Aggregation in Probabilistic Databases

Robert Fink, University of Oxford rf@robertfink.de

Télécom ParisTech DBWeb Seminar January 26th 2012

Outline

Biomass Plants in the US



AES Mendota Biomass Facility
APS Biomass I Biomass Facility
Aberdeen Biomass Facility
Acme Landfill Biomass Facility
Adrian Energy Associates LLC Biomass
Facility
Agrilectric Power Partners Ltd Biomass
Facility
Al Turi Biomass Facility
Alabama Pine Pulp Biomass Facility
Albany Landfill Gas Utilization Project
Biomass Facility
Alexandria Biomass Facility
Altamont Gas Recovery Biomass
Facility
American Canyon Power Plant Biomass
Facility
American Ref-Fuel of Delaware Valley

Name	Alabama Pine Pulp Biomass Facility
Facility	Alabama Pine Pulp
Sector	Biomass
Location	Monroe County, Alabama
Coordinates	31.5119068°, -87.460397° © Display map
Generating Capacity (MW)	32.085
Commercial Online Date	1991

Governors in US States



Governor's Name	State	Time in Office	Party
Gov. Robert Bentley	Alabama	(2010 -) (2011 -)	Republican
Gov. Bob Riley	Alabama	(2003 - 2011)	Republican
Gov. Donald Eugene Siegelman	Alabama	(1998 -) (1999 - 2003)	Democrat
Gov. James Elisha Folsom	Alabama	(1993 - 1995)	Democrat
Gov. Harold Guy Hunt	Alabama	(1990 -) (1987 - 1993)	Republican











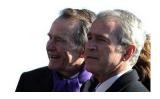






Who is responsible for a larger capacity of biogas plants, Democrats or Republicans?

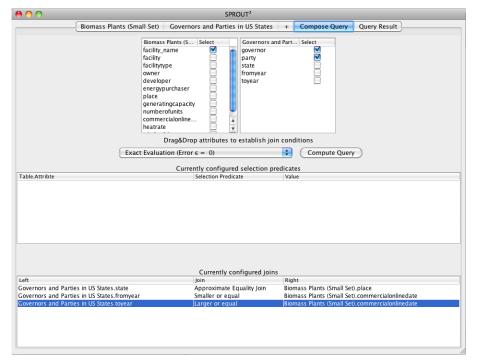




Biomass Plants (Small Set) Governors and Parties in US States + Compose Query Query Result

Biomass Plants (Smal	l Set)			•			Retrieve Goo	gle Fusion [*]	Гable	
acility_name facility	facilitytype	owner	developer	energypur	place	generating	number of	commerci	heatrate	windturbin
ES Mend AES Mend					Fresno Co	25		1989-01	17,873.6	
PS Bioma APS Bioma					Arizona	2.85		2006-01	8,911	
Aberdeen Aberdeen		Sierra Paci			Aberdeen,	12		0001-01		
Acme Lan Acme Lan	Landfill Gas				Contra Co	0.27		2003-01	12,916.67	
drian En Adrian En	Landfill Gas				Lenawee	2.4		1994-01	13,170.6	
grilectric Agrilectric					Calcasieu	12.2		1984-01	17,327.1	
l Turi Bio Al Turi	Landfill Gas				Orange Co	4.4		1988-01	15,600.2	
Nabama Alabama					Monroe C	32.085		1991-01	15,826.23	
lbany La Albany La	Landfill Gas				Albany Co	1.8		1998-01	11.913.9	
lexandri Alexandria		Indeck			Alexandri			0001-01		
Altamont Altamont	Landfill Gas				Alameda	2.6		2002-01	10,500	
American American	Landfill Gas				Napa Cou	1.4		1985-01		
merican American					Delaware	80		1991-01	18.674.9	
merican American					Essex Cou	60		1990-01		
merican American					Nassau Co	67.7		1989-01		
merican American					Niagara C			1980-01		
merican American					New Lond			1991-01		
Arbor Hills Arbor Hills	Landfill Gas				Washtena	19		1996-01		
Archbald Archbald					Lackawan			1988-01		
Ashland Bi Ashland		Boralex			Ashland			0001-01		
Atlantic Cit Atlantic Ci	Landfill Gas				New Jersey			2004-01	12.916.67	
Atlantic Co Atlantic C					Atlantic Co			2005-01		
von Ener Avon Ener					Cook Cou			1997-01		
J Gas Re BJ Gas Re					Gwinnett	2.4		1993-01		
KK Landf BKK Landfill					Los Angel			1993-01		
alefill Lan Balefill La					Bergen Co			1998-01		
arre Bio Barre	Landfill Gas				Worcester			1996-01		
aton Rog Baton Roque		Agrilectric			Lake Charl			0001-01		
lavarian L Bavarian L		rigi neeti ie			Boone Cou			2003-01		
av Front Bay Front	Luncilli Gas				Ashland C			1952-01		
ay Resou Bay Resou	Municipal				Bay Count			1987-01		
layport Bi Bayport	unicipal	Alan King			Bayport,			0001-01	23,140	
erlin Bio Berlin	Landfill Gas	- King			Green Lak			2001-01	10 583	
erlin Gor Berlin Gor					Coos Coun			1948-01		
ieber Pla Bieber Plant					Bieber, Ca			0001-01	15,020.23	
liodyne B Biodyne B					Will Count			2001-01	12 536 1	
Control of the property of the	Eurium das			base Queries				2001 01	12,550.1	





Bion	nass Plants (Small Set)	Governors and Parties in US States	+ Compose Query	Query Result
998960.facility_name	998960.generatingcapaci	ty 998555.governor	998555.party	Confidence
Atlantic City Landfi Biomass	1.44	Gov. James E. McGreevey	Democrat	Certain
Chicopee II LFG Biomass Fa	ci 5.42	Gov. Mitt Romney	Republican	Certain
Central Minn. Ethano Bioma	s 0.95	Gov. Tim Pawlenty	Republican	Certain
Central LF Biomass Facility	2.375	Gov. Don Carcieri	Republican	Certain
Dairyland PPA Landfi Bioma	is 2.85	Gov. Tim Pawlenty	Republican	Certain
Blue Spruce Farm Ana Bion	a 0.257	Gov. Jim Douglas	Republican	Certain
APS Biomass I Biomass Fac	lity 2.85	Gov. Janet Napolitano	Democrat	Certain
Chicopee II LFG Biomass Fa	ci 5.42	Gov. Jane Maria Swift	Republican	Certain
Crapo Hill Landfill Biomass	F 3.04	Gov. Mitt Romney	Republican	Certain
Coventry LFG Biomass Facil	ity 4.56	Gov. Jim Douglas	Republican	Certain
Atlantic City Landfi Biomass	1.44	Gov. Richard J. Codev	Democrat	Certain
Brickyard Energy Partners		Gov. Howard Dean M.D.	Democrat	High Confidence
Brickvard Recycling Biomas	s 0.19	Gov. Jim Douglas	Republican	High Confidence
Altamont Gas Recovery Bio	m 2.6	Gov. Donald Eugene Siegelm	Democrat	High Confidence
Covanta Marion Inc. Biomas	s 11.5	Gov. Bruce Edward Babbitt	Democrat 0.73254585	High Confidence
Covanta Marion Inc. Biomas	s 11.5	Gov. Joseph Edward Brennan	Democrat 0.73234383	High Confidence
American Ref-Fuel of Delay	v 80	Gov. Michael Newbold Castle	Republican	High Confidence
Biodyne Peoria Biomass Fa	ility 4	Gov. Zell Miller	Democrat	High Confidence
Brent Run Generating Statio	n 2.4	Gov. Don Sundquist	Republican	High Confidence
C & C Electric Biomass Faci	ity 2.7	Gov. Pete Wilson	Republican	High Confidence
Arbor Hills Biomass Facility	19	Gov. Michael Lowry	Democrat	High Confidence
Covanta Marion Inc. Biomas	s 11.5	Gov. Harry Roe Hughes	Democrat	High Confidence
Coyote Canyon Steam Plant	17	Gov. Neil Goldschmidt	Democrat	High Confidence
Al Turi Biomass Facility	4.4	Gov. Neil Goldschmidt	Democrat	High Confidence
Altamont Gas Recovery Bio	m 2.6	Gov. Frank H. Murkowski	Republican	High Confidence
Altamont Gas Recovery Bio	m 2.6	Gov. Tony Knowles	Democrat	High Confidence
Al Turi Biomass Facility	4.4	Gov. Mario Matthew Cuomo	Democrat	High Confidence
Alabama Pine Pulp Biomas	i 32.085	Gov. Harold Guy Hunt	Republican	High Confidence
Blackburn Landfill Co-Gene	r 2.9	Gov. James B. Hunt Jr.	Democrat	High Confidence
American Ref-Fuel of Essex	60	Gov. Thomas H. Kean	Republican	High Confidence
American Ref-Fuel of Essex	c 60	Gov. Jim Florio	Democrat	High Confidence
Colville Indian Power & Ven	e 12.5	Gov. Francis Anthony Keating	Republican	High Confidence
Alabama Pine Pulp Biomas	32.085	Gov. Barbara Roberts	Democrat	High Confidence
Alabama Pine Pulp Biomas	32.085	Gov. Neil Goldschmidt	Democrat	High Confidence
Albany Landfill Gas Utilizati	o 1.8	Gov. George E. Pataki	Republican	High Confidence
Charlotte Motor Speedway	Bi 4.3	Gov. James B. Hunt Jr.	Democrat	High Confidence
Century Flooring Co Biomas	s 1.7	Gov. John Carlin	Democrat	High Confidence
APS Riomass I Riomass Faci	litv 2.85	Gov Michael F Fasley	Democrat	High Confidence
	Sho	ow Database Query for Confidence	Computation	

Democratic Biomass Capacity \geq Republican Biomass Capacity $17 + 5 + 9 \geq 8 + 14 + 2$

Democratic Biomass Capacity \geq Republican Biomass Capacity $\Phi = [x_1 \otimes 17 + x_2 \otimes 5 + x_3 \otimes 9 \geq x_4 \otimes 8 + x_5 \otimes 14 + x_6 \otimes 2]$

Democratic Biomass Capacity
$$\geq$$
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■ Assume *x_i* are Boolean random variables

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- Then the sum expression $\alpha = x_1 \otimes 17 + x_2 \otimes 5 + x_3 \otimes 9$ is a \mathbb{N} -valued random variable

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- Hence Φ is a \mathbb{B} -valued random variable

Democratic Biomass Capacity
$$\geq$$
 Republican Biomass Capacity
$$\Phi = [x_1 \otimes 17 + x_2 \otimes 5 + x_3 \otimes 9 > x_4 \otimes 8 + x_5 \otimes 14 + x_6 \otimes 2]$$

- Assume x_i are Boolean random variables
- Then the sum expression $\alpha = x_1 \otimes 17 + x_2 \otimes 5 + x_3 \otimes 9$ is a N-valued random variable
- Hence Φ is a \mathbb{B} -valued random variable
- $P_{\Phi}[\top]$ is the probability that a random choice of possible values for the variables x_i satisfies the inequality
- In this example, $P_{\Phi}[\top]$ is the probability that Democrats are responsible for more biomass capacity than Republicans

Outline

What do we mean by + in $\Phi_1 \otimes 17 + \Phi_2 \otimes 5?$ Well, it depends . . .

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Aggregation modelled by commutative monoids

- Carrier M, e.g. \mathbb{N} or \mathbb{R}
- Binary operation $M \times M \rightarrow M$
- Neutral element $0 \in M$
- Examples for aggregation monoids: SUM $(\mathbb{N},+,0)$, MIN (\mathbb{N},\min,∞) , MAX $(\mathbb{N},\max,-\infty)$, PROD, COUNT (special case of SUM)

What are Φ_1, Φ_2 in $\Phi_1 \otimes 17 + \Phi_2 \otimes 5$?

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$$\Phi_1 \otimes 17 + \Phi_2 \otimes 5^\circ$$

■ Consider Query:

$$AGG_B[(R \cup S) \bowtie_A T]$$

R	S	T
А Ф	А Ф	А В Ф
1 x ₁	1 y ₁	1 17 z ₁
2 x ₂		2 5 z ₂

What are
$$\Phi_1, \Phi_2$$
 in $\Phi_1 \otimes 17 + \Phi_2 \otimes 5$?

$$\Phi_1 \otimes 17 + \Phi_2 \otimes 57$$

Consider Query:

$$AGG_B[(R \cup S) \bowtie_A T]$$

Tuples annotations modelled by semirings

■
$$(R \cup S) \bowtie_A T$$
 yields

$$\begin{array}{c|cccc} (R \cup S) \bowtie_A T \\ \hline A & B & \Phi \\ \hline 1 & 17 & (x_1 + y_1) \cdot z_1 \\ 2 & 5 & x_2 \cdot z_2 \\ \end{array}$$

Aggregation on top of this table yields:

$$((x_1 + y_1) \cdot z_1) \otimes 17 + (x_2 \cdot z_2) \otimes 5$$

where the meaning of + depends on the aggregation monoid

Semimodule

- Algebraic framework introduced by Amsterdamer et al. [2011]
- The algebraic structure combining semirings and monoids is called semimodule
- Generalisation of vector space. "Scalars": tuple annotations, "Vectors": aggregation values
- Semimodule expressions represent data values conditioned on tuple annotations

Semiring and semimodule expressions are random variables

- Semimodule: Random variable over aggregation domain
- Semiring expressions: ?
 - So far in probabilistic databases: Boolean random variable
 - However: B is in general not large enough for aggregation; need larger semiring, for example natural numbers

Aggregation Needs Semirings Larger Than B

Pro	ducerEU	ProducerUS			Product	S
Α	ф	Α	Ф	Α	Price	Φ
1	<i>x</i> ₁	1	<i>y</i> ₁	1	17	<i>z</i> ₁
2	<i>x</i> ₂			2	5	<i>z</i> ₂

- Query: SUM_{Price} [(ProducerEU ∪ ProducerUS) ⋈_A Products] asking for total price of products sold by all producers
- Resulting expression: $((x_1 + y_1) \cdot z_1) \otimes 17 + (x_2 \cdot z_2) \otimes 5$
- Valuation $\nu: x_1, x_2, y_1, z_1, z_2 \mapsto \top$ yields $\top \otimes 17 + \top \otimes 5 = 22$ Arguably not the expected result

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- Valuation $\nu: x_1, x_2, y_1, z_1, z_2 \mapsto \top$ yields $\top \otimes 17 + \top \otimes 5 = 22$ Arguably not the expected result
- Boolean semiring is not large enough for SUM
- Better choice: Semiring \mathbb{N} . Identify $\bot \sim 0$, $\top \sim 1$.
- Valuation $\nu: x_1, x_2, y_1, z_1, z_2 \mapsto 1$ yields $((1+1)\cdot 1)\otimes 17 + (1\cdot 1)\otimes 5 = 2\otimes 17 + 1\otimes 5 = 39.$

A More Formal View: **Expressions, Random Variables**

■ The probability space *induced by* **X** has as samples the set of valuations from **X** to *S*,

$$\Omega = \{\nu : \mathbf{X} \to \mathcal{S}\}$$

■ Every expression $\Phi \in K$ is an S-valued random variable over Ω with probability distribution

$$P_{\Phi}[s] = P(\{
u \in \Omega \mid
u(\Phi) = s\}) = \sum_{\substack{
u \in \Omega: \\
u(\Phi) = s}} P(
u)$$

for every $s \in S$

Outline

The pvc-tables Representation System

Ingredients for pvc-tables

- A set X of variable symbols
- Tuples contain constants or semimodule expressions over X
- Every tuple is annotated with a semiring expression over X

Queries

- \blacksquare Query Q maps pvc-table database D to pvc-table Q(D)
- Annotations are propagated via query operators
- Expressions concisely encode probability distributions of answers

Properties of pvc-tables

- Polynomial overhead (Amsterdamer et al. [2011]): $|Q(D)| \in \mathcal{O}\big(\text{poly}(|D|)\big) \qquad \qquad \text{(unlike pc-tables)}$
- Completeness: Every finite probability distribution over relations (with set or bag semantics) can be represented by pvc-tables

The pvc-tables Representation System

Semantics: Set vs Bag & Deterministic vs Probabilistic Different choices for the semiring and the probability distributions of the annotation variables give rise to different database semantics.

Database Sem	antics	Semiring	Probability Distributions
Deterministic	Set	\mathbb{B}	$P_x[\top] = 1 \text{ or } P_x[\bot] = 1$
Deterministic	Bag	\mathbb{N}	$\exists n \in \mathbb{N} : P_x[n] = 1$
Probabilistic	Set	\mathbb{B}	$P_x[\top], P_x[\bot] \in [0,1]$
Probabilistic	Bag	N	$\forall n \in \mathbb{N}: P_x[n] \in [0,1]$

Outline

Query Evaluation in pvc-tables (1)

Step 1: Construction of Expressions

Alongside (standard) query evaluation, compute annotations.

- Project, Union, Cartesian Product: Construction of semiring expressions (· for joint, and + for alternative use of data)
- Aggregation (with grouping): Construct semimodule expressions $(\sum_{AGG} \Phi \otimes \nu)$

R		
АВ	Φ	
a 1	<i>x</i> ₁	select AGG(B) from R group by
a 2	x ₂	
b 3	<i>x</i> ₃	
b 4	<i>x</i> ₄	

	pvc-table	
Α	AGG(B)	Ф
а	$x_1 \otimes 1 + x_2 \otimes 2$	$[x_1 + x_2 \neq 0]$
b	$x_3 \otimes 3 + x_4 \otimes 4$	$[x_3+x_4\neq 0]$

Query Evaluation in pvc-tables (1)

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	п		_
Α	В	Ф	-
а	1	<i>x</i> ₁	select AGG(B) from R gro
а	2	x_2	
b	3	<i>x</i> ₃	
b	4	<i>x</i> ₄	

pc-table		
Α	SUM(B)	Φ
а	0	$\bar{x}_1 \cdot \bar{x}_2$
а	1	$x_1 \cdot \bar{x}_2$
а	2	$\bar{x}_1 \cdot x_2$
а	3	$x_1 \cdot x_2$
b	0	$\bar{x}_3 \cdot \bar{x}_4$

Exponential overhead! Lechtenbörger et al. [2002]

Query Evaluation in pvc-tables (2)

Step 2: Probability Computation

Problem: Given a tuple, compute its probability distribution.

Idea: Tuple probability is equivalent to joint probability distribution of its semimodule expressions and annotation expression as obtained from evaluation step 1.

Approach: Compile expressions into a tractable form consisting of *independent* and *mutually exclusive* sub-expressions.

Compilation: Independent Decomposition

Consider semiring expression $\Phi = x + y$. Since x, y are independent random variables, the probability distribution of Φ is given by the convolution of x and y.

If
$$x, y$$
 are in \mathbb{N} : $P_{x+y}[n] = \sum_{\substack{i,j \in \mathbb{N} \\ i+j=n}} P_x[i]P_y[j]$

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If
$$x,y$$
 are Boolean:
$$P_{x+y}[\bot] = \sum_{\substack{a,b \in \{\bot,\top\}\\a\lor b=\bot}} P_x[a]P_y[b]$$

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$$P_{x+y}[\bot] = \sum_{\substack{a,b \in \{\bot,\top\}\\a\lor b = \bot}} P_x[a]P_y[b] = P_x[\bot]P_y[\bot]$$

$$P_{x+y}[\top] = \sum_{\substack{a,b \in \{\bot,\top\}\\a\lor b = \top}} P_x[a]P_y[b]$$

$$= P_x[\top]P_y[\top] + P_x[\bot]P_y[\top] + P_x[\top]P_y[\bot]$$

$$= 1 - P_x[\bot]P_y[\bot]$$

Compilation: Independent Decomposition

The applicability of convolution is not limited to "sums"; convolution is equally well defined for other binary operations:

Convolution for algebraic operations

- Semiring expressions: $\Phi \cdot \Psi$, $\Phi + \Psi$
- Semimodule expressions: $\alpha + \beta$
- Mixed semiring and semimodule expressions: $\Phi \otimes \alpha$
- \blacksquare Convolution is also applicable to comparisons of expressions, such as $\alpha \leq \beta$

Compilation: Mutually Exclusive Expressions

What if there are no independent sub-expressions? Example: $\alpha = a(b+c) \otimes 10 + c \otimes 20$

Idea: Instantiate one of the variables to create mutually exclusive sub-expressions.

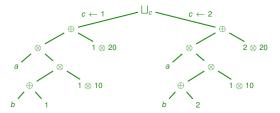
$$P(\alpha) = P_{c}[1] \cdot P(a(b+1) \otimes 10 + 1 \otimes 20) + P_{c}[2] \cdot P(a(b+2) \otimes 10 + 2 \otimes 20) + P_{c}[3] \cdot P(a(b+3) \otimes 10 + 3 \otimes 20) + \dots$$

Need to consider all possible values of c with non-zero probability. In particular: For Boolean variables, the above construction yields Shannon's expansion.

Decomposition Trees (d-trees)

Decomposition gives rise to a tree whose nodes explain the decomposition steps taken. For example, \bigsqcup for mutex decomposition, \oplus for convolution w.r.t. +, \otimes for convolution w.r.t. \otimes , etc.

Example: $\alpha = a(b+c) \otimes 10 + c \otimes 20$



Tractable Probability Computation for d-trees

The probability distribution P_d of a d-tree d whose nodes have probability distributions p_1, \ldots, p_n can be computed in time $\mathcal{O}(\prod |p_i|)$.

Specific polynomial time cases

- For MIN and MAX monoids combined with any semiring
- For SUM monoid: If monoid values and size of probability distributions of semiring expressions are bounded by constants
 - This subsumes COUNT aggregation

Further Applications of d-trees

- Approximate probability computation by partial expansion of d-tree (Olteanu et al. [2010], Fink et al. [2011])
- Sensitivity analysis and explanation of query results (Kanagal et al. [2011])
- Conditioning probabilistic databases (Koch and Olteanu [2008])

Tractability for query evaluation on probabilistic databases is considered with respect to data complexity:

For which class of queries can probability distributions of query answers be computed in *polynomial-time data complexity* for any tuple-independent database?

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For which class of queries can probability distributions of query answers be computed in *polynomial-time data complexity* for any tuple-independent database?

- Syntactic characterisation of tractable queries with aggregates
 - There are known classes of tractable non-aggregate queries with polynomial-time d-tree compilation, e.g. hierarchical queries
 - Extend these classes by adding nested aggregation without breaking the tractable (e.g. hierarchical) property

Example 1

select R.A from R where R.B =
$$\left(\text{select MIN}(S.B) \text{ from S}\right)$$

where S.C = R.C

Tractable sub-queries without aggregation: select S.B from S where S.C = B.C

Example 2

select 1 where

$$\left(\text{select MIN(R.A) from R} \right) <= \quad \left(\text{select COUNT(*) from S,T} \right. \\ \left. \text{where S.A=T.A} \right)$$

Tractable sub-queries without aggregation:

select 1 where (select R.A from R) select 1 from S.T where S.A=T.A

select 1 where (select R.A from R) <= (select 1 from S,T where S.A=T.A)

Outline

Performance Analysis

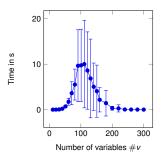


Figure: Varying the number of variables for a randomly generated semimodule expression (L=90, #cl=2, #l=2, maxv=5, c=3, #runs=40, AGGL=MIN)

$$\left[\sum_{\mathsf{AGGL}}^{\mathit{L}} \Phi_{\mathit{i}} \otimes \mathit{v}_{\mathit{i}} \; = \; \mathit{c}\right]$$

Performance Analysis

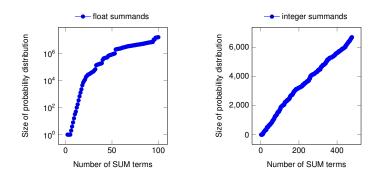


Figure: Size of the probability distributions for SUM semimodule expressions of varying size. When summing float numbers from a fixed range, the size of the probability distribution grows potentially exponentially in the number of terms, while summing integers from a fixed range it grows linearly.

Performance Analysis

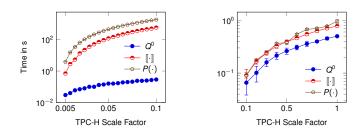


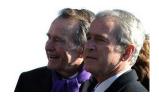
Figure: TPC-H Queries Q1 (modified) and Q2. For each query, the graphs compare the execution times (1) on a deterministic database (Q^0) without expression or probability computation, (2) of the computation of the expressions ($[\![\cdot]\!]$), and (3) of probability computation for the result tuples $(P(\cdot))$).





Who is responsible for a larger capacity of biogas plants, Democrats or Republicans?



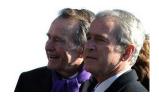






Who is responsible for a larger capacity of biogas plants, Democrats or Republicans?











Who is responsible for a larger capacity of biogas plants, Democrats or Republicans?





End. ?

Definitions

Monoid

A monoid is a set M with an operation $+: M \times M \to M$ and a neutral element $0 \in M$ that satisfy the following axioms for all $m_1, m_2, m_3 \in M$:

$$(m_1 + m_2) + m_3 = m_1 + (m_2 + m_3)$$

 $0 + m_1 = m_1 + 0 = m_1$

A monoid is commutative if $m_1 + m_2 = m_2 + m_1$

Semiring

A commutative semiring is a set S together with operations $+, \cdot : S \times S \to S$ and neutral elements $0, 1 \in S$ such that (S, +, 0) and $(S, \cdot, 1)$ are commutative monoids and the following holds for all $s_1, s_2, s_3 \in S$:

$$s_1 \cdot (s_2 + s_3) = (s_1 \cdot s_2) + (s_1 \cdot s_3)$$
$$(s_1 + s_2) \cdot s_3 = (s_1 \cdot s_3) + (s_2 \cdot s_3)$$
$$0 \cdot s_1 = s_1 \cdot 0 = 0$$

Definitions

Semimodule

Let $(S, +_S, 0_S, \cdot_S, 1_S)$ be a commutative semiring. As S-semimodule M consists of a commutative monoid $(M, +_M, 0_M)$ and a binary operation $\otimes : S \times M \to M$ such that for all $s_1, s_2 \in S$ and $m_1, m_2 \in M$ we have

$$s_{1} \otimes (m_{1} +_{M} m_{2}) = s_{1} \otimes m_{1} +_{M} s_{1} \otimes m_{2}$$

$$(s_{1} +_{S} s_{2}) \otimes m_{1} = s_{1} \otimes m_{1} +_{M} s_{2} \otimes m_{1}$$

$$(s_{1} \cdot_{S} s_{2}) \otimes m_{1} = s_{1} \otimes (s_{2} \otimes m_{1})$$

$$s_{1} \otimes 0_{M} = 0_{K} \otimes m_{1} = 0_{M}$$

$$1_{S} \otimes m_{1} = m_{1}$$

Further Experiments

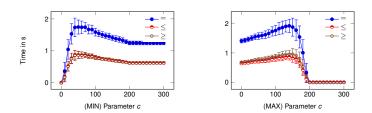
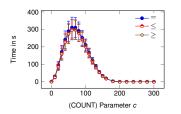


Figure: Experiment A: Varying the constant c for different aggregation monoids and comparison operators θ . #v=25, L=200, R=0, #cl=3, #l=3, maxv=200.

Further Experiments



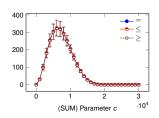


Figure: Experiment A: Varying the constant c for different aggregation monoids and comparison operators θ . #v=25, L=200, R=0, #cl=3, #l=3, maxv=200.

References