

Probability and Computing  
Hilary Term 2018  
Exercise Sheet 5

1. Consider a random walk on a *directed* graph with vertices  $\{0, 1, \dots, n\}$ , and directed edges  $\{(i, 0), (i, i+1) : i < n\} \cup \{(n, 0), (n, n)\}$ . At vertex  $i$  each edge is taken with probability  $1/d(i)$  where  $d(i)$  is the outdegree of  $i$ . Find the stationary distribution of the corresponding Markov chain.
2. A cat and a mouse each independently take a random walk on a connected, undirected, non-bipartite graph  $G = (V, E)$  with  $m$  edges and  $n$  vertices. They start at the same time on different nodes, and each makes one transition in each time step. The cat eats the mouse if they are ever at the same node at some time step.
  - (a) Suppose that the cat starts at vertex  $u$  and the mouse starts at vertex  $v$ . Show that there is an even-length walk from  $u$  to  $v$  of length at most  $2n$ .
  - (b) Show an upper bound of  $O(m^2n)$  on the expected time before the cat eats the mouse. Carefully justify your answer.
3. Consider a random walk on an undirected tree  $T(V, E)$ . Show that for each edge  $\{u, v\}$ ,  $h_{uv} = 2k - 1$ , where  $k$  is the number of nodes of the maximum-size subtree that contains  $u$  but not  $v$  (i.e., the number of nodes of the connected component of  $u$  that results after removing edge  $\{u, v\}$ ). Conclude that for every edge  $\{u, v\}$ ,  $h_{uv} + h_{vu} = 2|E|$ .

More generally, for a random walk on a connected graph  $G(V, E)$ , show that

$$h_{uv} + h_{vu} = 2|E|,$$

when edge  $\{u, v\}$  is a bridge. An edge  $\{u, v\}$  is called a *bridge* if removing it disconnects the graph.

4. Suppose that we are given  $n$  records,  $R_1, R_2, \dots, R_n$ . The records are kept in some order. We use the *move-to-front heuristic*: after accessing  $R_j$  we move it to the front of list. Assume that the cost of accessing the  $j$ -th record in the order, including move-to-front, is  $j$ . For example, if we access  $R_1$  in the list  $R_2, R_4, R_3, R_1$  the total cost is 4 and the resulting list is  $R_1, R_2, R_4, R_3$ .

Suppose that at each step record  $R_i$  is accessed with probability  $p_i > 0$ , where  $\sum_{i=1}^n p_i = 1$ . We can then think of the various orders of the records as states of a Markov chain. Give the stationary distribution of this chain. Assuming that the chain is in its stationary distribution, give an expression for the expected cost per access. Your expression for the expectation should be easily computable in time that is polynomial in  $n$ , given the  $p_i$ .

5. Give an FPRAS for the variant of DNF counting in which we are given a DNF formula and an integer  $k$  and we want to estimate the number of satisfying assignments in which exactly  $k$  of the variables are true.