{1},π_{2} π_{1},π_{2} π_{2} π_{1},π_{2} π_{1},π_{2} π_{2} π_{1},π_{2} π_{2} π_{1},π_{2} π_{2} π_{1},π_{2} π_{2}

Assume that in each region we can check for the existence of / construct feedback controllers driving all states in finite time to a subset of facets (including the empty set - controller making the region an invariant) $\Box \neg (\pi_2 \land \neg \pi_4 \land \neg \pi_3)) \land \\ \diamondsuit (\pi_4 \land \diamondsuit (\pi_1 \land \diamondsuit (\pi_1 \land (\pi_3 \land \pi_4) \land \diamondsuit (\pi_1 \land \neg \pi_3))))) \\ \uparrow$

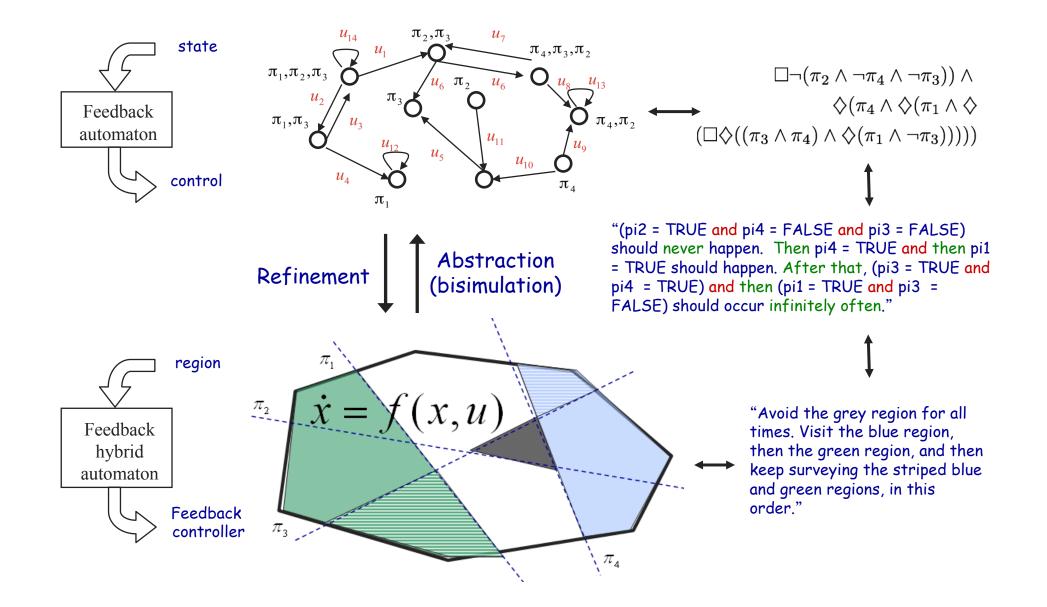
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Conservative TL control for small & simple dynamical systems



Conservative TL control for small & simple dynamical systems

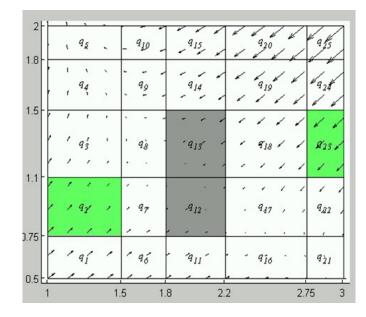
Dynamics and partitions allowing for easy construction of bisimilar abstractions

Library of controllers for polytopes $u \in U \subset \Re^m$ polyhedral U $\dot{x} = Ax + b + Bu$ $x \in \Re^n$ Control-to-facet Stay-inside Control-to-set-of-facets Control-to-face Stay-inside $u \in U \subset \Re^m$ $g(x) = \sum_{i_1, \dots, i_N \in \{0,1\}} C_{i_1, \dots, i_N} x_1^{i_1} \dots x_n^{i_n}$ $\dot{x} = g(x) + Bu$ $x \in \Re^n$ Control-to-facet Stay-inside Control-to-set-of-facets

checking for existence of controllers amounts to checking the non-emptiness of polyhedral sets in U
if controllers exist, they can be constructed everywhere in the polytopes by using simple formulas

L.C.G.J.M. Habets and J. van Schuppen, Automatica 2005 M. Kloetzer, et.al, CDC 2006 C. Belta and L.C.G.J.M. Habets, IEEE TAC, 2006

Conservative TL control for small & simple dynamical systems



Multi-affine dynamics

$$\dot{x}_1 = 2 - x_1 x_2 + u_1$$

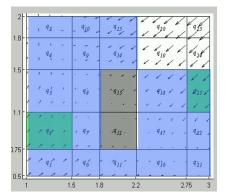
$$\dot{x}_2 = 1 + x_2 - x_1 x_2 + u_2,$$

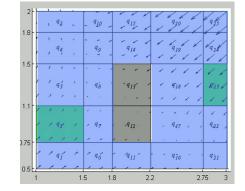
$$x \in [1,3] \times [0.5,2],$$

$$u \in [-1.5, 1.5] \times [-1.5, 1.5]$$

"visit the green regions, in any order, while avoiding the grey regions"

$$\Diamond q_2 \land \Diamond q_{25} \land \Box \neg (q_{12} \lor q_{13})$$





Control to one facet Deterministic quotient

Control to sets of facets Non-deterministic quotient

Initial states from which control strategies exist

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TL verification and control for finite systems

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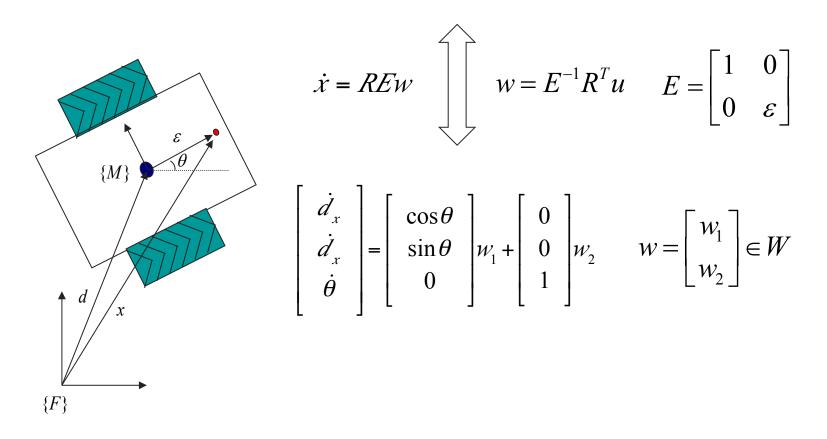
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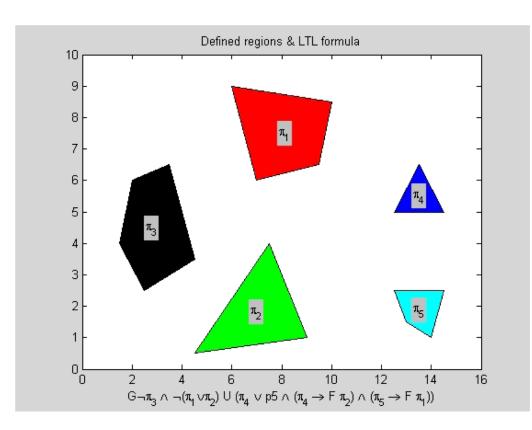
Mapping complex dynamics to simple dynamics: I/O linearization

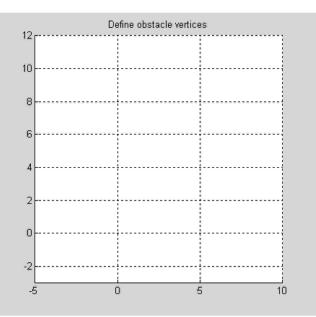
Fully actuated point $\dot{x} = u$ $u \in U$ U can be derived from W



J. Desai, J.P. Ostrowski, and V. Kumar. ICRA, 1998.

"Always avoid black. Avoid red and green until blue or cyan are reached. If blue is reached then eventually visit green. If cyan is reached then eventually visit red."



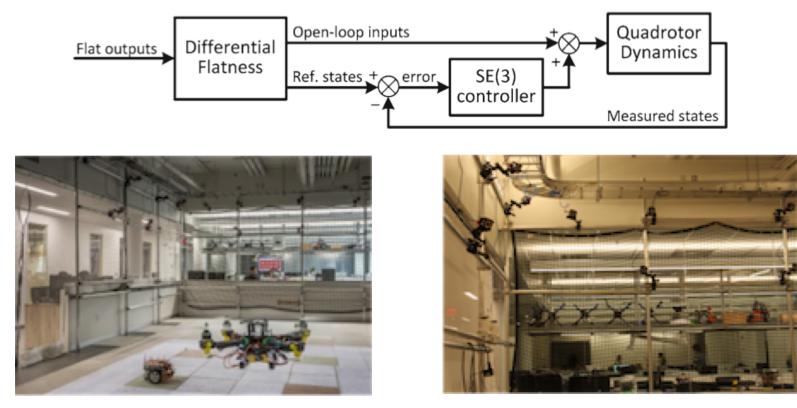


Mapping complex dynamics to simple dynamics: differential flatness

Quadrotor dynamics

- Nonlinear control system with 12 states (position, rotation, and their derivatives) with 4 inputs (total thrust force from rotors and 3 torques)
- Differentially flat with 4 flat outputs (position and yaw)
- Up to four derivatives of the flat output and necessary to compute the original state and input

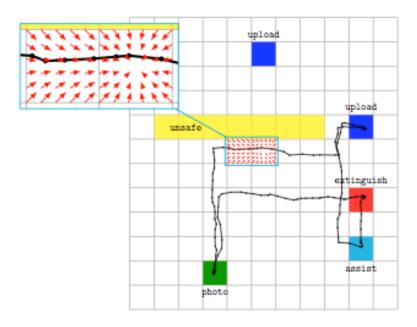
Mellinger and Kumar, 2011.; Hoffmann, Waslander, and Tomlin, 2008.; Leahy, Zhou, Vasile, Schwager, Belta, 2015



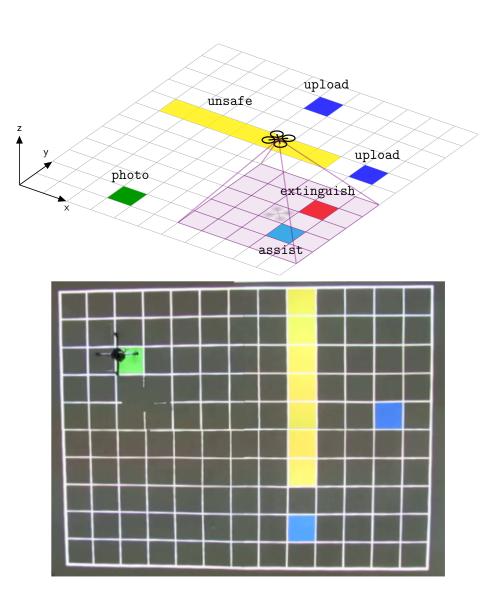
http://sites.bu.edu/robotics/

Persistent surveillance with global and local specs

Global spec: "Keep taking photos and upload current photo before taking another photo. Unsafe regions should always be avoided. Local spec: If fires are detected, then they should be extinguished. If survivors are detected, then they should be provided medical assistance. If both fires and survivors are detected locally, priority should be given to the survivors."

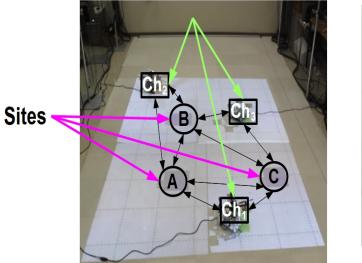


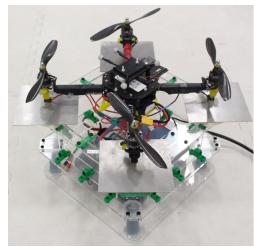
Ulusoy, Belta, RSS 2013, IJRR 2014



Persistent surveillance with deadlines and resource constraints

Charging Stations





Additional constraints:

- operation time
- charging time
- timed temporal specs

Mission Specification: Time Window Temporal Logic (TWTL)

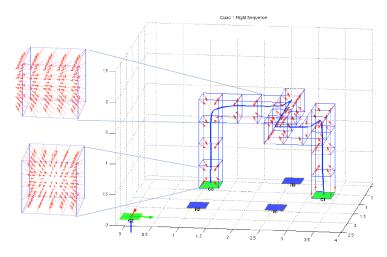
"Service site A for 2 time units within [0, 30] and site C for 3 time units within [0, 19]. In addition, within [0, 56], site B needs to be serviced for 2 time units followed by either A or C for 2 time units within [0, 10]."

$$\phi_{tw} = [H^2 A]^{[0,30]} \wedge [H^2 B [H^2 A \vee C]^{[0,10]}]^{[0,58]} \wedge [H^3 C]^{[0,19]}$$

Vasile, Aksaray, Belta, Theoretical Computer Science, 2017 Vasile and Belta, Robotics: Science and Systems (RSS) 2014

Persistent surveillance with deadlines and resource constraints

"Service site A for 2 time units within [0, 30] and site C for 3 time units within [0, 19]. In addition, within [0, 56], site B needs to be serviced for 2 time units followed by either A or C for 2 time units within [0, 10]."



Loop 1 (playback speed x1)

> Vasile, Aksaray, Belta, Theoretical Computer Science, 2017 (accepted) Vasile and Belta, Robotics: Science and Systems (RSS) 2014

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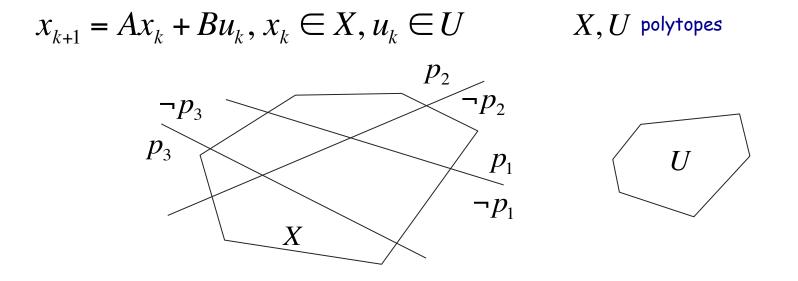
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Less conservative TL control for small and simple dynamics



Problem Formulation: Find a set of initial states and a state-feedback control strategy such that all trajectories of the closed loop system originating there satisfy an scLTL formula over a set of linear predicates

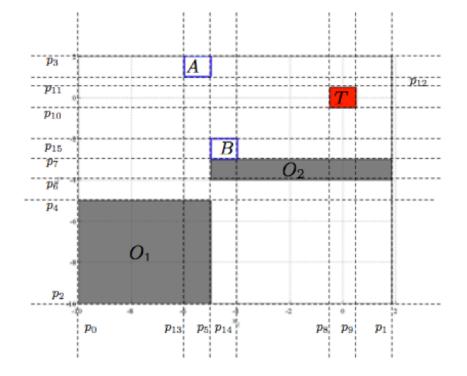
Language-guided Approach:

- Automaton-based partitioning and iterative refinement
- Polyhedral Lyapunov functions used to construct polytope-to-polytope controllers
- Solution is complete! (modulo linear partition and polyhedral Lyapunov functions)

Less conservative TL control for small and simple dynamics Example

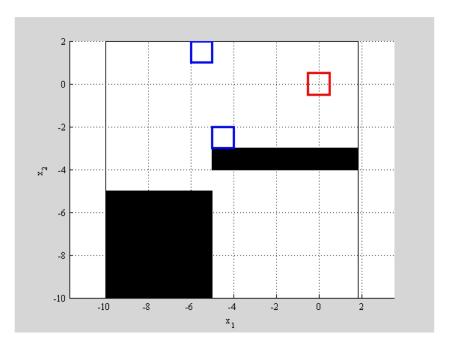
$$x_{k+1} = Ax_k + Bu_k, \quad x_k \in \mathbb{X}, \ u_k \in \mathbb{U}$$

"Visit region A or region B before reaching the target T while always avoiding the obstacles"



$$A = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \quad B = \begin{bmatrix} 0.5 \\ 1 \end{bmatrix}$$

 $\Phi_{2} = ((p_{0} \land p_{1} \land p_{2} \land p_{3} \land \neg (p_{4} \land p_{5}) \land \neg (\neg p_{5} \land \neg p_{6} \land p_{7})) \mathscr{U}$ $(\neg p_{8} \land p_{9} \land \neg p_{10} \land p_{11})) \land (\neg (\neg p_{8} \land p_{9} \land \neg p_{10} \land p_{11}) \mathscr{U} ((p_{5} \land \neg p_{12} \land \neg p_{13}) \lor (\neg p_{5} \land \neg p_{7} \land p_{14} \land p_{15})))$



E. Aydin Gol, M. Lazar, C. Belta., HSCC 2012, IEEE TAC 2014

Less conservative optimal TL control for small and simple dynamics

Optimal TL control

$$x_{k+1} = Ax_k + Bu_k, \quad x_k \in \mathbb{X}, \ u_k \in \mathbb{U}$$

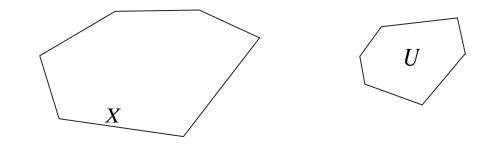
Initial state: x_0

Reference trajectories:

 $x_0^r, x_1^r \dots \\ u_0^r, u_1^r, \dots$

Observation horizon : N

Standard Model Predictive Control (MPC, Receding Horizon)



$$C(x_k, \mathbf{u}_k) = (x_{k+N} - x_{k+N}^r)^\top L_N(x_{k+N} - x_{k+N}^r) + \sum_{i=0}^{N-1} \{ (x_{k+i} - x_{k+i}^r)^\top L(x_{k+i} - x_{k+i}^r) + (u_{k+i} - u_{k+i}^r)^\top R(u_{k+i} - u_{k+i}^r) \},$$

Less conservative optimal TL control for small and simple dynamics



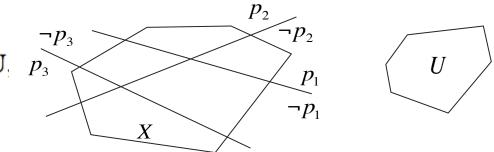
$$x_{k+1} = Ax_k + Bu_k, \quad x_k \in \mathbb{X}, \ u_k \in \mathbb{U}$$

Initial state: x_0 Reference trajectories: $x_0^r, x_1^r \dots$

 u_0^r, u_1^r, \dots

Observation horizon : ${\cal N}$

Standard Model Predictive Control (MPC, Receding Horizon)



$$C(x_k, \mathbf{u}_k) = (x_{k+N} - x_{k+N}^r)^\top L_N(x_{k+N} - x_{k+N}^r) + \sum_{i=0}^{N-1} \{ (x_{k+i} - x_{k+i}^r)^\top L(x_{k+i} - x_{k+i}^r) + (u_{k+i} - u_{k+i}^r)^\top R(u_{k+i} - u_{k+i}^r) \},\$$

Problem Formulation: Find an optimal state-feedback control strategy such that the trajectory originating at x_0 satisfies an scLTL formula over linear predicates p_i

Language-guided MPC Approach:

- Work on the refined automaton from the above TL control problem
- Enumerate paths of length given by the horizon and compute the costs.
- Terminal constraints ensuring the acceptance condition of the automaton: Lyapunovlike energy function
- Solve QP to find the optimal path

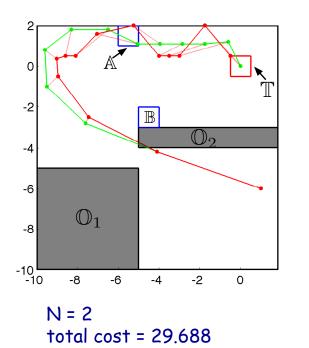
Less conservative optimal TL control for small and simple dynamics

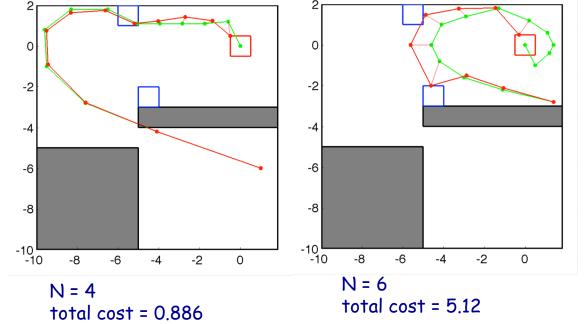
Example

$$x_{k+1} = Ax_k + Bu_k, \quad x_k \in \mathbb{X}, \ u_k \in \mathbb{U}$$

$$A = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \quad B = \begin{bmatrix} 0.5 \\ 1 \end{bmatrix}$$

- "Visit region A or region B before reaching the target while always avoiding the obstacles"
- Minimize the quadratic cost with $L=L_N=0.5I_2$, R=0.2





Reference trajectory violates the specification

Reference trajectory Controlled trajectory

TL verification and control for finite systems

Conservative TL control for small & simple dynamical systems

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Less conservative TL control for large & (possibly) complex dynamical systems

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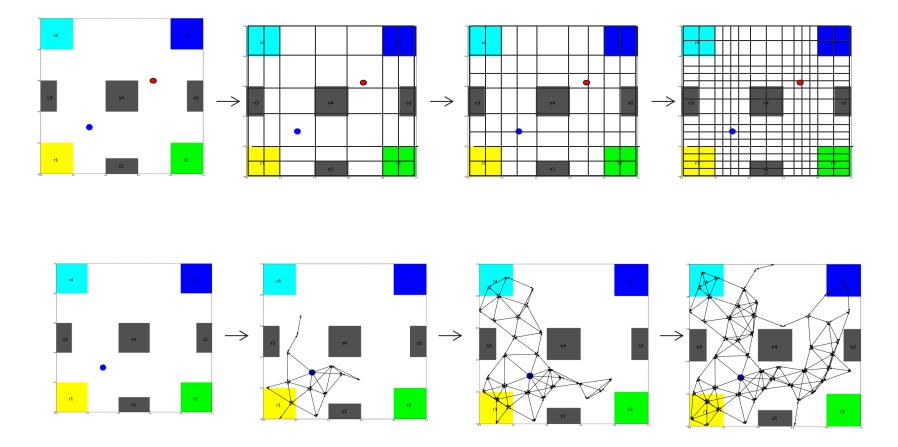
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Iterative Partition vs. Sampling in Motion Planning

Rapidly-exploring Random Trees (RRT) Rapidly-exploring Random Graphs (RRG) Steve LaValle, 1998

Karaman and Frazzoli, 2010



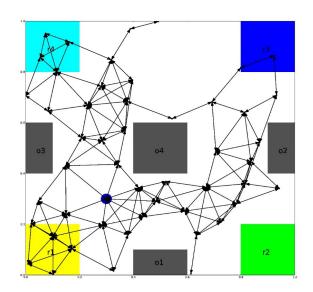
Only "go from A to B specs" can be handled with these techniques Can we extend sampling-based techniques to temporal logic specs?

Construct a transition system that contains a path satisfying the formula

- 1. LTL formula is translated to a Büchi automaton;
- 2. A transition system is incrementally constructed from the initial configuration using an RRG¹-based algorithm;
- 3. The product automaton is updated incrementally and used to check if there is a trajectory that satisfies the formula

Important Properties

- Probabilistically complete
- Scales incrementally (i.e., with the number of added samples at an iteration) - based on incremental Strongly Connected Component (SCC) algorithm²



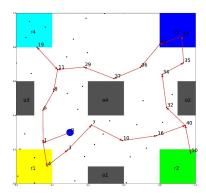
¹S. Karaman and E. Frazzoli. IJRR , 2011. ²Bernhard Haeupler, et al.. ACM Trans. Algorithms, 2012.

C. Vasile and C. Belta. IROS 2013

Case study 1: 2D configuration space, 20 runs Average execution time: 6.954 sec "Visit regions r1 r2 r3 and r4 infinitely many tim

"Visit regions *r1*, *r2*, *r3* and *r4* infinitely many times while avoiding regions *o1*, *o2*, *o3*, *o4* and *o5*"

 $\phi_1 = \mathbf{G}(\mathbf{F}r1 \wedge \mathbf{F}r2 \wedge \mathbf{F}r3 \wedge \mathbf{F}r4 \wedge \neg (o1 \vee o2 \vee o3 \vee o4))$



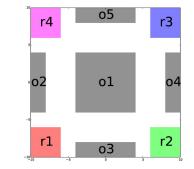
Case study 2: 10-dimensional configuration space, 20 runs Average execution time: 16.75 sec "Visit 3 regions r1, r2, r3 infinitely often while avoiding obstacle of $\phi_2 = \mathbf{G}(\mathbf{F}r1 \wedge \mathbf{F}r2 \wedge \mathbf{F}r3 \wedge \neg o1)$

Case study 3: 20-dimensional configuration space, 20 runs Average execution time: 7.45 minutes "Visit 2 regions (r1, r2) infinitely often" $\phi_3 = \mathbf{G}(\mathbf{F}r1 \wedge \mathbf{F}r2)$

Platform: Python2.7 on an iMac - 3.4 GHz Intel Core i7, 16GB of memory

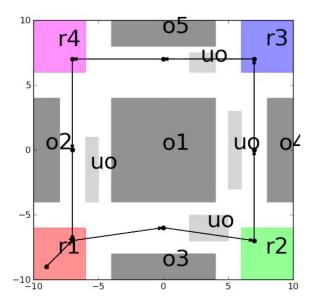
C. Vasile and C. Belta. IROS 2013

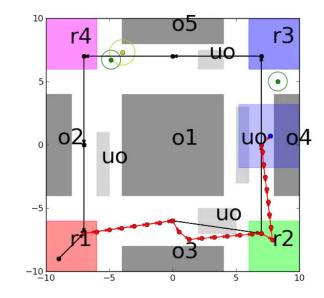
- Global mission specification: "visit regions r1, r2, r3 and r4 infinitely many times while avoiding regions o1, o2, o3, o4 and o5"
- Local mission specification: "Extinguish *fires* and provide medical assistance to *survivors*, with priority given to *survivors*, while avoiding *unsafe areas*"



Off-line part: generate a global transition system that contains a path satisfying the global spec

On-line (reactive) part: generate a local plan that does not violate the global spec





Fires and survivors are sensed locally. These service requests have given service radii.

C. Vasile and C. Belta, ICRA 2014

Spec: Maximize the probability of satisfying: "Always avoid all obstacles and Visit Marsh Plaza, Kenmore Square, Fenway Park, and Audubon Circle infinitely often and Bridge 2 should only be used for Northbound travel and Bridges 1 should only be used for Southbound travel. Uncertainty should always be below 0.9 m² and when crossing bridges it should be below 0.6 m²."

- Noisy controllers and sensors
- Unknown map
- No GPS



Approach:

- Generate a map of the unknown environment using purely vision and homography-based formation control with multiple quadrotors
- Label the map and define Gaussian Distribution Temporal Logic (GDTL) spec
- Synthesize control policy using GDTL Feedback Information RoadMaps (GDTL-FIRM)
- Simultaneously track and localize the ground robot with a single aerial vehicle using a homography - based pose estimation and position-based visual servoing control

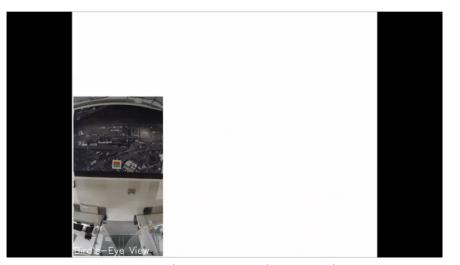
E. Cristofalo, K. Leahy, C.-I. Vasile, E. Montijano, M. Schwager and C. Belta, ISER 2016.

C. I. Vasile, K. Leahy, E. Cristofalo, A. Jones, M. Schwager and C. Belta, CDC 2016



Map unknown environment

Spec: "Always avoid all obstacles and Visit Marsh Plaza, Kenmore Square, Fenway Park, and Audubon Circle infinitely often and Bridge 2 should only be used for Northbound travel and Bridges 1 should only be used for Southbound travel. Uncertainty should always be below 0.9 m² and when crossing bridges it should be below 0.6 m²."



Localization and control

E. Cristofalo, K. Leahy, C.-I. Vasile, E. Montijano, M. Schwager and C. Belta, *ISER* 2016. C. I. Vasile, K. Leahy, E. Cristofalo, A. Jones, M. Schwager and C. Belta, CDC 2016

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Less conservative optimal TL control for large & simple dynamics

Signal Temporal Logic: Boolean (Qualitative) and Quantitative Semantics

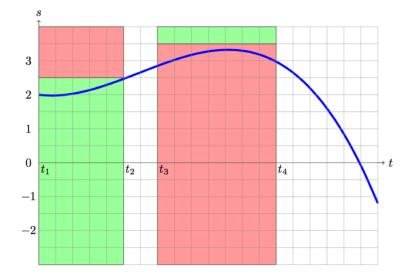
- Temporal operators are timed
- Semantics defined over signals
- Has qualitative semantics: real-valued function $~
 ho(s,\phi)$

$$\Box_{[t_1,t_2]}(s \le 2.5) \quad \Diamond_{[t_3,t_4]}(s > 3.5)$$

Boolean: True Quantitative: 0.01 Boolean: False Quantitative: -0.2

$$\Box_{[t_1,t_2]}(s \le 2.5) \land \Diamond_{[t_3,t_4]}(s > 3.5)$$

Boolean: False Quantitative: -0.2



- Boolean satisfaction of STL formulae over linear predicates can be mapped to feasibility of mixed integer linear equalities / inequalities (MILP feasibility)
- Robustness is piecewise affine in the integer and continuous variables

Less conservative optimal TL control for large & simple dynamics

Optimal STL Control

 $\min_{u^H} J(x^H, u^H) \quad \text{(any linear cost)}$

subject to

dynamics

 $x^+ = f(x,u)$ (any MLD system, e.g., piecewise affine) correctness

 x^{H}, u^{H} satisfy STL formula over linear predicates

Reduces to solving a MILP!

Less conservative optimal TL control for large & simple dynamics Planar Robot Example

$$x^{+} = x + u$$

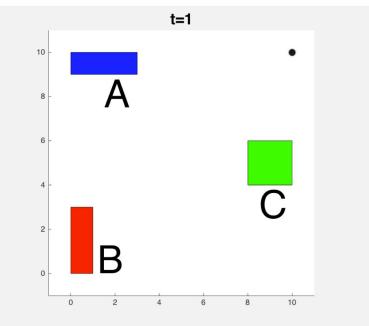
$$\varphi = \Box_{[40,50]} A \land \Diamond_{[0,40]} \Box_{[0,10]} B \land \Diamond_{[0,30]} C \qquad H = 50$$

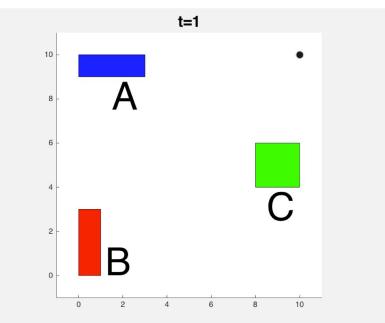
Minimum Fuel Only

$$J = \sum_{\tau=0}^{H-1} \left| u[\tau] \right|$$



$$J = -1000\rho + \sum_{\tau=0}^{H-1} |u[\tau]|$$





Sadraddini and Belta, Allerton Conference, 2015

Less conservative optimal TL control for large & simple dynamics

STL Model Predictive Control (MPC)

Repetitive tasks in infinite time: global STL formulas: $\Box_{[0,\infty]} \varphi$

$$\begin{split} u^{H}[t] &= \operatorname{argmin} \quad J(x^{H}[t], u^{H}[t]) \\ & \text{subject to} \quad x^{+} = f(x, u) \\ J &= J_{c} \\ J &= \rho \\ J &= -M(\rho - \|\rho\|) + J_{c} \end{split}$$

M is a large number. When ho < 0 , effectively maximize 2M
ho

Terminal constraints are guaranteed!

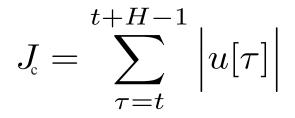
Less conservative optimal TL control for large & simple dynamics

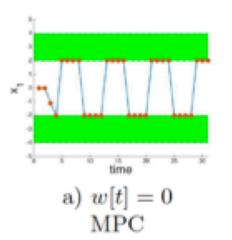
Example: Double Integrator

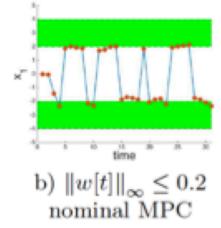
$$x^{+} = \left(\begin{array}{cc} 1 & 0.5\\ 0 & 0.8 \end{array}\right) x + \left(\begin{array}{c} 0\\ 1 \end{array}\right) + w$$

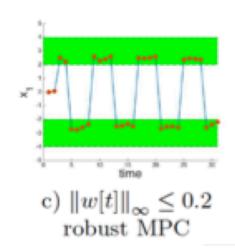
Spec: $\Box_{[0,\infty]} \left(\Diamond_{[0,4]} ((x_1 \le 4) \land (x_1 \ge 2)) \land \Diamond_{[0,4]} ((x_1 \ge -4) \land (x_1 \le -2)) \right)$

Minimize fuel consumption. If the spec becomes infeasible, maximize robustness.



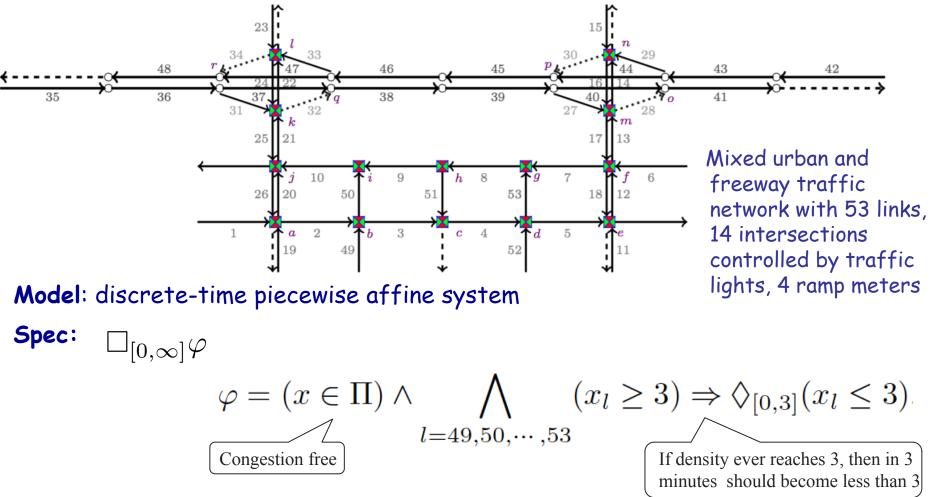






Sadraddini and Belta 2015

Less conservative *optimal* TL control for large & simple dynamics Example: Traffic network



Cost: delay over a given horizon

Takes less than 5 sec. to compute a optimal robust control strategy (MILP in 212 dimensions)

Sadraddini and Belta 2016

Summary

- Automata (Buchi, Rabin) games can be adapted to produce conservative TL control strategies for simple and small dynamical systems
- The above can be extended to conservative strategies for large and complicated systems by using I/O linearization techniques
- Partition refinement can be used to reduce conservatism for simple and small dynamical systems -> connection between optimality and TL correctness
- Sample-based techniques can be used to generate probabilistically complete TL strategies in high dimensions
- TL with quantitative semantics can be used for robust, provablycorrect optimal control in high dimensions

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Jana Tumova (now at KTH)







