CONTEXTUALITY OF TRANSFORMATIONS

Shane Mansfield

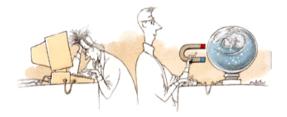
S. Mansfield and E. Kashefi. "Quantum advantage via sequential transformation contextuality".

Physical Review Letters 121 (23), 230401.



Contextuality Workshop Oxford June 6th 2019

Objective: Understand the **computational advantage** afforded by **quantum resources**



Objective: Understand the computational advantage afforded by quantum resources



If quantum systems really provide advantages over classical ones then it must come down to some behaviour that is inaccessible to classical resources

Objective: Understand the computational advantage afforded by quantum resources



If quantum systems really provide advantages over classical ones then it must come down to some behaviour that is inaccessible to classical resources

Questions we begin to have answers to:

- What are these **non-classical behaviours**? *Contextuality*, ...
- How do they relate to **computational advantage**? Specific examples, though not yet a systematic understanding...

Objective: Understand the computational advantage afforded by quantum resources



If quantum systems really provide advantages over classical ones then it must come down to some behaviour that is inaccessible to classical resources

Questions we begin to have answers to:

- What are these **non-classical behaviours**? *Contextuality*, ...
- How do they relate to **computational advantage**? Specific examples, though not yet a systematic understanding...

Where to go with this knowledge:

- Higher-level reasoning about contextuality in our resources and how to spend it!
- Systematic, structural path to quantum superiority

Non-Classicality

Non-classical behaviours should be established by no-go theorems

Behaviour	Familiar theorems
Nonlocality	Bell '64
Contextuality	Kochen, Specker '76
Non-macrorealism	Leggett, Garg '85
Generalised contextuality	Spekkens '05
ψ -ontology	Pusey, Barrett, Rudolph '12

Non-Classicality

Non-classical behaviours should be established by no-go theorems

Behaviour	Familiar theorems
Nonlocality	Bell '64
Contextuality	Kochen, Specker '76
Non-macrorealism	Leggett, Garg '85
Generalised contextuality	Spekkens '05
ψ -ontology	Pusey, Barrett, Rudolph '12
Dynamic contextuality	

• We posit a space of ontic states Λ

	Quantum mechanics	Ontological models
Preparation	ρ	$d_{\rho} \in P(\Lambda)$
Transformation	U	$f_U:\Lambda o \Lambda$
Measurement	М	$\xi_M:\Lambda o P(O)$

• We posit a space of ontic states Λ

	Quantum mechanics	Ontological models
Preparation	ρ	$d_{\rho} \in P(\Lambda)$
Transformation	U	$f_U:\Lambda \to \Lambda$
Measurement	М	$\xi_M:\Lambda o P(O)$

Empirical predictions

Empirical data $e_{\rho,U,M}$: P(O) should be reproduced as a weighted average over ontic states,

$$e_{
ho,U,M} = \sum_{oldsymbol{\lambda} \in \Lambda} d_{
ho}(oldsymbol{\lambda}) \, \xi_M(f_U(oldsymbol{\lambda}))$$

• We posit a space of ontic states Λ

	Quantum mechanics	Ontological models
Preparation	ρ	$d_{\rho} \in P(\Lambda)$
Transformation	U	$f_U:\Lambda \to \Lambda$
Measurement	М	$\xi_M:\Lambda \to P(O)$

Empirical predictions

Empirical data $e_{\rho,U,M}$: P(O) should be reproduced as a weighted average over ontic states,

$$e_{
ho,U,M} = \sum_{oldsymbol{\lambda} \in \Lambda} d_{
ho}(oldsymbol{\lambda}) \, \xi_M(f_U(oldsymbol{\lambda}))$$

• No-go theorems arise when ontological models and the Born Rule make differing predictions

• We posit a space of ontic states Λ

	Quantum mechanics	Ontological models
Preparation	ρ	$d_{\rho} \in P(\Lambda)$
Transformation	U	$f_U:\Lambda \to \Lambda$
Measurement	М	$\xi_M:\Lambda o P(O)$

Empirical predictions

Empirical data $e_{\rho,U,M}$: P(O) should be reproduced as a weighted average over ontic states,

$$e_{
ho,U,M} = \sum_{oldsymbol{\lambda} \in \Lambda} d_{
ho}(oldsymbol{\lambda}) \, \xi_M(f_U(oldsymbol{\lambda}))$$

- No-go theorems arise when ontological models and the Born Rule make differing predictions
- These require additional structural assumptions

• We posit a space of ontic states Λ

	Quantum mechanics	Ontological models
Preparation	ρ	$d_{\rho} \in P(\Lambda)$
Transformation	U	$f_U:\Lambda \to \Lambda$
Measurement	М	$\xi_M:\Lambda \to P(O)$

Empirical predictions

Empirical data $e_{\rho,U,M}$: P(O) should be reproduced as a weighted average over ontic states,

$$e_{
ho,U,M} = \sum_{\lambda \in \Lambda} d_{
ho}(\lambda) \, \xi_M(f_U(\lambda))$$

- No-go theorems arise when ontological models and the Born Rule make differing predictions
- These require additional structural assumptions

*No further assumptions about *features* at this stage (cf. Spekkens)

*Non*contextuality

*Non*contextuality

• Context: a set of compatible measurements

$$C = \{M_1, \ldots, M_n\}$$

Noncontextuality

• Context: a set of compatible measurements

$$C = \{M_1, \ldots, M_n\}$$

• Ontological representations respect compatibility

$$\xi_C(\lambda) = \prod_{M \in C} \xi_M(\lambda)$$

*Non*contextuality

• Context: a set of compatible measurements

$$C = \{M_1,\ldots,M_n\}$$

• Ontological representations respect compatibility

$$\xi_C(\lambda) = \prod_{M \in C} \xi_M(\lambda)$$

• ... and are context-independent; e.g. if $M \in C, C'$

$$\xi_{M^{(C)}} = \xi_{M^{(C')}}$$

Noncontextuality of transformations

Noncontextuality of transformations

• A context is a convex decomposition of a fixed transformation, e.g.

$$T = \frac{1}{2}U_a + \frac{1}{2}U_A$$
(C)
$$T = \frac{1}{3}U_a + \frac{1}{3}U_b + \frac{1}{3}U_c.$$
(C')

Noncontextuality of transformations

• A context is a convex decomposition of a fixed transformation, e.g.

$$T = \frac{1}{2}U_a + \frac{1}{2}U_A$$
(C)
$$T = \frac{1}{3}U_a + \frac{1}{3}U_b + \frac{1}{3}U_c.$$
(C')

• Ontological representations respect convex decompositions, e.g.

$$f_T = \frac{1}{2}f_{U_a} + \frac{1}{2}f_{U_A} = \frac{1}{3}f_{U_a} + \frac{1}{3}f_{U_b} + \frac{1}{3}f_{U_c}$$

Noncontextuality of transformations

• A context is a convex decomposition of a fixed transformation, e.g.

$$T = \frac{1}{2}U_a + \frac{1}{2}U_A$$
(C)
$$T = \frac{1}{3}U_a + \frac{1}{3}U_b + \frac{1}{3}U_c.$$
(C')

• Ontological representations respect convex decompositions, e.g.

$$f_T = \frac{1}{2}f_{U_a} + \frac{1}{2}f_{U_A} = \frac{1}{3}f_{U_a} + \frac{1}{3}f_{U_b} + \frac{1}{3}f_{U_c}$$

• ... and are context-independent

$$f_{U_a^{(C)}} = f_{U_a^{(C')}}$$

Noncontextuality of transformations in sequence

Noncontextuality of transformations in sequence

• A context is a sequence $U_n \circ U_{n-1} \circ \cdots \circ U_1$

Noncontextuality of transformations in sequence

- A context is a sequence $U_n \circ U_{n-1} \circ \cdots \circ U_1$
- Ontological representations respect sequentiality

$$f_{U_n \circ \cdots \circ U_1} = f_{U_n} \circ \cdots \circ f_{U_1}$$

Noncontextuality of transformations in sequence

- A context is a sequence $U_n \circ U_{n-1} \circ \cdots \circ U_1$
- Ontological representations respect sequentiality

$$f_{U_n \circ \cdots \circ U_1} = f_{U_n} \circ \cdots \circ f_{U_1}$$

• ... and are context-independent

$$f_{U^{(C)}}=f_{U^{(C^\prime)}}$$

Noncontextuality of transformations in sequence

- A context is a sequence $U_n \circ U_{n-1} \circ \cdots \circ U_1$
- Ontological representations respect sequentiality

$$f_{U_n \circ \cdots \circ U_1} = f_{U_n} \circ \cdots \circ f_{U_1}$$

• ... and are context-independent

$$f_{U^{(C)}}=f_{U^{(C^\prime)}}$$

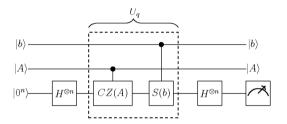
"Contextuality in our sense implies that the system of study cannot have an ontology in which transformations correspond to modular, composable operations on ontic states, such that they are well-defined independently of which transformations may have been performed previously or will be performed subsequently."

Example: Quantum Advantage in Shallow Circuits

Bravyi, Gossett, König, Science, 2017

There exists a task for which

Circuit-depth complexity	
Classical	log
Quantum	constant

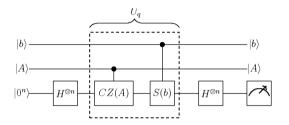


Example: Quantum Advantage in Shallow Circuits

Bravyi, Gossett, König, Science, 2017

There exists a task for which

Circuit-depth complexity	
Classical	log
Quantum	constant



Contextuality analysis

Suppose sequential noncontextuality:

$$f_{U_n \circ \cdots \circ U_1} = f_{U_n} \circ \cdots \circ f_{U_1}$$

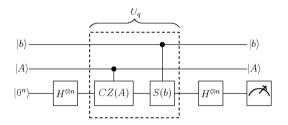
then we should have constant depth in the classical case too!

Example: Quantum Advantage in Shallow Circuits

Bravyi, Gossett, König, Science, 2017

There exists a task for which

Circuit-depth complexity	
Classical	log
Quantum	constant



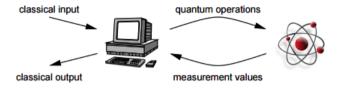
Contextuality analysis

Suppose sequential noncontextuality:

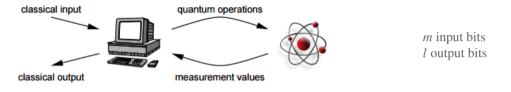
$$f_{U_n \circ \cdots \circ U_1} = f_{U_n} \circ \cdots \circ f_{U_1}$$

then we should have constant depth in the classical case too!

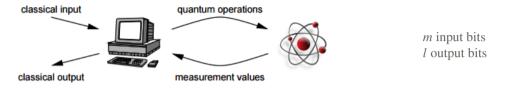
• Sequential contextuality is necessary for advantage







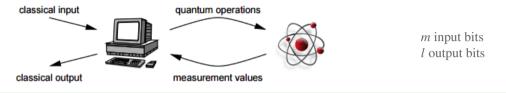
- Classical control computer
 - Determines the sequence of measurements
 - In *this example* can perform \mathbb{Z}_2 -linear computations only
- Power to compute non-linear functions may reside in certain (quantum) resources



- Classical control computer
 - Determines the sequence of measurements
 - In *this example* can perform \mathbb{Z}_2 -linear computations only
- Power to compute non-linear functions may reside in certain (quantum) resources

Raussendorf, PRA, 2013

- ℓ 2-MBQC *deterministically* computes a non-linear Boolean function $f: 2^m \longrightarrow 2^l$ iff the resource is *maximally contextual*
- Bounded error non-linear computations require *contextuality*



- Classical control computer
 - Determines the sequence of measurements
 - In *this example* can perform \mathbb{Z}_2 -linear computations only
- Power to compute non-linear functions may reside in certain (quantum) resources

Abramsky, Barbosa, M, PRL, 2017



quantifiable relationship!

Is Dynamic Contextuality Useful?

Dunjko, Kapourniotis, Kashefi, QIC, 2016.



Classical control $(\oplus L)$:

- Classical inputs $a, b \in \mathbb{Z}_2$
- Controls transformations
- Announces meas. outcome $\{+1 \mapsto 0, -1 \mapsto 1\}$

Quantum resource:

- Prepare qubit in state $|+\rangle$
- Transformations

 $U_0 = V_0 = W_0 = I$ $U_1 = V_1 = W_1 = R_z(\pi/2)$

$$f(a,b) = a \otimes_2 b$$

Is Dynamic Contextuality Useful?

Dunjko, Kapourniotis, Kashefi, QIC, 2016.



Classical control $(\oplus L)$:

- Classical inputs $a, b \in \mathbb{Z}_2$
- Controls transformations
- Announces meas. outcome $\{+1 \mapsto 0, -1 \mapsto 1\}$

Quantum resource:

- Prepare qubit in state $|+\rangle$
- Transformations

$$U_0 = V_0 = W_0 = I$$

 $U_1 = V_1 = W_1 = R_z(\pi/2)$

$$f(a,b) = a \otimes_2 b$$

- Boosts computational power: $\oplus L \longrightarrow P$
- Contextuality in the traditional sense cannot arise with a single qubit!
- So what, if anything, is the non-classical behaviour?

An Appropriate Ontology

$$|+\rangle - U(a) - V(b) - W(a \oplus b) - \checkmark \sigma_X$$

• Classical computers can do $a \otimes b$, and constitute reasonable ontologies

An Appropriate Ontology

$$|+\rangle - U(a) - V(b) - W(a \oplus b) - \checkmark \sigma_X$$

- Classical computers can do $a \otimes b$, and constitute reasonable ontologies
- Artificial problem: boost $\oplus L \longrightarrow P$
- Meaningless without $\oplus L$ restriction on ontological models
- $\oplus L$ circuits are built from NOTs and CNOTs

An Appropriate Ontology

$$|+\rangle - U(a) - V(b) - W(a \oplus b) - \checkmark \sigma_X$$

- Classical computers can do $a \otimes b$, and constitute reasonable ontologies
- Artificial problem: boost $\oplus L \longrightarrow P$
- Meaningless without $\oplus L$ restriction on ontological models
- $\oplus L$ circuits are built from NOTs and CNOTs

Commutative \oplus *L***-ontological models:**

Ontic states	$\Lambda = (\mathbb{Z}_2)^n$
Transformations	$f_U(oldsymbol{\lambda}) = (I \oplus A_U) oldsymbol{\lambda} \oplus oldsymbol{u}$
Measurements	$\xi_M(oldsymbol{\lambda}) = [(I \oplus A_M) oldsymbol{\lambda} \oplus oldsymbol{u}] \cdot oldsymbol{\delta}$

Parity Proof of Contextuality

• \oplus *L*-ontological realisation of the protocol requires the following equations to be satisfied:

 $\begin{aligned} [\boldsymbol{\lambda} \oplus A_U(0)\boldsymbol{\lambda} \oplus A_V(0)\boldsymbol{\lambda} \oplus A_W(0)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(0) \oplus \boldsymbol{v}(0) \oplus \boldsymbol{w}(0) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 0\\ [\boldsymbol{\lambda} \oplus A_U(0)\boldsymbol{\lambda} \oplus A_V(1)\boldsymbol{\lambda} \oplus A_W(1)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(0) \oplus \boldsymbol{v}(1) \oplus \boldsymbol{w}(1) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 0\\ [\boldsymbol{\lambda} \oplus A_U(1)\boldsymbol{\lambda} \oplus A_V(0)\boldsymbol{\lambda} \oplus A_W(1)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(1) \oplus \boldsymbol{v}(0) \oplus \boldsymbol{w}(1) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 0\\ [\boldsymbol{\lambda} \oplus A_U(1)\boldsymbol{\lambda} \oplus A_V(1)\boldsymbol{\lambda} \oplus A_W(0)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(1) \oplus \boldsymbol{v}(1) \oplus \boldsymbol{w}(0) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 1\end{aligned}$

Parity Proof of Contextuality

• \oplus *L*-ontological realisation of the protocol requires the following equations to be satisfied:

 $\begin{aligned} &[\boldsymbol{\lambda} \oplus A_U(0)\boldsymbol{\lambda} \oplus A_V(0)\boldsymbol{\lambda} \oplus A_W(0)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(0) \oplus \boldsymbol{v}(0) \oplus \boldsymbol{w}(0) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} = 0 \\ &[\boldsymbol{\lambda} \oplus A_U(0)\boldsymbol{\lambda} \oplus A_V(1)\boldsymbol{\lambda} \oplus A_W(1)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(0) \oplus \boldsymbol{v}(1) \oplus \boldsymbol{w}(1) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} = 0 \\ &[\boldsymbol{\lambda} \oplus A_U(1)\boldsymbol{\lambda} \oplus A_V(0)\boldsymbol{\lambda} \oplus A_W(1)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(1) \oplus \boldsymbol{v}(0) \oplus \boldsymbol{w}(1) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} = 0 \\ &[\boldsymbol{\lambda} \oplus A_U(1)\boldsymbol{\lambda} \oplus A_V(1)\boldsymbol{\lambda} \oplus A_W(0)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(1) \oplus \boldsymbol{v}(1) \oplus \boldsymbol{w}(0) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} = 1 \end{aligned}$

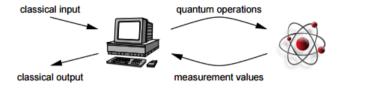
- System of equations is not jointly satisfiable
- Sum RHS: odd
- Sum LHS: even (each vector appears even number of times)

Parity Proof of Contextuality

• \oplus *L*-ontological realisation of the protocol requires the following equations to be satisfied:

 $\begin{aligned} [\boldsymbol{\lambda} \oplus A_U(0)\boldsymbol{\lambda} \oplus A_V(0)\boldsymbol{\lambda} \oplus A_W(0)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(0) \oplus \boldsymbol{v}(0) \oplus \boldsymbol{w}(0) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 0\\ [\boldsymbol{\lambda} \oplus A_U(0)\boldsymbol{\lambda} \oplus A_V(1)\boldsymbol{\lambda} \oplus A_W(1)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(0) \oplus \boldsymbol{v}(1) \oplus \boldsymbol{w}(1) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 0\\ [\boldsymbol{\lambda} \oplus A_U(1)\boldsymbol{\lambda} \oplus A_V(0)\boldsymbol{\lambda} \oplus A_W(1)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(1) \oplus \boldsymbol{v}(0) \oplus \boldsymbol{w}(1) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 0\\ [\boldsymbol{\lambda} \oplus A_U(1)\boldsymbol{\lambda} \oplus A_V(1)\boldsymbol{\lambda} \oplus A_W(0)\boldsymbol{\lambda} \oplus A_M\boldsymbol{\lambda} \oplus \boldsymbol{u}(1) \oplus \boldsymbol{v}(1) \oplus \boldsymbol{w}(0) \oplus \boldsymbol{m}] \cdot \boldsymbol{\delta} &= 1\end{aligned}$

- System of equations is not jointly satisfiable
- Sum RHS: odd
- Sum LHS: even (each vector appears even number of times)
- Such a realisation is necessarily dynamically contextual

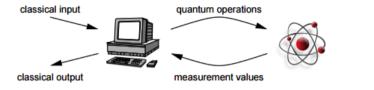




*l*2-TBQC

 $\oplus L$ control computer with access to a (quantum) resource

- Fixed preparation and 2-outcome measurement
- Controlled unitary operations



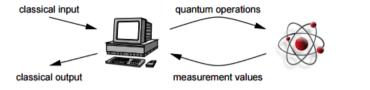
m input bits 1 output bit

*l*2-TBQC

 $\oplus L$ control computer with access to a (quantum) resource

- Fixed preparation and 2-outcome measurement
- Controlled unitary operations

• Examples: DKK, CHSH*



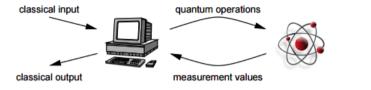


l2-TBQC

 $\oplus L$ control computer with access to a (quantum) resource

- Fixed preparation and 2-outcome measurement
- Controlled unitary operations

- Examples: DKK, CHSH*
- 12-MBQC $\neq 12$ -TBQC





l2-TBQC

 $\oplus L$ control computer with access to a (quantum) resource

- Fixed preparation and 2-outcome measurement
- Controlled unitary operations

- Examples: DKK, CHSH*
- 12-MBQC $\neq 12$ -TBQC
- In general we could shift perspective and express any MBQC as a TBQC

Empirical model: for each context C, a distribution e_C over possible outcomes

 $e = \{e_C\}$

Cf. Abramsky, Brandenburger, NJP, 2011.

Empirical model: for each context C, a distribution e_C over possible outcomes

 $e = \{e_C\}$

Cf. Abramsky, Brandenburger, NJP, 2011.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

E.g. CHSH* strategy

Empirical model: for each context C, a distribution e_C over possible outcomes

 $e = \{e_C\}$

Cf. Abramsky, Brandenburger, NJP, 2011.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

E.g. CHSH* strategy

context		outcome		
	а	b	o = 0	o = 1
C_0	0	0	3/4	1/4
C_1	0	1	1/4	3/4
C_2	1	0	1/4	3/4
<i>C</i> ₃	1	1	0	1

Empirical model: for each context C, a distribution e_C over possible outcomes

 $e = \{e_C\}$

Cf. Abramsky, Brandenburger, NJP, 2011.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

E.g. CHSH* strategy

context		outcome		
	а	b	o = 0	o = 1
C_0	0	0	3/4	1/4
C_1	0	1	1/4	3/4
C_2	1	0	1/4	3/4
<i>C</i> ₃	1	1	0	1

$$e_{C_0} = (3/4, 1/4)$$

$$e_{C_1} = (1/4, 3/4)$$

$$e_{C_2} = (1/4, 3/4)$$

$$e_{C_3} = (0, 1)$$

Quantifying Contextuality

Noncontextual fraction NCF(e) NCF(e) is max ω over all decompositions

$$e = \omega e^{\mathsf{NC}} + (1 - \omega)e'$$

s.t. e^{NC} is noncontextual

Contextual fraction CF(e)

$$CF := 1 - NCF(e)$$

• $CF(e), NCF(e) \in [0,1]$

Quantifying Hardness

Distance on functions: Given $f, g : (\mathbb{Z}_2)^r \to \mathbb{Z}_2$,

$$d(f,g) := 2^{-r} |\{ i \mid f(i) \neq g(i) \}|$$

Fraction of inputs for which outputs differ

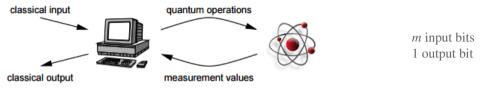
Non-linearity of a function $f : (\mathbb{Z}_2)^r \to \mathbb{Z}_2$,

$$\mathbf{v}(f) := \min \left\{ d(f,g) \mid g : (\mathbb{Z}_2)^r \to \mathbb{Z}_2 \text{ linear} \right\}$$

Distance to nearest linear function

Cf. Abramsky, Barbosa, M, PRL, 2017.

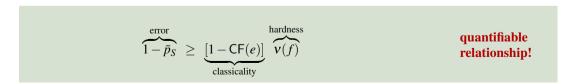
Advantage in *l*2-TBQC



l2-TBQC

 $\oplus L$ control computer with access to a (quantum) resource

- Fixed preparation and 2-outcome measurement
- Controlled unitary operations



Example: The CHSH* Game

Henaut, Catani, Browne, Pappa, M, PRA, 2018.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

- Task: compute $a \otimes_2 b$
- Maximise success probability in various regimes
- Tsirelson bound for qubits
- Similar for qutrits, with \otimes_3 , etc.
- Dimensional witness!

	$p_{\rm success}^{\rm max}$
bit	0.75
Spekkens toy bit	0.75
stabiliser qubit	0.75
qubit	0.85
qutrit	1

Classical Erasure

Henaut, Catani, Browne, Pappa, M, PRA, 2018.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

• Classically, can compute $a \otimes_2 b$ with *l*2-operations and **erasure**

 $U_0 = I$ $U_1 = NOT$ $V_0 = RESET_0$ $V_1 = I$

Classical Erasure

Henaut, Catani, Browne, Pappa, M, PRA, 2018.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

• Classically, can compute $a \otimes_2 b$ with *l*2-operations and **erasure**

$$U_0 = I$$
 $U_1 = \text{NOT}$ $V_0 = \text{RESET}_0$ $V_1 = I$

• Undesirable for an ontological model!

 $f_I \neq I$

Classical Erasure

Henaut, Catani, Browne, Pappa, M, PRA, 2018.

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

• Classically, can compute $a \otimes_2 b$ with *l*2-operations and **erasure**

$$U_0 = I$$
 $U_1 = \mathsf{NOT}$ $V_0 = \mathsf{RESET}_0$ $V_1 = I$

• Undesirable for an ontological model!

 $f_I \neq I$

• Expected *erasure cost* per run, averaged over pairs of inputs, to compute a function g, with 1- and 2-bit gates coincides with non-linearity measure:

Contextuality-Erasure Tradeoff

 $|+\rangle - U(a) - V(b) - \checkmark \sigma_X$

Contextuality-Erasure Tradeoff

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

Landauer's Principle

Erasure of a bit results in an entropy increase of at least $kT \ln 2$ in the non-information-bearing degrees of freedom of the system

Contextuality-Erasure Tradeoff

$$|+\rangle - U(a) - V(b) - \checkmark \sigma_X$$

Landauer's Principle

Erasure of a bit results in an entropy increase of at least $kT \ln 2$ in the non-information-bearing degrees of freedom of the system

The Combined Perspective

If an *l*2-TBQC is run *n* times with uniformly random inputs and the overall change in environmental entropy is ΔS , then

$$\overline{\varepsilon} \ge \left[\mathsf{NCF}(e) - \frac{\Delta S}{n \, kT \ln 2}\right] \tilde{v}(g)$$

Equivalently,

$$\mathsf{CF}(e) \geq 1 - \frac{\overline{\varepsilon}}{\widetilde{\nu}(g)} - \frac{\Delta S}{n \, kT \ln 2}$$

• Another way to be contextual! – **dynamic contextuality**

- Another way to be contextual! **dynamic contextuality**
- Quantifiably relates to quantum advantage in l2-TBQCs

- Another way to be contextual! dynamic contextuality
- Quantifiably relates to quantum advantage in l2-TBQCs
- Results parallel Anders and Browne, Raussendorf, ABM for BKS contextuality

- Another way to be contextual! **dynamic contextuality**
- Quantifiably relates to quantum advantage in l2-TBQCs
- Results parallel Anders and Browne, Raussendorf, ABM for BKS contextuality
- Available to single qubits

- Another way to be contextual! dynamic contextuality
- Quantifiably relates to quantum advantage in l2-TBQCs
- Results parallel Anders and Browne, Raussendorf, ABM for BKS contextuality
- Available to single qubits
- Dimensional/irreversibility/quantumness witnesses

- Another way to be contextual! dynamic contextuality
- Quantifiably relates to quantum advantage in l2-TBQCs
- Results parallel Anders and Browne, Raussendorf, ABM for BKS contextuality
- Available to single qubits
- Dimensional/irreversibility/quantumness witnesses

Some directions

• Where else does it play a role?

E.g. other single qubit advantages (Knill-Laflamme, Galvão), other informatic tasks, universal QC? Reinforcement learning? Indefinite causal structures?

- Another way to be contextual! **dynamic contextuality**
- Quantifiably relates to quantum advantage in l2-TBQCs
- Results parallel Anders and Browne, Raussendorf, ABM for BKS contextuality
- Available to single qubits
- Dimensional/irreversibility/quantumness witnesses

Some directions

- Where else does it play a role?
 - E.g. other single qubit advantages (Knill-Laflamme, Galvão), other informatic tasks, universal QC? Reinforcement learning? Indefinite causal structures?
- Experimental tests

- Another way to be contextual! dynamic contextuality
- Quantifiably relates to quantum advantage in l2-TBQCs
- Results parallel Anders and Browne, Raussendorf, ABM for BKS contextuality
- Available to single qubits
- Dimensional/irreversibility/quantumness witnesses

Some directions

- Where else does it play a role?
 - E.g. other single qubit advantages (Knill-Laflamme, Galvão), other informatic tasks, universal QC? Reinforcement learning? Indefinite causal structures?
- Experimental tests
- Circuit contextuality: generalise BKS and dynamic contextuality

• Classicality is characterised by the presence of **structure preserving partial functors** from category of quantum circuits to category of calssical circuits

- Classicality is characterised by the presence of **structure preserving partial functors** from category of quantum circuits to category of calssical circuits
- I.e. 'Shape' of circuit is preserved and components appearing in different circuits are represented in same way

- Classicality is characterised by the presence of **structure preserving partial functors** from category of quantum circuits to category of calssical circuits
- I.e. 'Shape' of circuit is preserved and components appearing in different circuits are represented in same way
- Kind of structure identifies kind of classicality

- Classicality is characterised by the presence of **structure preserving partial functors** from category of quantum circuits to category of calssical circuits
- I.e. 'Shape' of circuit is preserved and components appearing in different circuits are represented in same way
- Kind of structure identifies kind of classicality
- Can subsume existing notions of classicality

- Classicality is characterised by the presence of **structure preserving partial functors** from category of quantum circuits to category of calssical circuits
- I.e. 'Shape' of circuit is preserved and components appearing in different circuits are represented in same way
- Kind of structure identifies kind of classicality
- Can subsume existing notions of classicality

	Components	Composition
Locality	М	\otimes
Noncontextuality (BKS)	М	×
Measurement NC (Spekkens*)	М	×
Preparation NC (Spekkens*)	Р	$+_{\lambda}$
Transformation NC (Spekkens*)	Т	$+_{\lambda}$
Preparation Independence (PBR)	Р	\otimes_{\min}
Subsystem Condition (SM)	Р	\otimes
Dynamic NC	Т	0