

# Tensor Comprehensions in SaC

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# Tensors, Tensor Notation, Einstein Notation, Ricci Calculus,... and their applications

- N-dimensional index spaces for data and operations on them
- Notation omits ranges and boundaries whenever “obvious”
- Typically nested!

- General Activation Formula:  $a_j^{[l]} = g^{[l]}(\sum_k w_{jk}^{[l]} a_k^{[l-1]} + b_j^{[l]}) = g^{[l]}(z_j^{[l]})$

- $J(x, W, b, y)$  or  $J(\hat{y}, y)$  denote the cost function.

## Examples of cost function:

- $J_{CE}(\hat{y}, y) = -\sum_{i=0}^m y^{(i)} \log \hat{y}^{(i)}$

- $J_1(\hat{y}, y) = \sum_{i=0}^m |y^{(i)} - \hat{y}^{(i)}|$

Isn't that just array comprehensions?  
--- or .... getting to the point?

[ expr | id<sub>1</sub> ← generator<sub>1</sub>, ..., id<sub>n</sub> ← generator<sub>n</sub> ]

Here we need lower  
and upper bounds

We need to  
“compose” the  
generators!

We want non-inherently-  
sequential generators!

# Transposition

Mathematics:

$$a^T_{i,j} = a_{j,i}$$

Composition is always orthogonal

"typical" array comprehension:

```
aT = [ a[j,i] | i in 0 .. shape(a)[1], j in 0 .. shape(a)[0] ]
```

SaC 1.0:

```
aT = { [i,j] -> a[j,i] };
```

Set-notation is mapped into data-parallelism

Lower and upper bounds are inferred if possible

# Element-wise addition

Mathematics:

$$c_{\vec{T}} = a_{\vec{T}} + b_{\vec{T}}$$

scalar addition!

"typical" array comprehension:

```
c = [ a[i] + b[i] | i in 0 .. shape(a)[0] ]
```

recursive call !

SaC 1.0:

```
c = { iv -> a[iv] + b[iv] };
```

scalar addition!

vector of indices!

```
c = { [i] -> a[i] + b[i] };
```

recursive call !

## Our physics example

$$a_j^{[l]} = g^{[l]} \left( \sum_k w_{jk}^{[l]} a_k^{[l-1]} + b_j^{[l]} \right)$$

in SaC 1.0:

```
a = { [j] -> g * (sum( {[k] -> w[j,k]*a[k] }) + b[j]) };
```

# Concatenation

Bound inference fails !

```
c = { [i] -> ( i < len(a) ? a[i] : b[i-len(a)] ) };
```

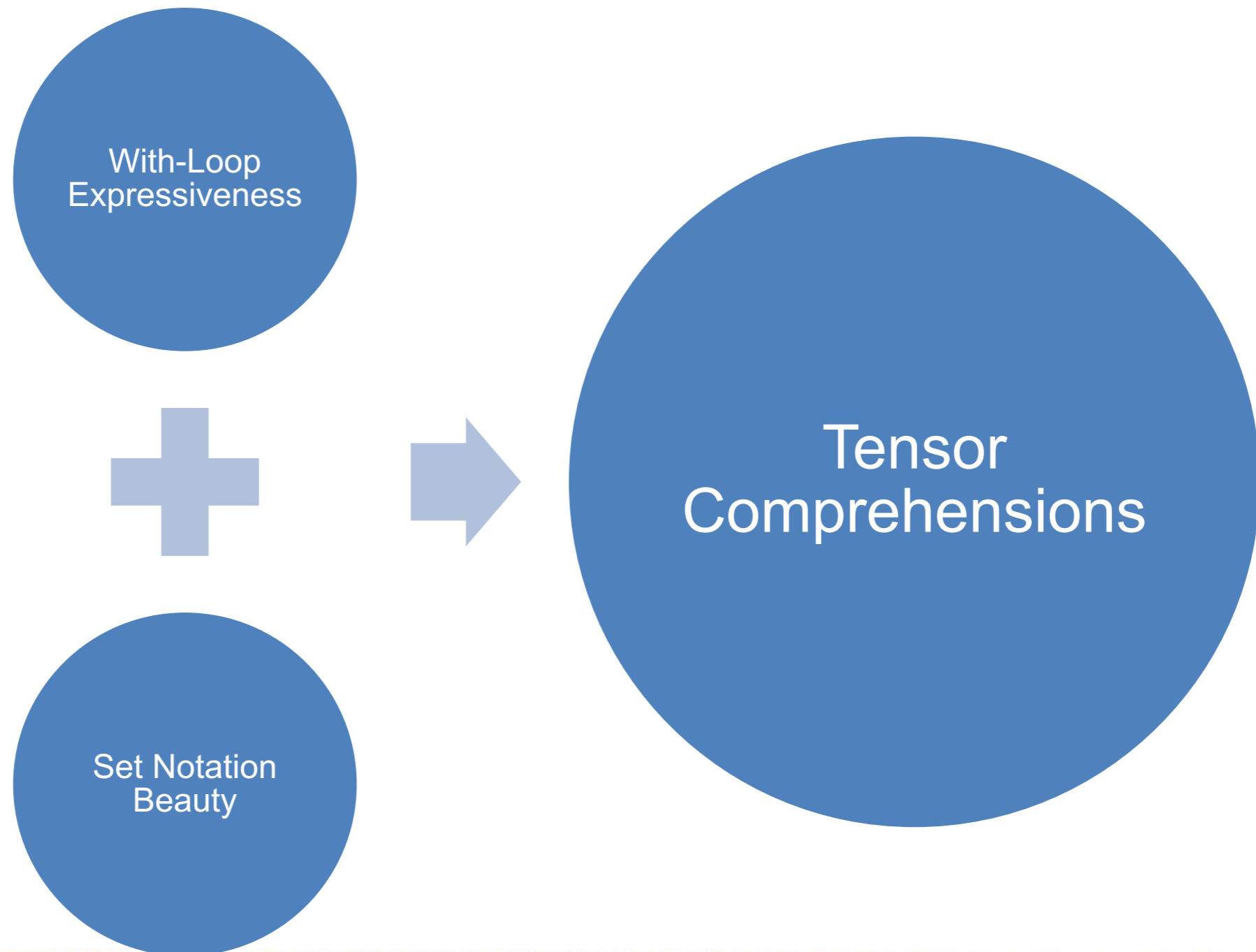
Conditional is ugly!

SaC 1.4, Tensor Comprehensions!

```
c = { [i] -> a[i] ;  
      [i] -> b[i-len(a)] | [i] < shape(a) + shape(b) };
```



# Aim



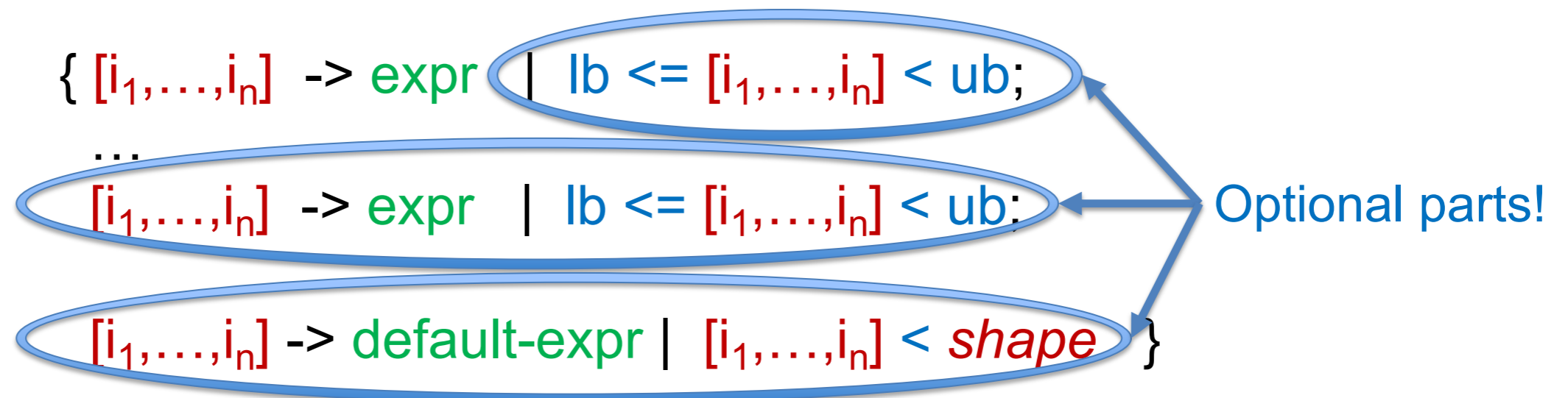
# Tensor Comprehensions: Full With-Loop expressiveness in Set Notation!

```
with {  
  ( lb <= [i1,...,in] < ub ) : expr;  
  ...  
  ( lb <= [i1,...,in] < ub ) : expr;  
} : genarray( shape, default-expr)
```



```
{ [i1,...,in] -> expr | lb <= [i1,...,in] < ub;  
  ...  
  [i1,...,in] -> expr | lb <= [i1,...,in] < ub;  
  [i1,...,in] -> default-expr | [i1,...,in] < shape }
```

## Beauty Measure #1: Make some parts optional & use the Set Notation inference



## Beauty Measure #2: Extend the Inference

Example: `take (int[.] s, int[*] a):`

`{ iv -> a[iv] | iv < s }`

default-expr is missing!

Type-inference cannot help!

Consider `take ([0], a)` where `a::int[0,7]` !

New Inference:

`{ iv -> a[iv] | iv < s;`  
`iv -> genarray (drop (shape (s), shape (a)), 0) }`

## Key Idea of the Inference

$\{ \text{iv} \rightarrow a[\text{iv}] \mid \text{iv} < s \}$



Generate default from one expression

$\{ \text{iv} \rightarrow a[\text{iv}] \mid \text{iv} < s;$   
 $\text{iv} \rightarrow \text{genarray}(\text{shape}(a[0*s]), \text{zero}(a[0*s])) \}$



Rewrite to manifest some laziness

$\{ \text{iv} \rightarrow a[\text{iv}] \mid \text{iv} < s;$   
 $\text{iv} \rightarrow \text{genarray}(\text{drop}(\text{shape}(s), \text{shape}(a)), 0) \}$

# Leveraging Demand Analysis

```
genarray (shape (a[0*s]), zero (a[0*s]))
```

=> Analysis of selection yields: in order to compute the shape of `a[0*s]`, we only need to know the shape of `a` and the shape of `s`!  
(for details see “A Binding Scope Analysis for Generic Programs on Arrays”, IFL’05)

=> A systematic rewrite of the definition of selection yields that  
`shape (a[0*s]) = sel_s( shape(s), shape(a)) = drop(shape(s), shape(a))`  
(for details see “Tensor Comprehensions in SaC”, IFL’19)

Hence, we get overall:

```
genarray (drop (shape (s), shape (a)), 0)
```

# Conclusions

- Array Comprehensions in the context of n-dimensional arrays / arbitrary tensors with homogeneous nesting is surprisingly challenging!
  - Full Expressiveness leads to very extensive specifications.
  - Range inference is non-trivial.
  - Default element inference is even harder.
- The Tensor Comprehensions presented here offer:
  - Full expressiveness
  - Flexibility in the degree of specificational demand
  - Novel default element inference that is independent of the type system
- Leads to a mechanism for manifesting laziness with an eager execution mechanism

