

QUANTUM MEASUREMENTS FROM A LOGICAL POINT OF VIEW

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PRESENTATION OUTLINE

A Logic about Measurement

Models

Two Theorems on Models

Sequences of Outcomes

A LOGIC ABOUT MEASUREMENT

Logical Constructor

$$\text{Mes}(s, \mathcal{O}, p, t)$$

- s - Measured system
- \mathcal{O} - Measured **sharp** observable
- $p \in \mathcal{O}$ - Outcome
- t - Resulting system

$$\neg(\exists s, \mathcal{O}, t: \text{Mes}(s, \mathcal{O}, \perp, t))$$

$$\forall p \neq \perp, \exists s, \mathcal{O}, t: \text{Mes}(s, \mathcal{O}, p, t)$$

$$\forall s, \mathcal{O}, \exists p, t: \text{Mes}(s, \mathcal{O}, p, t)$$

Definition

$$s \blacktriangleright p \stackrel{\Delta}{\iff} \neg(\exists \mathcal{O}, t: \text{Mes}(s, \mathcal{O}, p^\perp, t))$$

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Meaning in a Hilbert Space

If a system s is in a state $|\varphi\rangle$,

then $s \blacktriangleright p$ corresponds to $|\varphi\rangle \in p$, or $\Pi_p|\varphi\rangle = |\varphi\rangle$.

$$s \blacktriangleright \top$$

$$\neg(s \blacktriangleright \perp)$$

$$\forall p \neq \perp, \exists s: s \blacktriangleright p$$

Claim

The **certainty/impossibility** of an outcome is **independent** of the measured observable.

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Motivation: the Born Rule

The probability of obtaining outcome P in state $|\varphi\rangle$ is $\langle\varphi|\Pi_P|\varphi\rangle$.

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If $s \blacktriangleright p$ and $p \leq q$, then $s \blacktriangleright q$.

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Justification

$$\{p, p^\perp \wedge q, q^\perp\}$$

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Justification

p and q are compatible iff Π_p and Π_q commute

If $\Pi_p|\varphi\rangle = |\varphi\rangle$, then $\Pi_p(\Pi_q|\varphi\rangle) = \Pi_q|\varphi\rangle$.

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$$\forall a, b \in L, \quad a \& b \triangleq b \wedge (a \vee b^\perp)$$

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“Simplifications”

Weak Noncontextuality 2 $s \blacktriangleright p$ and $s \blacktriangleright q \implies s \blacktriangleright p \& q$

Compatible Preservation $s \blacktriangleright p$ and $\text{Mes}(s, \mathcal{O}, q, t) \implies t \blacktriangleright p \& q$

Axioms of \mathcal{T}_L

$$s \blacktriangleright \top$$

$$\neg(s \blacktriangleright \perp)$$

$$p \neq \top \implies \exists s: s \blacktriangleright p$$

$$s \blacktriangleright p \text{ and } p \leq q \implies s \blacktriangleright q$$

$$s \blacktriangleright p \text{ and } s \blacktriangleright q \implies s \blacktriangleright p \& q$$

$$s \blacktriangleright p \text{ and } \text{Mes}(s, \mathcal{O}, q, t) \implies t \blacktriangleright p \& q$$

MODELS

Reminder

$$\forall s, p, \quad s \blacktriangleright p \stackrel{\Delta}{\iff} \forall t, \neg(\exists \mathcal{O}: \text{Mes}(s, \mathcal{O}, p^\perp, t))$$

$$\forall s, p, q, t, \quad s \blacktriangleright p \text{ and } (\exists \mathcal{O}: \text{Mes}(s, \mathcal{O}, q, t)) \implies t \blacktriangleright p \& q$$

Definition

A **model** of \mathcal{T}_L is a pair $\mathfrak{G} = (A, M)$ consisting of

- a set A
- a relation M on $A \times L \times A$

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$\exists \mathcal{O} : \text{Mes}(s, \mathcal{O}, p, t)$ translates as $M(a, p, b)$

$$a \blacktriangleright_{\mathfrak{G}} p \stackrel{\Delta}{\iff} \neg(\exists b \in A : M(a, p^\perp, b))$$

$$\forall a \in A, a \blacktriangleright_{\mathcal{G}} \top$$

$$\forall a \in A, \neg(a \blacktriangleright_{\mathcal{G}} \perp)$$

$$\forall p \in L, p \neq \top \implies \exists a \in A: a \blacktriangleright_{\mathcal{G}} p$$

$$\forall a \in A, \forall p, q \in L, a \blacktriangleright_{\mathcal{G}} p \text{ and } p \leq q \implies a \blacktriangleright_{\mathcal{G}} q$$

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$$\forall a, b \in A, \forall p, q \in L, a \blacktriangleright_{\mathcal{G}} p \text{ and } M(a, q, b) \implies b \blacktriangleright_{\mathcal{G}} p \& q$$

Definition

Given a Hilbert space \mathcal{H} , we define the model $\mathfrak{H}_{\mathcal{H}} = (A_{\mathfrak{H}}, M_{\mathfrak{H}})$ by

$$A_{\mathfrak{H}} \stackrel{\Delta}{=} \{|\varphi\rangle \mid \langle\varphi|\varphi\rangle = 1\}$$

$$M_{\mathfrak{H}}(|\varphi\rangle, \rho, |\psi\rangle) \stackrel{\Delta}{\iff} \Pi_{\rho}|\psi\rangle \neq |0\rangle \text{ and } |\varphi\rangle = \frac{\Pi_{\rho}|\psi\rangle}{\|\Pi_{\rho}|\psi\rangle\|}$$

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Verification Relation

$$|\varphi\rangle \blacktriangleright_{\mathfrak{H}} p \iff |\varphi\rangle \in p$$

Definition

Given an *orthomodular lattice* L , we define the model $\mathfrak{L}_L = (A_{\mathfrak{L}}, M_{\mathfrak{L}})$ by

$$A_{\mathfrak{L}} \stackrel{\Delta}{=} L^* \quad \text{where} \quad L^* \stackrel{\Delta}{=} L \setminus \{\perp\}$$

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TWO THEOREMS ON MODELS

Definition

Given a model $\mathfrak{G} = (A, M)$ of \mathcal{T}_L , for all $a \in A$,

$$\llbracket a \rrbracket_{\mathfrak{G}} \triangleq \{p \in L \mid a \blacktriangleright_{\mathfrak{G}} p\}$$

Theorem¹

If $\mathfrak{G} = (A, M)$ is a model of $\mathcal{T}_L(\mathcal{H})$ with $\dim \mathcal{H} \geq 3$, then for all $a \in A$, $\llbracket a \rrbracket_{\mathfrak{G}}$ contains at most one vector ray.

¹arXiv:quant-ph/0610066

Kochen-Specker Theorem

Pick **exactly** one element in each maximal orthogonal family of vectors

Kochen-Specker 117 vectors in dimension 3

Peres 33 vectors in dimension 3

Cabello 17 vectors in dimension 4

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Pick **exactly** one element in each maximal orthogonal family of vectors

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Here

Pick **at most** one element in each maximal orthogonal family of vectors

2 vectors in dimension 3

Theorem

If $3 \leq \dim \mathcal{H} < \infty$ and $\mathfrak{G} = (A, M)$ is a model of $\mathcal{T}_{L(\mathcal{H})}$, then

$$\forall a \in A, \exists e(a) \in L(\mathcal{H}) : \llbracket a \rrbracket_{\mathfrak{G}} = e(a)^{\uparrow}$$

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If $3 \leq \dim \mathcal{H} < \infty$ and $\mathfrak{G} = (A, M)$ is a model of $\mathcal{T}_L(\mathcal{H})$, then

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Consequence

If $3 \leq \dim \mathcal{H} < \infty$, in $\mathcal{T}_L(\mathcal{H})$,

$$s \blacktriangleright p \text{ and } s \blacktriangleright q \implies s \blacktriangleright p \wedge q$$

$$\{[a]_{\mathfrak{G}} \mid a \in A\} \simeq L(\mathcal{H})$$

$$\{[[a]]_{\mathfrak{G}} \mid a \in A\} \simeq L(\mathcal{H})$$

[The Metaphysical Disaster is] the error of **equating properties** of a physical system on the one hand **with experimentally testable propositions** about the system on the other hand.

Unfortunately, this is precisely what is done in conventional Hilbert-space based quantum mechanics where both properties and experimentally testable propositions are represented by projection operators.

(David Foulis, private communication)

SEQUENCES OF OUTCOMES

A **word** on L is a **finite sequence of elements** of L

Definition

Given a model $\mathfrak{G} = (A, M)$ of \mathcal{T}_L , a word $\mathbf{p} = p_1 p_2 \cdots p_n$ is in $\ell(\mathfrak{G})$ iff

$\exists a_0, a_1, \dots, a_n \in A$:

$M(a_0, p_1, a_1)$ and \cdots and $M(a_{n-1}, p_n, a_n)$

$$a_0 \xrightarrow{p_1} a_1 \xrightarrow{p_2} a_2 \xrightarrow{p_3} \cdots \xrightarrow{p_{n-1}} a_{n-1} \xrightarrow{p_n} a_n$$

Theorem

If $3 \leq \dim \mathcal{H} < \infty$,

$$\ell(\mathcal{L}_{\mathcal{L}(\mathcal{H})}) = \ell(\mathcal{S}_{\mathcal{H}})$$

Theorem

If $3 \leq \dim \mathcal{H} < \infty$ and $\mathfrak{G} = (A, M)$ is a model of $\mathcal{T}_L(\mathcal{H})$, then

$$\ell(\mathfrak{G}) = \ell(\mathfrak{L}_L(\mathcal{H})) = \ell(\mathfrak{H}_{\mathcal{H}})$$

Theorem

If \mathcal{H} is such that $3 \leq \dim \mathcal{H} < \infty$, then for all model \mathfrak{G} of $\mathcal{T}_L(\mathcal{H})$,

$$l(\mathfrak{G}) = l(\mathfrak{L}_L(\mathcal{H})) = l(\mathfrak{H}_H)$$

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Theorem

If \mathcal{H} is such that $3 \leq \dim \mathcal{H} < \infty$, then for all model \mathfrak{G} of $\mathcal{T}_{L(\mathcal{H})}$,

$$l(\mathfrak{G}) = l(\mathfrak{L}_{L(\mathcal{H})}) = l(\mathfrak{H}_H) = l(\mathcal{H})$$

Given a word $p_1 p_2 \cdots p_n$ on $L(\mathcal{H})$,

$$\begin{aligned} p_1 p_2 \cdots p_n \in l(\mathcal{H}) &\iff p_1 \& p_2 \& \cdots \& p_n \neq \perp \\ &\iff \prod_{p_n} \prod_{p_{n-1}} \cdots \prod_{p_1} \neq 0 \end{aligned}$$

Theorem

If \mathcal{H} is such that $3 \leq \dim \mathcal{H} < \infty$, then for all model \mathfrak{G} of $\mathcal{T}_L(\mathcal{H})$,

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Can Quantum-Mechanical Description of Physical Reality be Considered Complete?

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Can Quantum-Mechanical Description of Physical Reality be Considered Complete?

Possibilistically yes.