LECTURE 8: SETS
LECTORE 6. SETS
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### 2 Defining Sets — Enumeration

- By convention, a set can be defined by enumerating its components in curly brackets.
- EXAMPLE.

 $Vowels == \{a, e, i, o, u\}$ 

EXAMPLE.

 $Weekend == \{Sat, Sun\}$ 

- The double-equals symbol (==) is read 'is defined to be'; it gives us a method for naming sets, so that we can subsequently use them.
- (We can't use sets unless they have been previously defined in some way.)

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#### 1 What is a Set?

- The concept of a set is used throughout mathematics; its formal definition matches closely our intuitive understanding of the word.
- **Definition**: A set is an unordered collection of distinct objects.
- We can build sets containing any objects that we like, but usually we consider sets whose elements have some property in common.
- The objects comprising a set are called its elements or members.
- Note that:

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- sets do not contain *duplicates*;
- the order of elements in a set is *not* significant.

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- In Z, another equivalent method for defining sets is allowed ...
- EXAMPLE.

- Another standard convention is to allow number ranges ...
- EXAMPLE.

Days == 1 ... 365

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• Some sets occur so often that they have been pre-defined and given standard names.

 $\mathbb{N}$  — the natural numbers;

$$IN = \{0, 1, 2, 3, \ldots\}$$

 $\mathbb{N}_1$  — the natural numbers greater than 0;

$$\mathbb{N}_1 = \{1, 2, 3, \ldots\}$$

Z — the *integers*;

$$Z = \{\ldots, -2, -1, 0, 1, 2, \ldots\}$$

**IB** — the *truth values*:

$$IB = \{true, false\}$$

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4

- Note that the r.h.s. of the expression  $x \in S$  must be an expression which returns a set of values of the same type as the expression on the l.h.s returns.
- EXAMPLE. This isn't allowed:

$$\{1, 2, 3\} \in \{4, 5, 6\}$$

But this is:

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$$\{1, 2, 3\} \in \{\{6, 7, 8\}, \{1, 2, 3\}\}\$$

and so is this:

$$\{1, 2, 3\} \in \{\{6, 2, 1\}\}\$$

The first legal expression above evaluates to *true*; the second evaluates to *false*.

• Note that Z differs in this respect from what might be called *naive set theory*: Z is said to be a *typed* set theory, (because we must check the types on the l.h.s and r.h.s. of our expressions.)

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6

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### 3 Relational Operators for Sets

• The set theory in Z provides us with a variety of *relational operators* for expressing the properties of sets.

# 3.1 Set Membership

• To express the fact that an object *x* is a *member* of set *S*, we write

$$x \in S$$

Note that an expression with this form is either *true* or *false*; the thing is either a member or not.

- EXAMPLES.
  - $-a \in \{a, e, i, o, u\}$
  - $-1 \in \mathbb{N}$
  - $-23 \in Days$

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# 3.2 Set Equality

• **Definition**: Let *S* and *T* be arbitrary sets of the same type. Then

$$S = T$$

is *true* iff *S* and *T* contain precisely the same members.

- Note that the *order* in which values occur in the set is not significant.
- EXAMPLES. The following is a true expression:

$$\{1, 2, 3\} = \{2, 3, 1\}$$

as is this

$$Days = 1..365$$

however, the following is not:

$$Days = \mathbb{N}$$
.

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#### 3.3 Subsets

• **Definition:** Let *S* and *T* be arbitrary sets of the same type. Then

$$S \subseteq T$$

is true iff every member of *S* is also a member of *y*.

- EXAMPLES. The following are true expressions:
  - $-\{1,2,3\} \subseteq \mathbb{N};$
  - {true} ⊆  $\mathbb{B}$ ;
  - **N** ⊆ Z;
  - $-\mathbb{N}_1\subseteq\mathbb{N};$
  - $-\{1,2,3\} \subseteq \{1,2,3\}.$

whereas the following are not:

- $-\{1,2,3,4\} \subseteq \{1,2,3\};$
- -Z ⊆  $\mathbb{N}$ .
- Note that equal sets are subsets of each other; we can express this fact in the following theorem about sets:

$$(S = T) \Rightarrow (S \subseteq T)$$

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8

### 4 The Empty Set

• There is a special set that has the property of having no members; this set is called the *empty set*, and is denoted either

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or (more usually)

Ø

• Note that if *x* is *any* value, then the expression

$$x \in \emptyset$$

must be false. (Nothing is a member of the empty set.)

• Similarly, if *S* is an arbitrary set, then

$$\emptyset \subseteq S$$

must be true; this includes of course

$$\emptyset \subseteq \emptyset$$

(The emptyset is a subset of every set.)

• But note that this is *not* a theorem:

$$\emptyset \subset S$$

Why?

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10

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# 3.4 Proper Subsets

- As we just saw, if two sets are equal, then they are subsets of each other. This may seem strange! We introduce another relational operator that more accurately captures our intuitive understanding of what 'subset' means.
- **Definition**: Let *S* and *T* be arbitrary sets of the same type. Then:

$$S \subset T$$

is true iff both  $S \subseteq T$  and  $S \neq T$ .

- If  $S \subset T$  then then S is said to be a *proper subset* of T.
- Note that we can formally define the ⊂ relation using concepts introduced earlier:

$$S \subset T \Leftrightarrow (S \subseteq T) \land (S \neq T)$$

• *Note the use of*  $\Leftrightarrow$  *to introduce definitions.* 

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## 5 The Powerset Operator

• **Definition:** Let *S* be an arbitrary set. Then the set of all subsets of *S* is given by

$$\mathbb{P} S$$
.

• EXAMPLES.

$$S = \{1\}$$

then

$$\mathbb{P} S = \{\emptyset, \{1\}\}\$$

– If

$$S = \{1, 2\}$$

then

$$\mathbf{IP} S = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}\$$

– If

$$S = \{a, b, c\}$$

then

$$\mathbb{P} S = ?$$

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$$T \in \mathbb{P} S \Leftrightarrow (T \subseteq S)$$

- Note that if *S* has *n* members, then  $\mathbb{P} S$  has  $2^n$  members.
- For this reason, the powerset of a set *S* is sometimes denoted by 2<sup>*S*</sup>. (But not by us.)

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12

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### 6.2 Set Intersection

- The *intersection* of two sets is a third set that contains only elements common to both
- **Definition**: Let *S* and *T* be arbitrary sets. Then

$$x \in (S \cap T)$$

iff both  $x \in S$  and  $x \in T$ .

• As a derived operator:

$$x \in (S \cap T) \Leftrightarrow (x \in S) \land (x \in T).$$

• EXAMPLES.

$$-\left\{ a,e,i\right\} \cap\left\{ o,u\right\} =\emptyset;$$

$$-a \neg \in \{a, e, i\} \cap \{e, i\};$$

$$-\mathbb{N}=\mathbb{N}\cap Z$$
.

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14

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# 6 Manipulating Sets

#### 6.1 Set Union

- The *union* of two sets is a third set that contains the members of both.
- ullet The symbol for union is  $\cup$ .
- **Definition:** Let *S* and *T* be arbitrary sets. Then

$$x \in (S \cup T)$$

is true iff either  $x \in S$  or  $x \in T$ .

• Note that we can introduce this as a derived operator:

$$x \in S \cup T \Leftrightarrow x \in S \lor x \in T$$
.

• EXAMPLES.

$$-\{a,e,i\} \cup \{o,u\} = \{a,e,i,o,u\};$$

$$-\{a,e\} \cup \{e,i\} = \{a,e,i\};$$

$$-\emptyset \cup \{a,e\} = \{a,e\}.$$

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#### 6.3 Set Difference

- The notation S \ T denotes the set obtained from S by removing from it all the elements of T tha occur in it; in other words, you take T away from S.
- **Definition:** Let *S* and *T* be arbitrary sets. Then

$$x \in (S \setminus T)$$

is true iff  $x \in S$  and  $x \notin T$ .

• As a derived operator:

. .

• EXAMPLES.

$$-\{a, e, i, o, u\} \setminus \{o\} = \{a, e, i, u\}$$

 $-\emptyset\setminus\{1,2,3\}=\emptyset.$ 

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#### 6.4 Theorems about Sets

Let S, T, and U be arbitrary sets of the same type. Then:

$$S \cup S = S$$

$$S \cap S = S$$

 $(S \cup T) \cup U = S \cup (T \cup U)$ 

 $(S\cap T)\cap U \,=\, S\cap (T\cap U)$ 

$$S \cup T = T \cup S$$

$$S \cap T = T \cap S$$

$$S \cup (T \cap U) = (S \cup T) \cap (S \cup U)$$

$$S \cap (T \cup U) = (S \cap T) \cup (S \cap U)$$

$$S \cup \emptyset = S$$

$$S \cap \emptyset = \emptyset$$

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16

## 7 Defining Sets: Comprehension

- We have seen how sets may be defined by *numeration* listing their contents.
- This technique is impractical for large sets!
- *Comprehension* is a way of defining sets:
  - in terms of their properties;
  - in terms of other sets.
- Comprehension is best explained via examples:
- 1.  $\{n : \mathbb{N} \mid (n \ge 0) \land (n \le 3)\} = \{0, 1, 2, 3\}$
- 2.  $\{n : Z \mid (n > 1) \land (n < 1)\} = \emptyset$
- 3.  $\{n: \mathbb{N} \mid (n \ge 0) \land (n \le 3) \bullet n + 1\} = \{1, 2, 3, 4\}$

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18

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# 6.5 Cardinality

- The *cardinality* of a set is the number of elements in it.
- **Definition**: If *S* is an arbitrary set, then the *cardinality* of *S* is denoted #*S*.
- Note that for an arbitrary set *S*:

$$\#S \in \mathbb{N}$$
.

• EXAMPLES.

$$-\#\{1,2,3\}=3;$$

$$-\#\emptyset = 0;$$

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$$-\#(\{1,2,3\}\cap\{2\})=1.$$

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- In general, a set comprehension expression will have three parts:
  - a signature specifies the base sets from which values are extracted;
  - a predicate the defining property of the set;
  - a term specifies how the actual values in the set are computed.
- Generally, we omit the term part we just have a signature and a predicate, as in examples (1) and (2), above.

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