

Categories of Relations as Models of Quantum Theory

Sean Tull

(with Chris Heunen)

University of Oxford

sean.tull@cs.ox.ac.uk

QPL 2015

Categories as Toy Quantum Models

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Examples

- **FHilb** - finite-dimensional Hilbert spaces & linear maps.

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Examples

- **FHilb** - finite-dimensional Hilbert spaces & linear maps.
- **Rel** - sets & *relations*.

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Examples

- **FHilb** - finite-dimensional Hilbert spaces & linear maps.
- **Rel** - sets & *relations*.
- **Spek** - Spekkens' toy model, subcategory of **Rel**.

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Examples

- **FHilb** - finite-dimensional Hilbert spaces & linear maps.
- **Rel** - sets & *relations*.
- **Spek** - Spekkens' toy model, subcategory of **Rel**.

A new class of models

Rel(C) - the category of relations of a regular category \mathbf{C} .

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Examples

- **FHilb** - finite-dimensional Hilbert spaces & linear maps.
- **Rel** - sets & *relations*.
- **Spek** - Spekkens' toy model, subcategory of **Rel**.

A new class of models

Rel(C) - the category of relations of a regular category \mathbf{C} .

- Surprising connections: mixing \rightsquigarrow groupoids & categorification!

Categories as Toy Quantum Models

Categorical Quantum Mechanics lets us see any dagger compact \mathbf{C} as a toy model of quantum theory.

Examples

- **FHilb** - finite-dimensional Hilbert spaces & linear maps.
- **Rel** - sets & *relations*.
- **Spek** - Spekkens' toy model, subcategory of **Rel**.

A new class of models

Rel(C) - the category of relations of a regular category \mathbf{C} .

- Surprising connections: mixing \rightsquigarrow groupoids & categorification!
- Quantum-like behaviour without superposition.

Dagger Compact Categories

Dagger Compact Categories

Dagger compact category \mathbf{D} :

Dagger Compact Categories

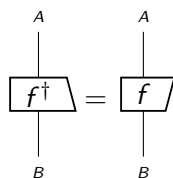
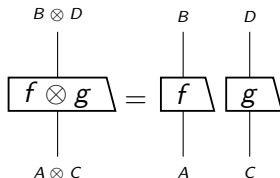
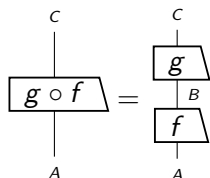
Dagger compact category \mathbf{D} :

- symmetric monoidal $A \otimes B \simeq B \otimes A$
- compact closed $A \dashv A^*$
- $\dagger : \mathbf{D}^{\text{op}} \rightarrow \mathbf{D}$ with $A^\dagger = A$

Dagger Compact Categories

Dagger compact category \mathbf{D} :

- symmetric monoidal $A \otimes B \simeq B \otimes A$
- compact closed $A \dashv A^*$
- $\dagger : \mathbf{D}^{\text{op}} \rightarrow \mathbf{D}$ with $A^\dagger = A$



Categories of Relations

Categories of Relations

A **relation** $R : A \dashv\vdash B$ in \mathbf{C} is a subobject $R \hookrightarrow A \times B$.

Categories of Relations

A **relation** $R : A \dashv\dashv B$ in \mathbf{C} is a subobject $R \hookrightarrow A \times B$.
 \mathbf{C} **regular** $\rightsquigarrow \mathbf{Rel}(\mathbf{C})$ dagger compact.

Categories of Relations

A **relation** $R : A \dashv\vdash B$ in \mathbf{C} is a subobject $R \twoheadrightarrow A \times B$.

\mathbf{C} **regular** $\rightsquigarrow \mathbf{Rel}(\mathbf{C})$ dagger compact.

Internal logic: can pretend we're in $\mathbf{Rel}(\mathbf{Set})$, use \wedge and \exists :

Categories of Relations

A **relation** $R : A \dashrightarrow B$ in \mathbf{C} is a subobject $R \hookrightarrow A \times B$.

\mathbf{C} **regular** $\rightsquigarrow \mathbf{Rel}(\mathbf{C})$ dagger compact.

Internal logic: can pretend we're in $\mathbf{Rel}(\mathbf{Set})$, use \wedge and \exists :

$$S \circ R = \{(a, c) \in A \times C \mid (\exists b \in B) R(a, b) \wedge S(b, c)\} \hookrightarrow A \times C$$

Categories of Relations

A **relation** $R : A \dashrightarrow B$ in \mathbf{C} is a subobject $R \hookrightarrow A \times B$.

\mathbf{C} **regular** $\rightsquigarrow \mathbf{Rel}(\mathbf{C})$ dagger compact.

Internal logic: can pretend we're in $\mathbf{Rel}(\mathbf{Set})$, use \wedge and \exists :

$$S \circ R = \{(a, c) \in A \times C \mid (\exists b \in B) R(a, b) \wedge S(b, c)\} \hookrightarrow A \times C$$

$$R^\dagger = \{(b, a) \in B \times A \mid R(a, b)\}, \quad \otimes \text{ from } \times \text{ in } \mathbf{C}$$

Categories of Relations

A **relation** $R : A \dashrightarrow B$ in \mathbf{C} is a subobject $R \hookrightarrow A \times B$.

\mathbf{C} **regular** $\rightsquigarrow \mathbf{Rel}(\mathbf{C})$ dagger compact.

Internal logic: can pretend we're in $\mathbf{Rel}(\mathbf{Set})$, use \wedge and \exists :

$$S \circ R = \{(a, c) \in A \times C \mid (\exists b \in B) R(a, b) \wedge S(b, c)\} \hookrightarrow A \times C$$

$$R^\dagger = \{(b, a) \in B \times A \mid R(a, b)\}, \quad \otimes \text{ from } \times \text{ in } \mathbf{C}$$

Examples

Regular: any topos, category of algebras, abelian category.

Categories of Relations

A **relation** $R : A \dashrightarrow B$ in \mathbf{C} is a subobject $R \rightarrow A \times B$.

\mathbf{C} **regular** $\rightsquigarrow \mathbf{Rel}(\mathbf{C})$ dagger compact.

Internal logic: can pretend we're in $\mathbf{Rel}(\mathbf{Set})$, use \wedge and \exists :

$$S \circ R = \{(a, c) \in A \times C \mid (\exists b \in B) R(a, b) \wedge S(b, c)\} \rightarrow A \times C$$

$$R^\dagger = \{(b, a) \in B \times A \mid R(a, b)\}, \quad \otimes \text{ from } \times \text{ in } \mathbf{C}$$

Examples

Regular: any topos, category of algebras, abelian category.

$\mathbf{Rel}(\mathbf{Set}) = \mathbf{Rel}$.

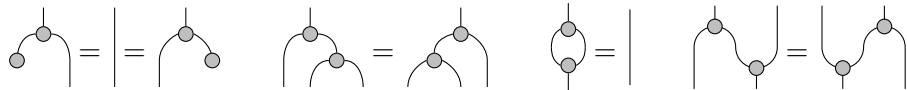
$\mathbf{Rel}(\mathbf{Grp})$: subgroups $R \leq G \times H$.

$\mathbf{Rel}(\mathbf{Vect}_k)$: subspaces $R \leq V \oplus W$ - see 'Categories in Control'.

C^* -Algebras become Groupoids

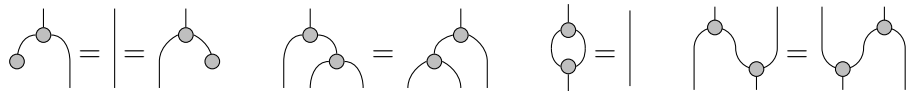
C*-Algebras become Groupoids

A special dagger Frobenius structure (A, μ, ν) satisfies:



C*-Algebras become Groupoids

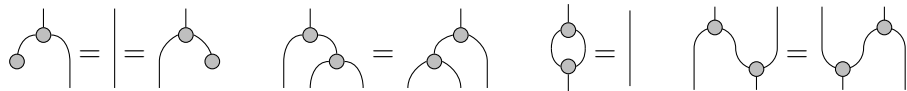
A special dagger Frobenius structure $(A, \circlearrowleft, \circlearrowright)$ satisfies:



In **FHilb**: finite-dimensional C*-algebras. In **Rel**: small groupoids!

C*-Algebras become Groupoids

A special dagger Frobenius structure (A, μ, ν) satisfies:



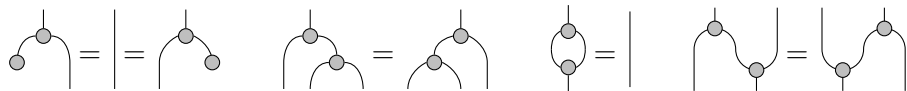
In **FHilb**: finite-dimensional C*-algebras. In **Rel**: small groupoids!

Theorem

Special dagger Frobenius structures μ, ν in $\mathbf{Rel}(\mathbf{C})$ are the same as internal groupoids in \mathbf{C} .

C*-Algebras become Groupoids

A special dagger Frobenius structure $(A, \circlearrowleft, \circlearrowright)$ satisfies:



In **FHilb**: finite-dimensional C*-algebras. In **Rel**: small groupoids!

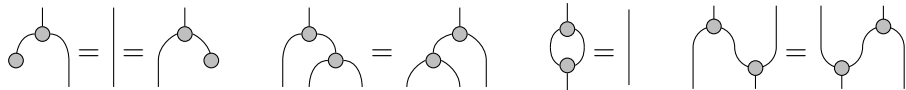
Theorem

Special dagger Frobenius structures \circlearrowleft in $\mathbf{Rel}(\mathbf{C})$ are the same as internal groupoids in \mathbf{C} .

$$\begin{array}{c}
 \longleftarrow t \quad \circlearrowleft i \\
 O \xrightarrow{u} A \xleftarrow{m} A \times_O A \\
 \longleftarrow s
 \end{array}$$

C*-Algebras become Groupoids

A special dagger Frobenius structure $(A, \circlearrowleft, \circlearrowright)$ satisfies:



In **FHilb**: finite-dimensional C*-algebras. In **Rel**: small groupoids!

Theorem

Special dagger Frobenius structures \circlearrowleft in **Rel(C)** are the same as internal groupoids in **C**.

Examples

Set: small groupoids.

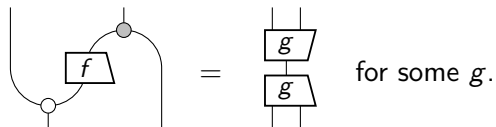
Grp: strict 2-groups (Baez-Lauda) \iff crossed modules.

Vect_k: 2-vector spaces (Baez-Crans).

Completely Positive Relations

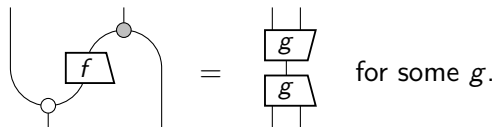
Completely Positive Relations

CP(D): Frobenius \circlearrowleft in **D** & **completely positive** $(A, \circlearrowleft) \xrightarrow{f} (B, \circlearrowleft)$



Completely Positive Relations

CP(D): Frobenius \circlearrowleft in **D** & **completely positive** $(A, \circlearrowleft) \xrightarrow{f} (B, \circlearrowleft)$

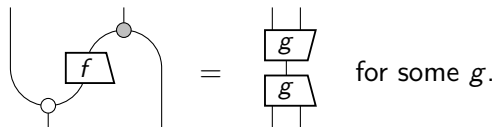


Examples

CP(FHilb): finite-dimensional C^* -algebras & completely positive maps.

Completely Positive Relations

CP(D): Frobenius \circlearrowleft in **D** & **completely positive** $(A, \circlearrowleft) \xrightarrow{f} (B, \circlearrowleft)$



Examples

CP(FHilb): finite-dimensional C^* -algebras & completely positive maps.

CP(Rel): groupoids & relations such that

$$R(a, b) \Rightarrow R(a^{-1}, b^{-1}) \wedge R(\text{id}_{\text{dom}(a)}, \text{id}_{\text{dom}(b)})$$

Mal'cev Categories

Mal'cev Categories

In **Grp** or **Vect**_k:

In **Grp** or **Vect**_k:

$$R(a, c) \wedge R(b, c) \wedge R(b, d) \implies R(a, d)$$

In **Grp** or **Vect**_k:

$$R(a, c) \wedge R(b, c) \wedge R(b, d) \implies R(a, d)$$

$$(a, c) - (b, c) + (b, d) = (a, d)$$

Mal'cev Categories

In **Grp** or **Vect_k**:

$$R(a, c) \wedge R(b, c) \wedge R(b, d) \implies R(a, d)$$

C is **Mal'cev** when holds $\forall R$.

Mal'cev Categories

In **Grp** or **Vect_k**:

$$R(a, c) \wedge R(b, c) \wedge R(b, d) \implies R(a, d)$$

C is **Mal'cev** when holds $\forall R$.

Theorem

*When **C** is Mal'cev regular we get an equivalence of categories*

$$\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Gpd}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$$

Mal'cev Categories

In **Grp** or **Vect_k**:

$$R(a, c) \wedge R(b, c) \wedge R(b, d) \implies R(a, d)$$

C is **Mal'cev** when holds $\forall R$.

Theorem

When **C** is Mal'cev regular we get an equivalence of categories

$$\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Gpd}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$$

Examples

Grp, **Vect_k**, Rings, Lie algebras. Any abelian category.

Mal'cev Categories

In **Grp** or **Vect_k**:

$$R(a, c) \wedge R(b, c) \wedge R(b, d) \implies R(a, d)$$

C is **Mal'cev** when holds $\forall R$.

Theorem

When **C** is Mal'cev regular we get an equivalence of categories

$$\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Gpd}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$$

Examples

Grp, **Vect_k**, Rings, Lie algebras. Any abelian category.

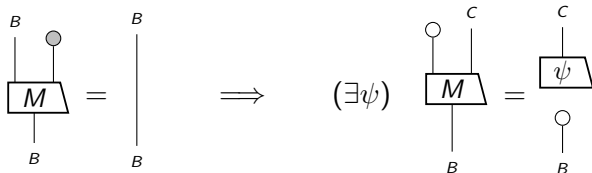
$$\mathbf{CP}(\mathbf{Rel}(\mathbf{Grp})) \simeq \mathbf{Rel}(\mathbf{CrossedModules})$$

$$\mathbf{CP}(\mathbf{Rel}(\mathbf{Vect}_k)) \simeq \mathbf{Rel}(\mathbf{2Vect}_k)$$

Quantum Properties of $\mathbf{Rel}(\mathbf{C})$

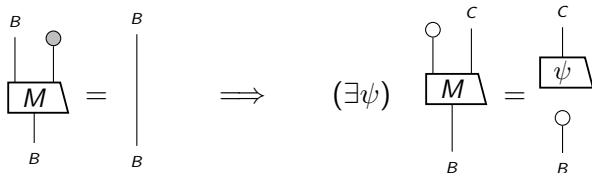
Quantum Properties of $\mathbf{Rel}(\mathbf{C})$

Heisenberg Uncertainty Principle:



Quantum Properties of $\mathbf{Rel}(\mathbf{C})$

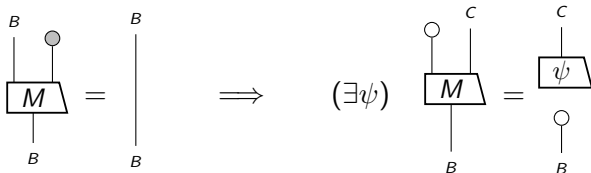
Heisenberg Uncertainty Principle:



Rel	Rel(C) Mal'cev	FHilb
		✓

Quantum Properties of $\mathbf{Rel}(\mathbf{C})$

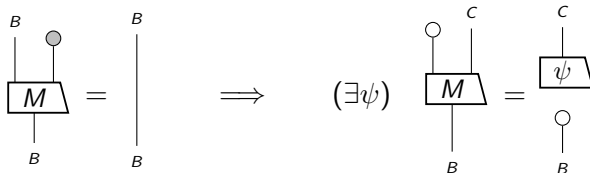
Heisenberg Uncertainty Principle:



Rel	Rel(C) Mal'cev	FHilb
\times		\checkmark

Quantum Properties of $\mathbf{Rel}(\mathbf{C})$

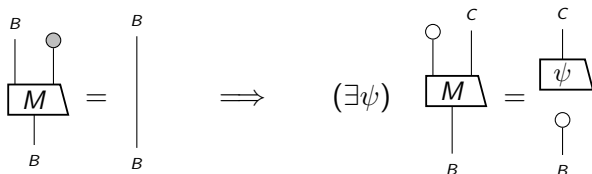
Heisenberg Uncertainty Principle:



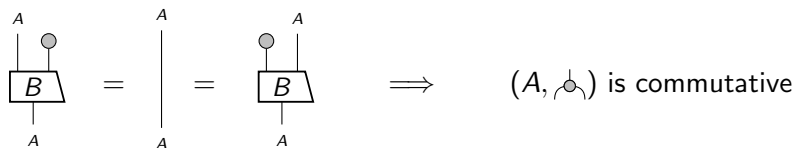
Rel	Rel(C) Mal'cev	FHilb
X	✓	✓

Quantum Properties of $\mathbf{Rel}(\mathbf{C})$

Heisenberg Uncertainty Principle:



No-Broadcasting:

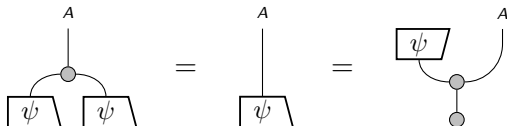


\mathbf{Rel}	$\mathbf{Rel}(\mathbf{C})$ Mal'cev	\mathbf{FHilb}
\times	\checkmark	\checkmark

Non-Quantum Features of $\mathbf{Rel}(\mathbf{C})$

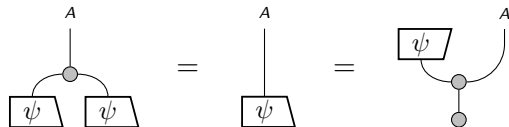
Non-Quantum Features of $\mathbf{Rel}(\mathbf{C})$

\mathbf{C} Mal'cev \rightsquigarrow any state is a **projection**:

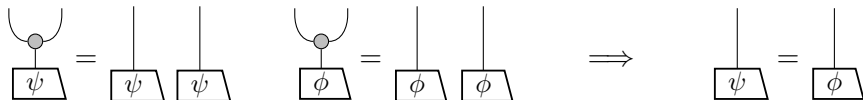


Non-Quantum Features of $\mathbf{Rel}(\mathbf{C})$

\mathbf{C} Mal'cev \rightsquigarrow any state is a **projection**:



\mathbf{C} has zero object (e.g. \mathbf{Grp} , \mathbf{Vect}_k) \rightsquigarrow no distinct classical data:



Summary

- **Rel(C)** gives us many new toy quantum models.

Summary

- **Rel(C)** gives us many new toy quantum models.
- Internal logic makes life easy: can pretend we're in **Rel**.

Summary

- $\mathbf{Rel}(\mathbf{C})$ gives us many new toy quantum models.
- Internal logic makes life easy: can pretend we're in \mathbf{Rel} .

When \mathbf{C} is Mal'cev:

- $\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$, mixing \rightsquigarrow categorification.

Summary

- $\mathbf{Rel}(\mathbf{C})$ gives us many new toy quantum models.
- Internal logic makes life easy: can pretend we're in \mathbf{Rel} .

When \mathbf{C} is Mal'cev:

- $\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$, mixing \rightsquigarrow categorification.
- $\mathbf{Rel}(\mathbf{C})$ lacks **superposition** + while retaining quantum-like behaviour.

Summary

- $\mathbf{Rel}(\mathbf{C})$ gives us many new toy quantum models.
- Internal logic makes life easy: can pretend we're in \mathbf{Rel} .

When \mathbf{C} is Mal'cev:

- $\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$, mixing \rightsquigarrow categorification.
- $\mathbf{Rel}(\mathbf{C})$ lacks **superposition** + while retaining quantum-like behaviour.

“least quantum”
 \mathbf{Rel}

\leftrightarrow
 $\mathbf{Rel}(\mathbf{C})$ Mal'cev

“most quantum”
 \mathbf{FHilb}

Summary

- $\mathbf{Rel}(\mathbf{C})$ gives us many new toy quantum models.
- Internal logic makes life easy: can pretend we're in \mathbf{Rel} .

When \mathbf{C} is Mal'cev:

- $\mathbf{CP}(\mathbf{Rel}(\mathbf{C})) \simeq \mathbf{Rel}(\mathbf{Cat}(\mathbf{C}))$, mixing \rightsquigarrow categorification.
- $\mathbf{Rel}(\mathbf{C})$ lacks [superposition](#) + while retaining quantum-like behaviour.

“least quantum”
 \mathbf{Rel}

\leftrightarrow
 $\mathbf{Rel}(\mathbf{C})$ Mal'cev

“most quantum”
 \mathbf{FHilb}

Thanks for listening!