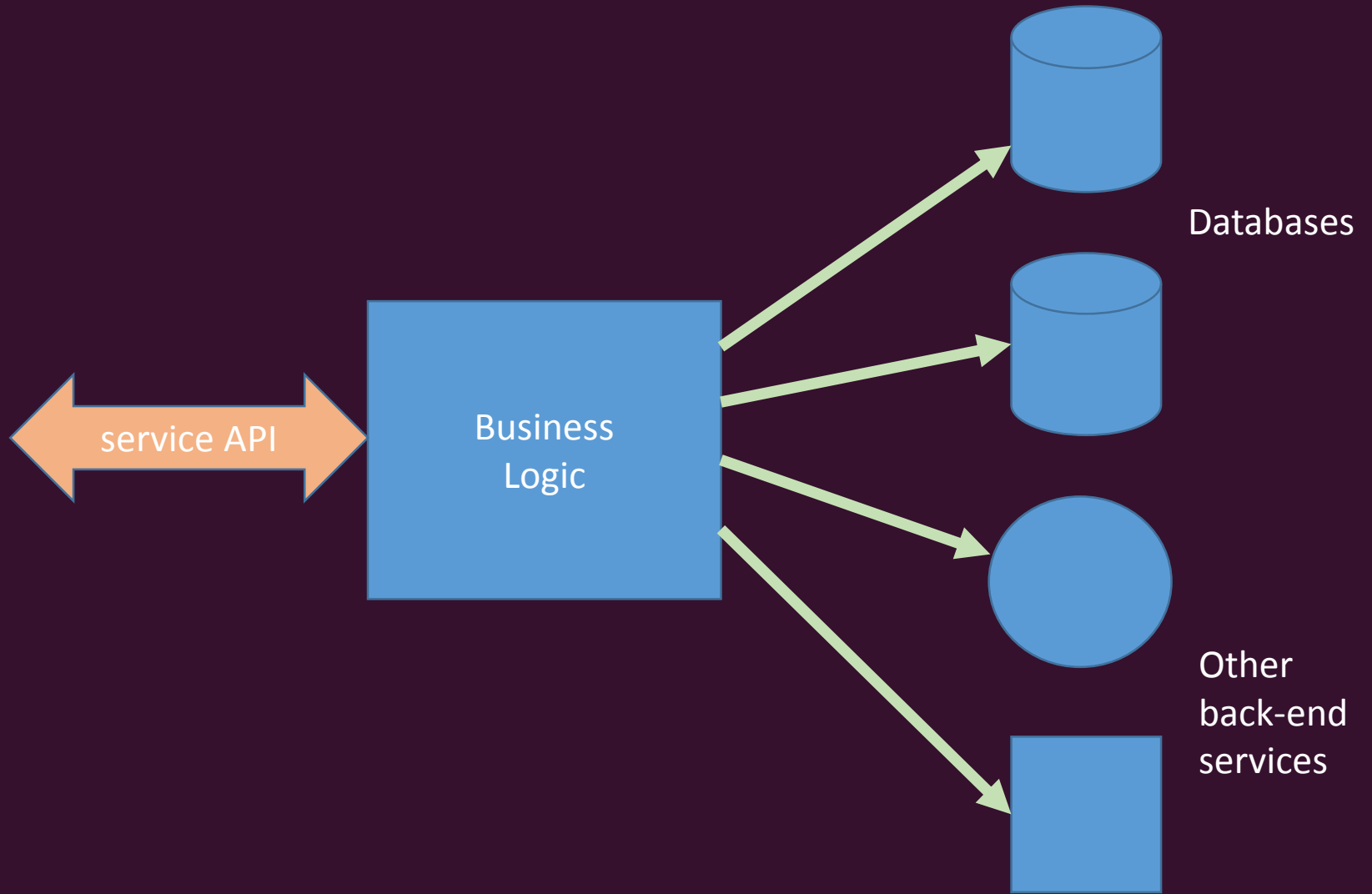


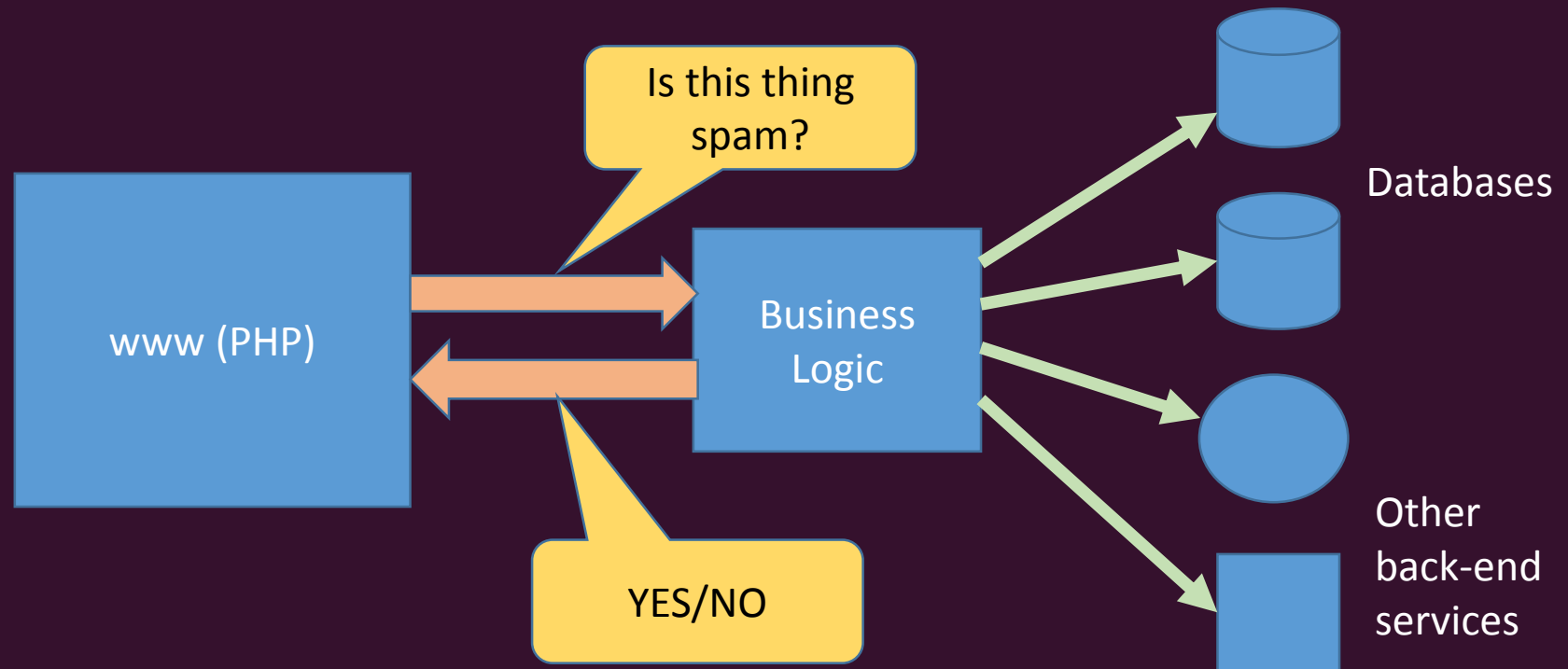
The Haxl Project at Facebook

Simon Marlow
Jon Coens
Louis Brandy
Jon Purdy
& others

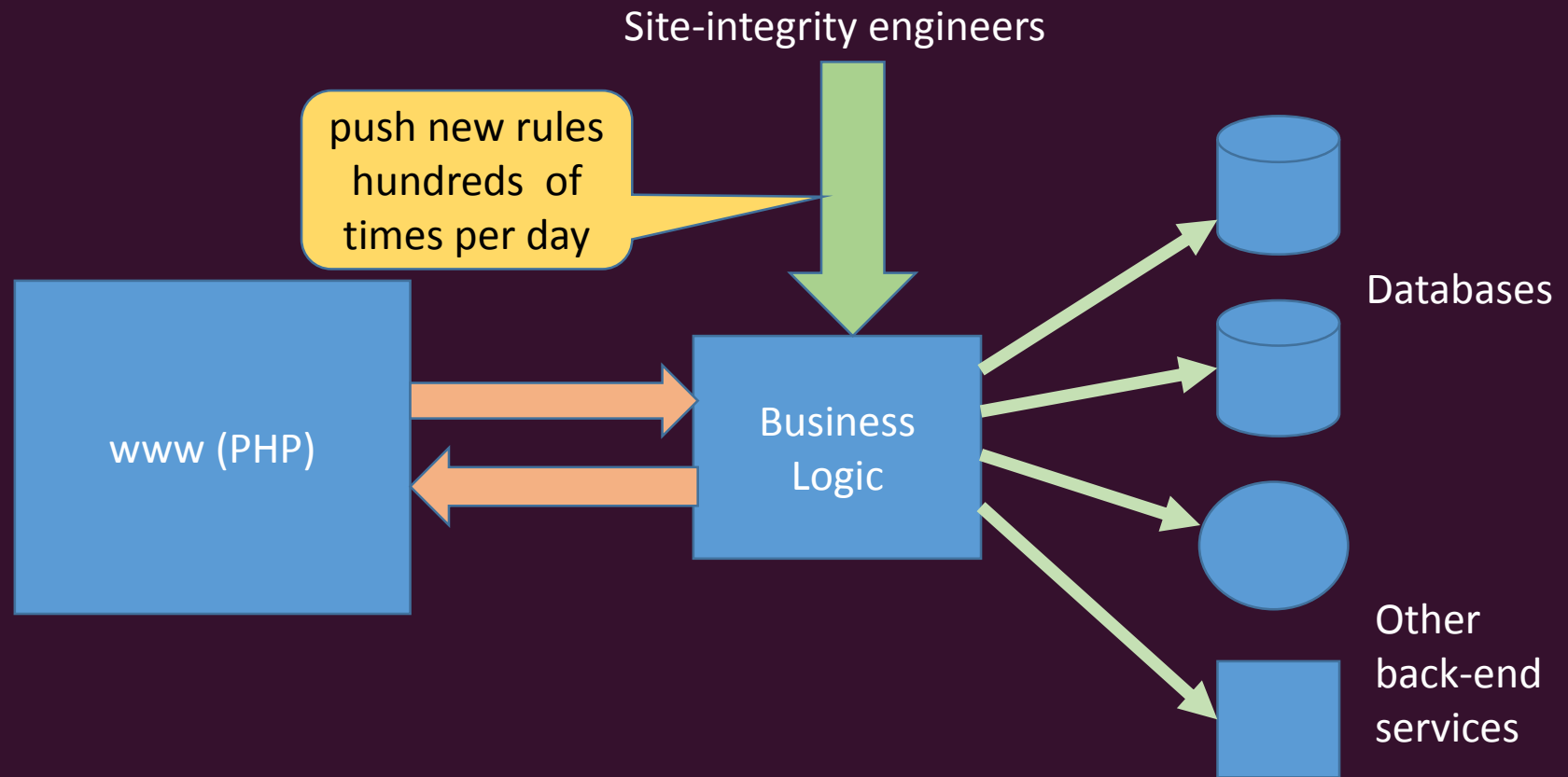




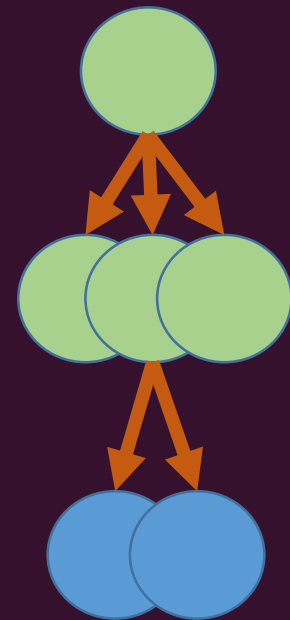
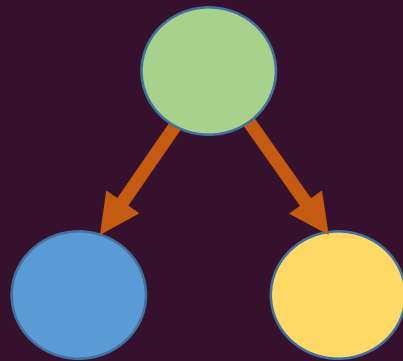
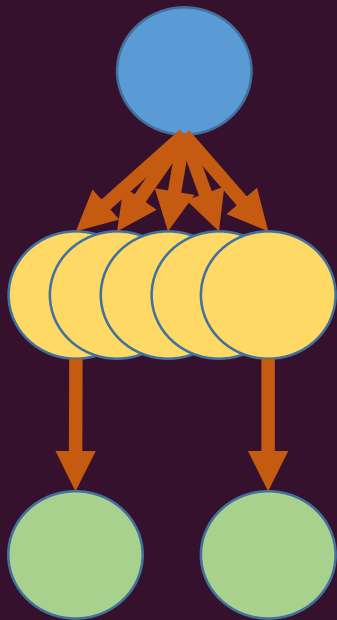
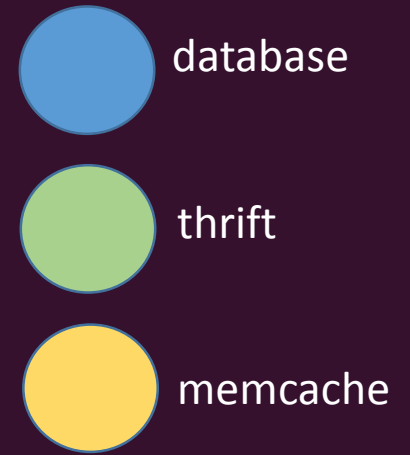
Use case: fighting spam



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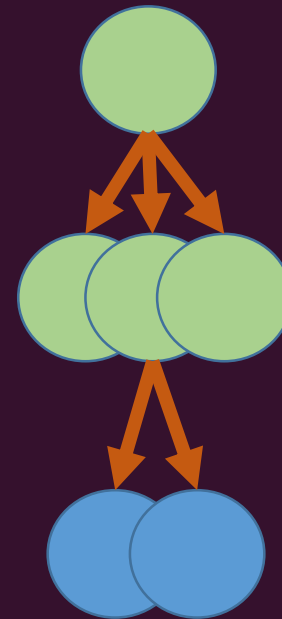
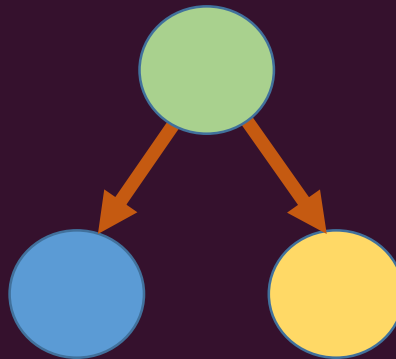
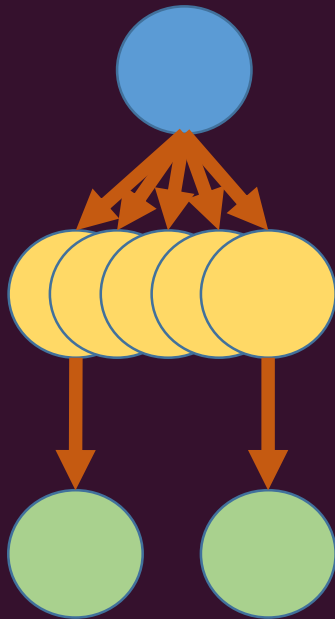
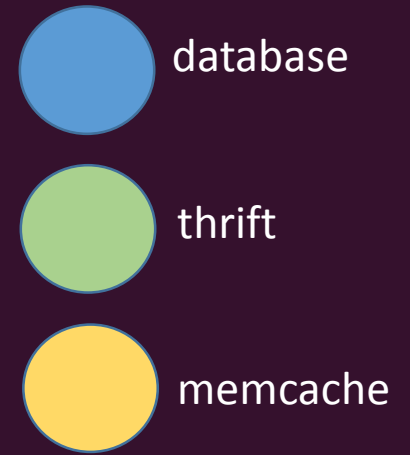


Data dependencies in a computation



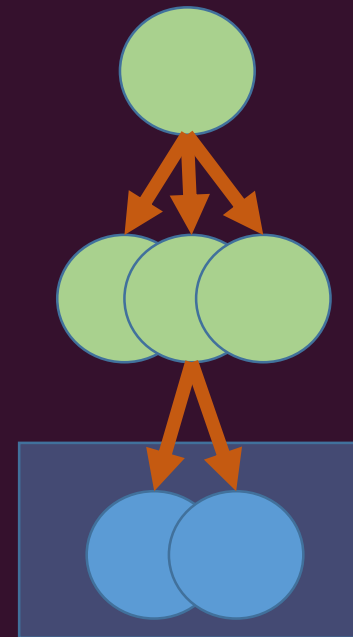
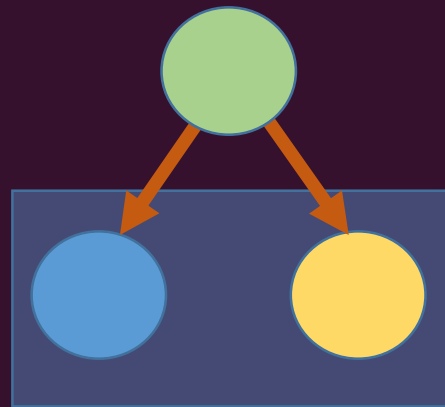
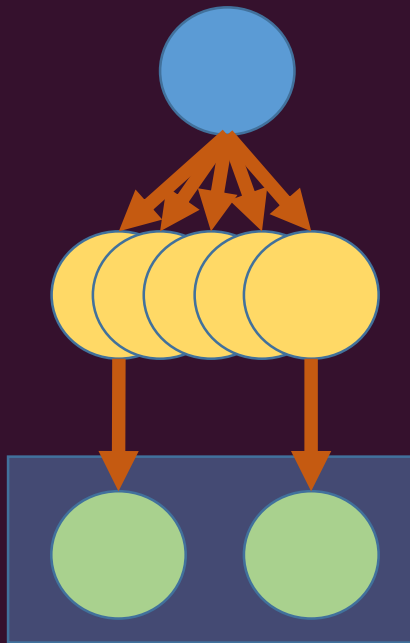
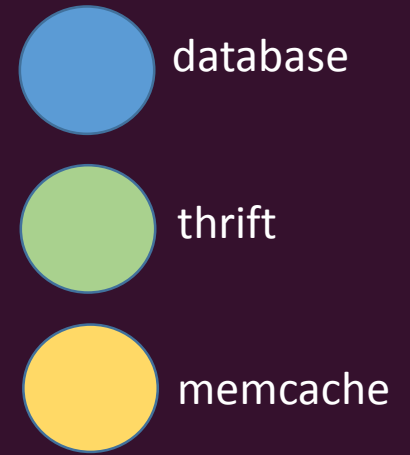
Code wants to be structured hierarchically

- abstraction
- modularity



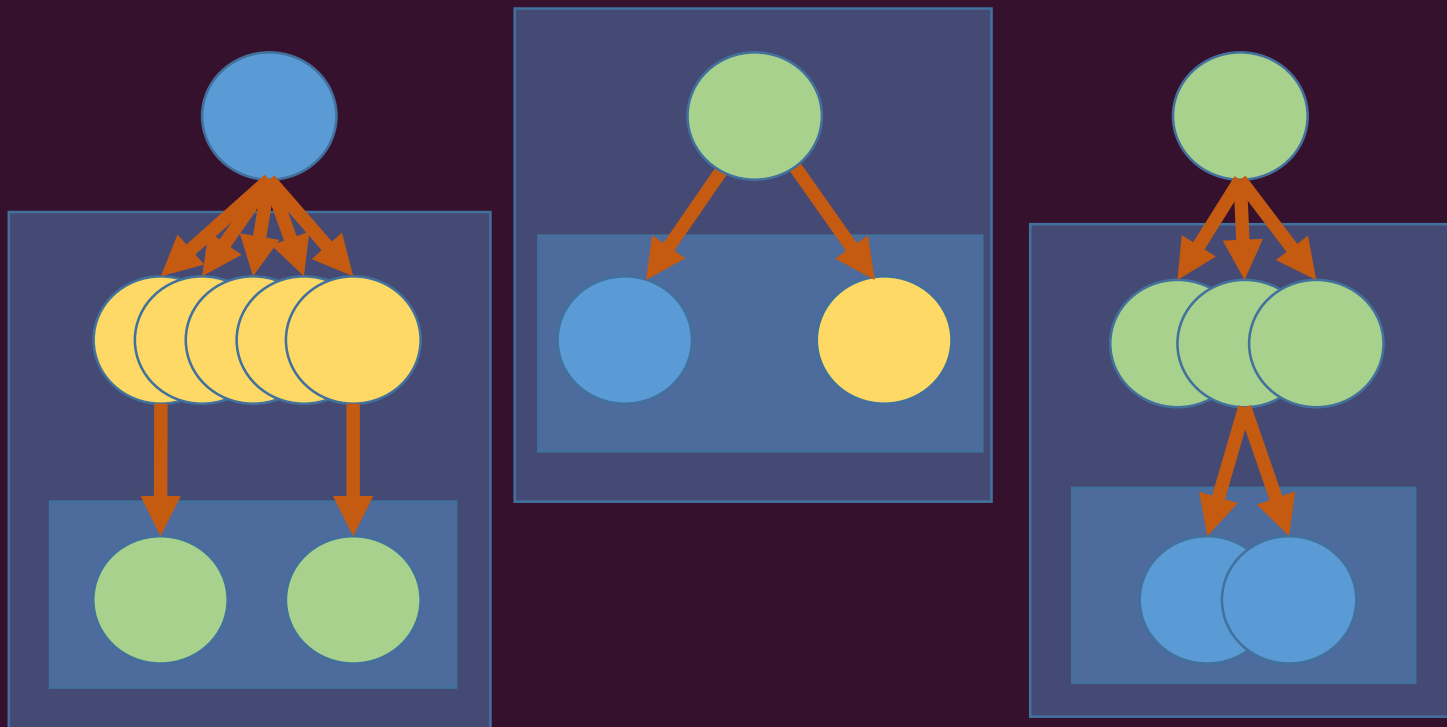
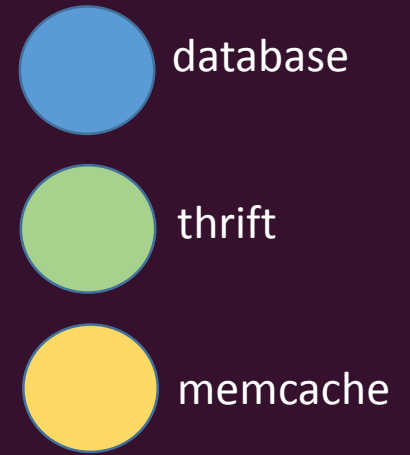
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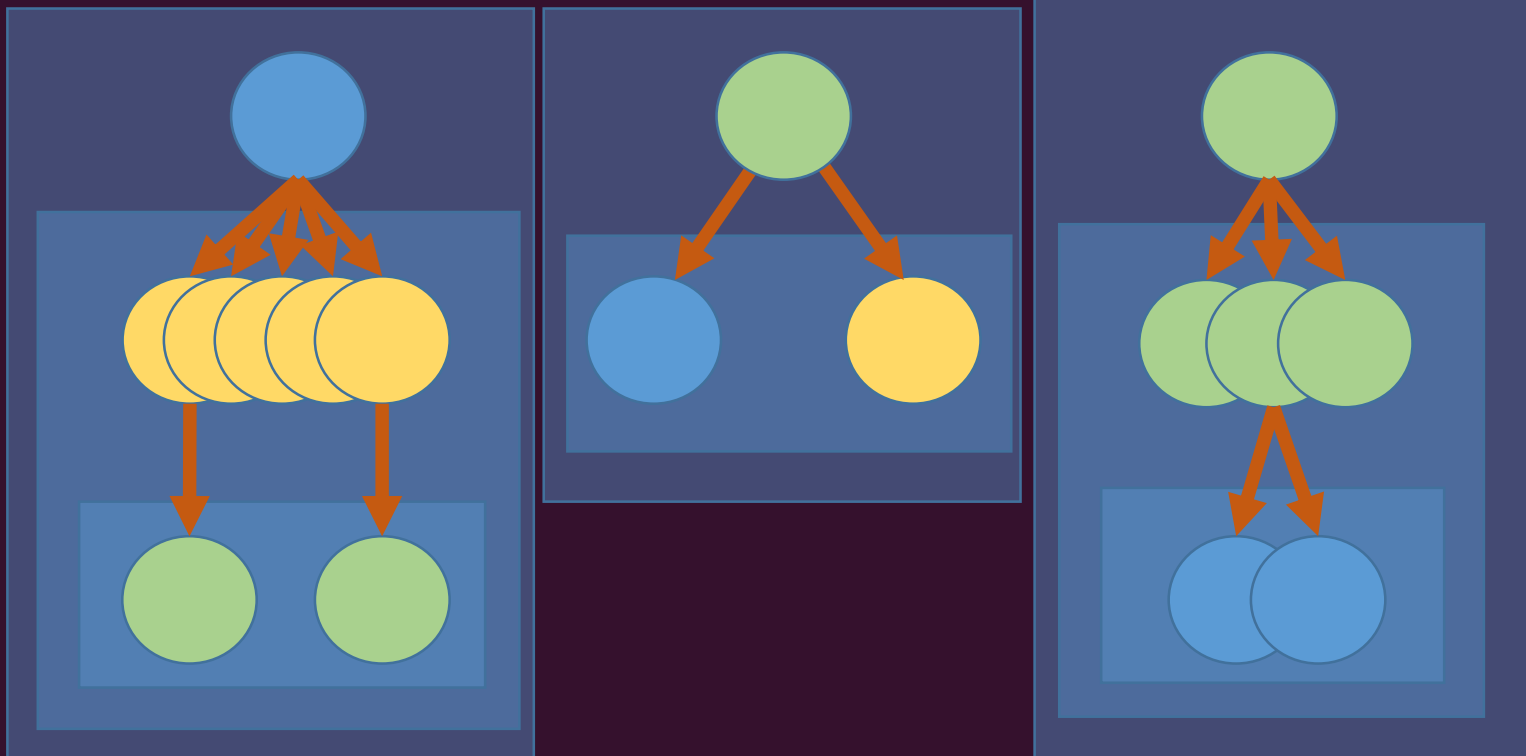
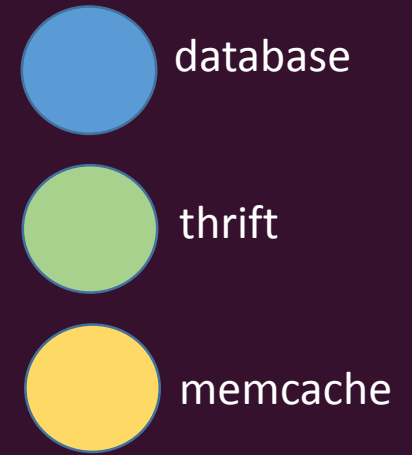
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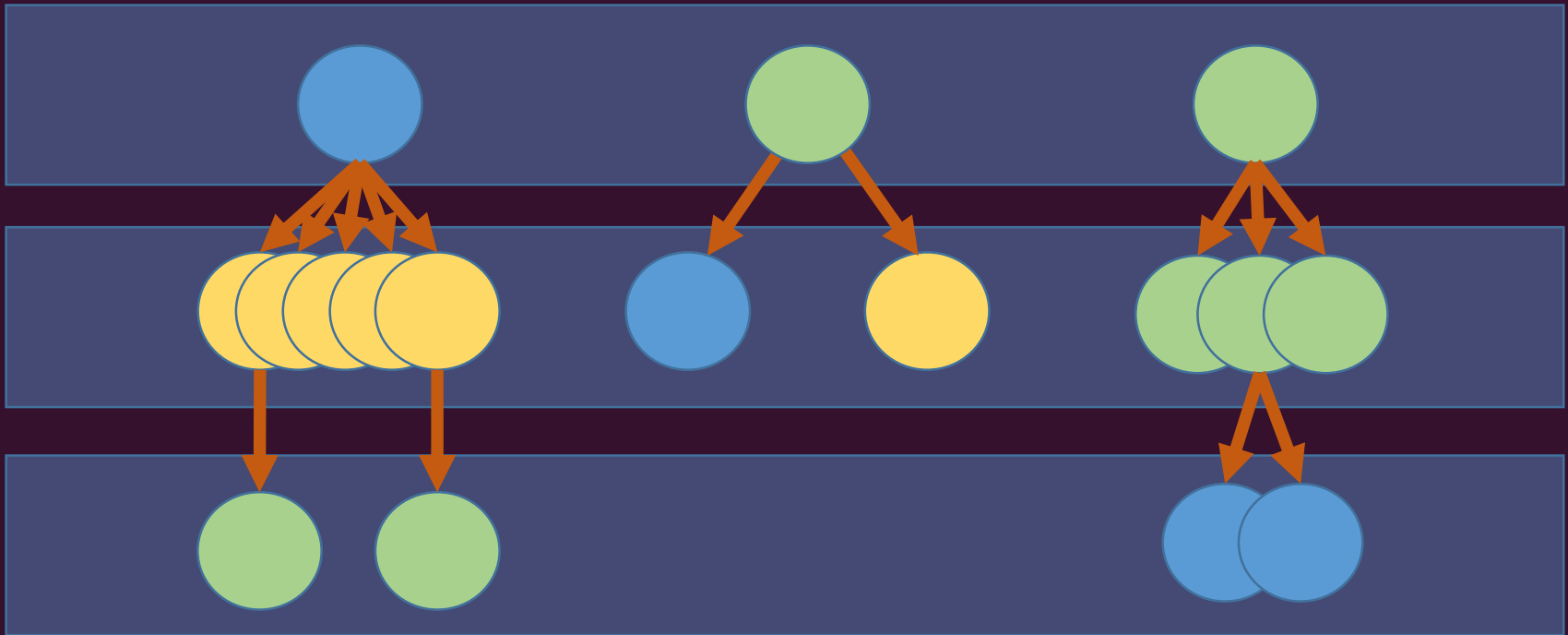
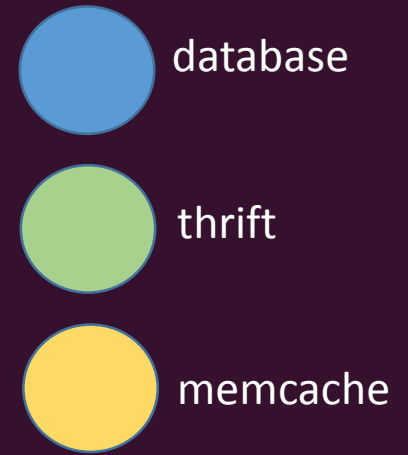
Code wants to be structured hierarchically

- abstraction
- modularity



Execution wants to be structured horizontally

- Overlap multiple requests
- Batch requests to the same data source
- Cache multiple requests for the same data



- Furthermore, each data source has different characteristics
 - Batch request API?
 - Sync or async API?
 - Set up a new connection for each request, or keep a pool of connections around?
- Want to abstract away from *all* of this in the business logic layer

But we know how to do this!

But we know how to do this!

- Concurrency.

Threads let us keep our abstractions & modularity while executing things at the same time.

- Caching/batching can be implemented as a service in the process
 - as we do with the IO manager in GHC

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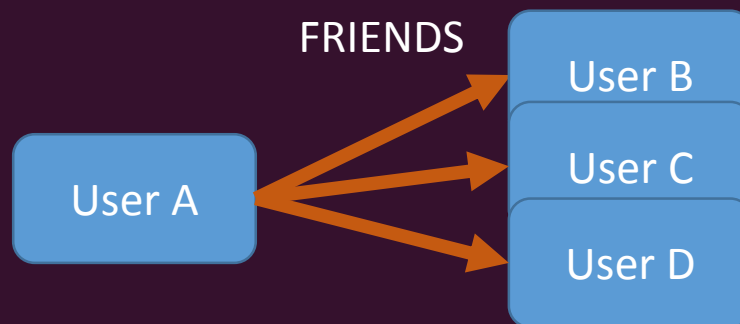
- Example...

- x and y are Facebook users
- suppose we want to compute the number of friends that x and y have in common
- simplest way to write this:

```
length (intersect (friendsOf x) (friendsOf y))
```


Brief detour: TAO

- TAO implements Facebook's data model
 - most important data source we need to deal with
- Data is a graph
 - Nodes are "objects", identified by 64-bit ID
 - Edges are "assocs" (directed; a pair of 64-bit IDs)
- Objects and assocs have a type
 - object fields determined by the type
- Basic operations:
 - Get the object with a given ID
 - Get the assocs of a given type from a given ID



- Back to our example

```
length (intersect (friendsOf x) (friendsOf y))
```

- **(friendsOf x)** makes a request to TAO to get all the IDs for which there is an assoc of type FRIEND (x,_).
- TAO has a multi-get API; very important that we submit **(friendsOf x)** and **(friendsOf y)** as a single operation.

Using concurrency

- This:

```
length (intersect (friendsOf x) (friendsOf y))
```

Using concurrency

- This:

```
length (intersect (friendsOf x) (friendsOf y))
```

- Becomes this:

```
do
  m1 <- newEmptyMVar
  m2 <- newEmptyMVar
  forkIO (friendsOf x >>= putMVar m1)
  forkIO (friendsOf y >>= putMVar m2)
  fx <- takeMVar m1
  fy <- takeMVar m2
  return (length (intersect fx fy))
```



OH GOD

MY EYES

- Using the async package:

```
do
```

```
  ax <- async (friendsOf x)  
  ay <- async (friendsOf y)  
  fx <- wait ax  
  fy <- wait ay  
  return (length (intersect fx fy))
```



- Using `Control.Concurrent.Async.concurrently`:

```
do
```

```
  (fx,fy) <- concurrently (friendsOf x) (friendsOf y)  
  return (length (intersect fx fy))
```




Why not concurrency?

- **friendsOf x** and **friendsOf y** are
 - obviously independent
 - obviously both needed
 - “pure”

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 - if **friendsOf x** is requested twice, we *must* get the same answer both times
 - *caching is a requirement*

Why not concurrency?

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 - obviously independent
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 - “pure”
- Caching is not just an optimisation:
 - if **friendsOf x** is requested twice, we *must* get the same answer both times
 - *caching is a requirement*
- we don't want the programmer to have to ask for concurrency here

- Could we use `unsafePerformIO`?

```
length (intersect (friendsOf x) (friendsOf y))  
friendsOf = unsafePerformIO ( .. )
```

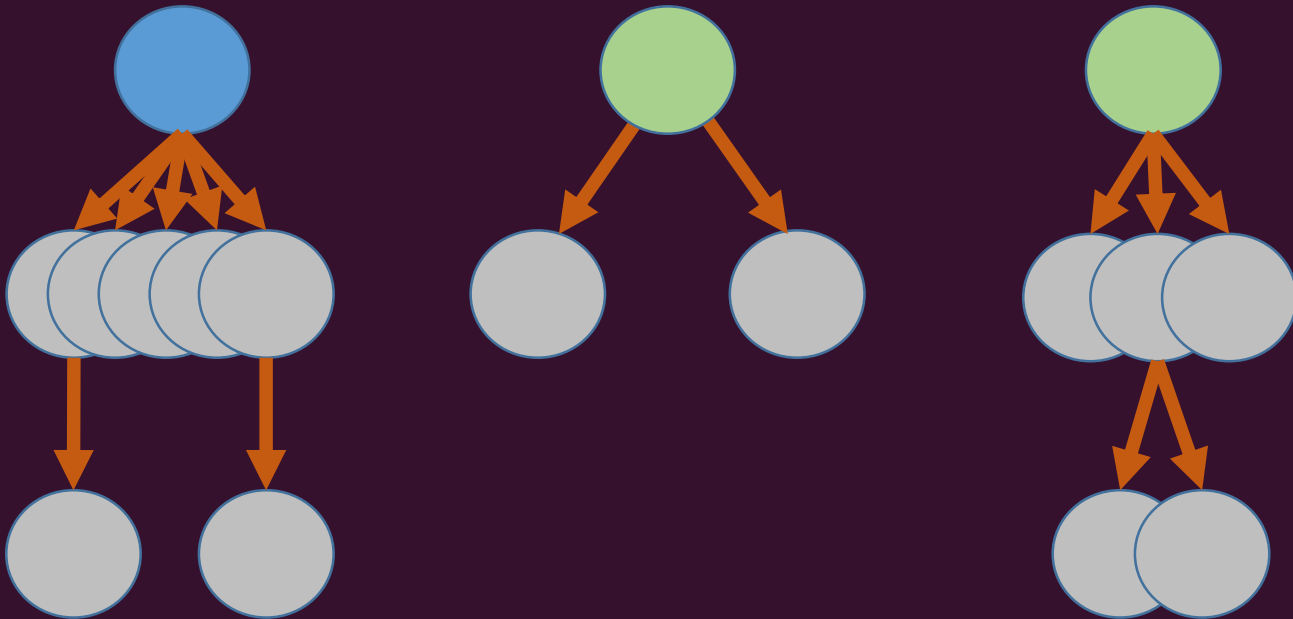
- we could do caching this way, but not concurrency. Execution will stop at the first data fetch.

Central problem

- Reorder execution of an expression to perform data fetching optimally.
- The programming model has no side effects (other than reading)

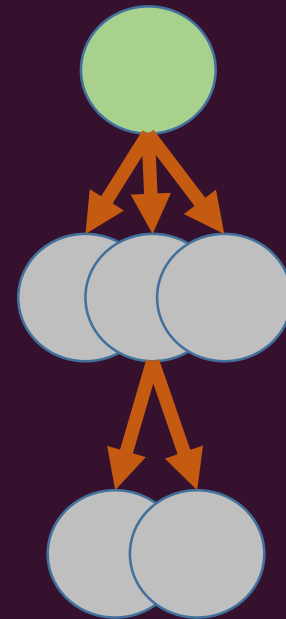
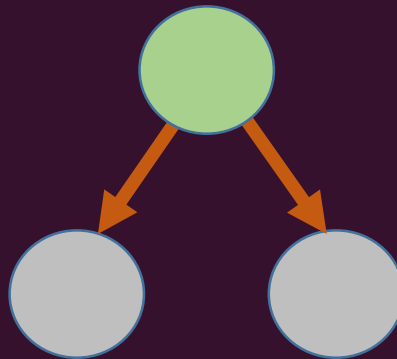
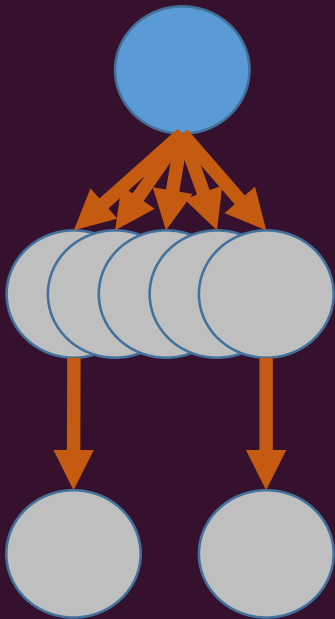
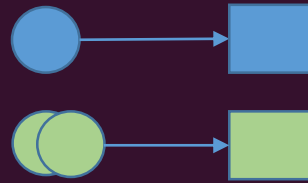
What we would like to do:

- explore the expression along all branches to get a set of data fetches



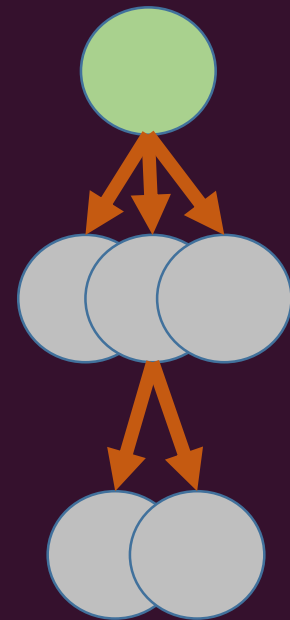
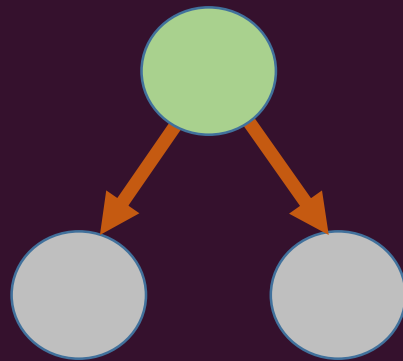
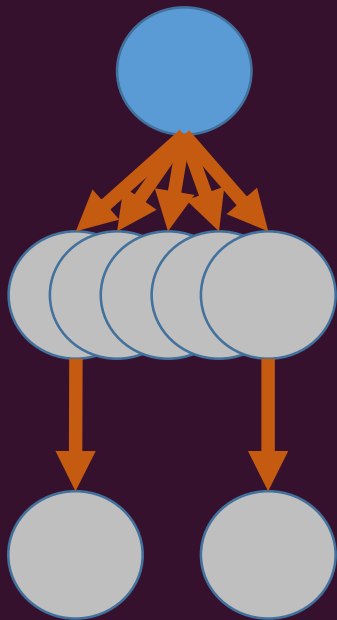
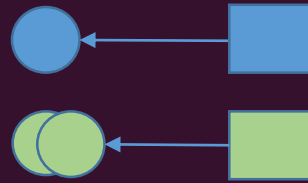
What we would like to do:

- submit the data fetches



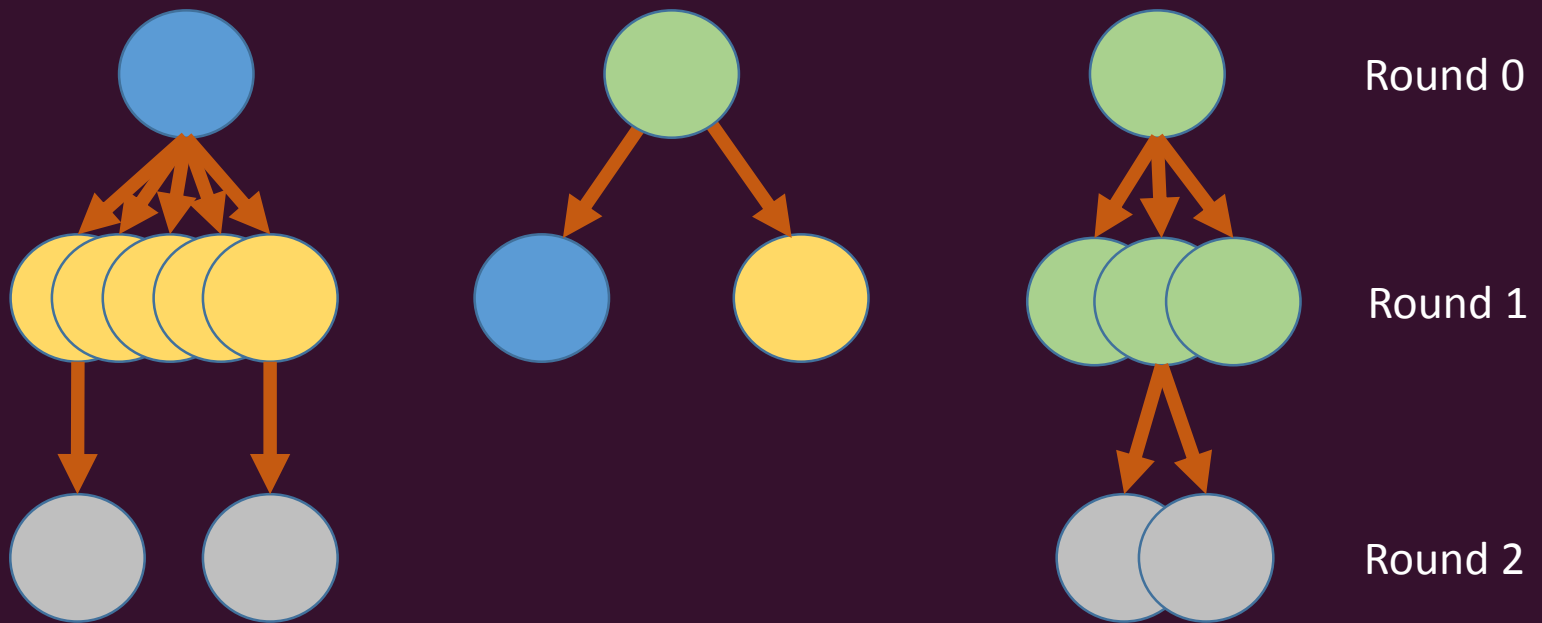
What we would like to do:

- wait for the responses



What we would like to do:

- now the computation is unblocked along multiple paths
- ... explore again
- collect the next batch of data fetches
- and so on



Fighting spam with pure functions

By Louis Brandy on Thursday, 24 January 2013 at 17:31

Like any popular Internet site, Facebook is a target for abuse. Our Site Integrity engineers rely on FXL, a domain-specific language forged in the fires of spam fighting at Facebook, to



- Facebook's existing solution to this problem: FXL
- Lets you write

```
Length(Intersect(FriendsOf(X),FriendsOf(Y)))
```

- And optimises the data fetching correctly.
- But it's an interpreter, and works with an explicit representation of the computation graph.

- We want to run compiled code for efficiency
- And take advantage of Haskell
 - high quality implementation
 - great libraries for writing business logic etc.
- So, how can we implement the right data fetching behaviour in a Haskell DSL?

Start with a concurrency monad

```
newtype Haxl a = Haxl { unHaxl :: Result a }

data Result a = Done a
              | Blocked (Haxl a)

instance Monad Haxl where
  return a = Haxl (Done a)
  m >>= k = Haxl $
    case unHaxl m of
      Done a      -> unHaxl (k a)
      Blocked r   -> Blocked (r >>= k)
```

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```



It's a
Free
Monad

- The concurrency monad lets us run a computation until it blocks, do something, then resume it
- But we need to know what it blocked on...
- Could add some info to the **Blocked** constructor

```
newtype Haxl a = Haxl { unHaxl :: Responses -> Result a }
```

```
data Result a = Done a  
              | Blocked Requests (Haxl a)
```

```
instance Monad Haxl where
```

```
  return a = Haxl $ \_ -> Done a
```

```
  Haxl m >>= k = Haxl $ \resps ->
```

```
    case m resps of
```

```
      Done a          -> unHaxl (k a) resps
```

```
      Blocked reqs r -> Blocked reqs (r >>= k)
```

```
addRequest      :: Request a -> Requests -> Requests
```

```
emptyRequests  :: Requests
```

```
fetchResponse  :: Request a -> Responses -> a
```

```
dataFetch      :: Request a -> Haxl a
```

```
dataFetch req = Haxl $ \_ ->
```

```
  Blocked (addRequest req emptyRequests) $ Haxl $ \resps ->
```

```
    Done (fetchResponse req resps)
```


- Ok so far, but we still get blocked at the first data fetch.

```
numCommonFriends x y = do  
  fx <- friendsOf x  
  fy <- friendsOf y  
  return (length (intersect fx fy))
```

Blocked here

- To explore multiple branches, we need to use **Applicative**

```
instance Applicative Haxl where
```

```
  pure = return
```

```
  Haxl f <*> Haxl a = Haxl $ \resps ->
```

```
    case f resps of
```

```
      Done f' ->
```

```
        case a resps of
```

```
          Done a'          -> Done (f' a')
```

```
          Blocked reqs a' -> Blocked reqs (f' <$> a')
```

```
      Blocked reqs f' ->
```

```
        case a resps of
```

```
          Done a'          -> Blocked reqs (f' <*> return a')
```

```
          Blocked reqs' a' -> Blocked (reqs <> reqs') (f' <*> a')
```

```
<*> :: Applicative f => f (a -> b) -> f a -> f b
```

- This is precisely the advantage of **Applicative** over **Monad**:
 - **Applicative** allows exploration of the structure of the computation
- Our example is now written:

```
numCommonFriends x y =  
  length <$> (intersect <$> friendsOf x <*> friendsOf y)
```

- Or:

```
numCommonFriends x y =  
  length <$> common (friendsOf x) (friendsOf y)  
  where common = liftA2 intersect
```

- Note that we still have the Monad!
- The Monad allows us to make decisions based on values when we need to.

```
do
  fs <- friendsOf x
  if simon `elem` fs
  then ...
  else ...
```



Blocked here

- Batching will not explore the then/else branches
 - exactly what we want.

- But it does mean the programmer should use Applicative composition to get batching.
- This is suboptimal:

```
do
  fx <- friendsOf x
  fy <- friendsOf y
  return (length (intersect fx fy))
```

- So our plan is to
 - provide APIs that batch correctly
 - translate do-notation into Applicative where possible
 - (forthcoming GHC extension)

- We really want bulk operations to benefit from batching.

```
friendsOfFriends id =  
  concat <$> (mapM friendsOf =<< friendsOf id)
```

- But this doesn't work: **mapM** uses **Monad** rather than **Applicative** composition.
- This is why **traverse** exists:

```
traverse :: (Traversable t, Applicative f)  
  => (a -> f b) -> t a -> f (t b)
```

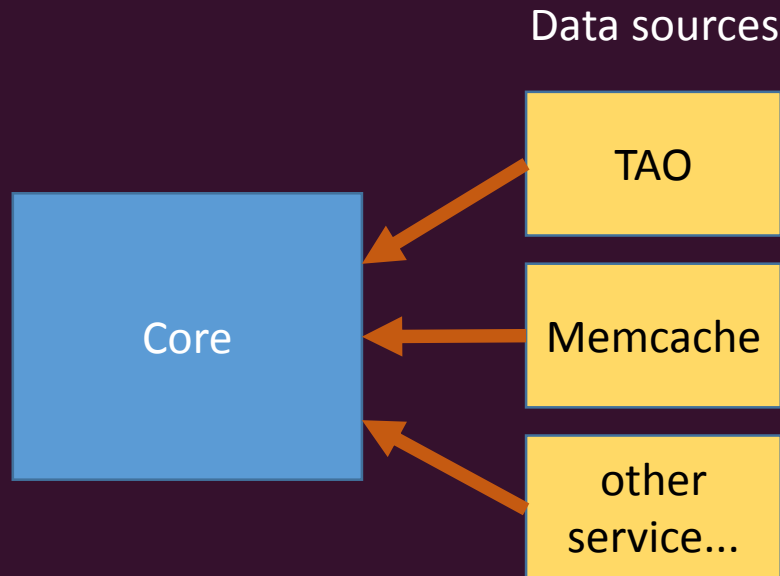
- So in our library, we make **mapM = traverse**
- Also: **sequence = sequenceA**
- Will be fixed once **Applicative** is a superclass of **Monad**

Implementation

- DataSource abstraction
 - Replaying requests
 - Scaling
 - Hot-code swapping
-
- Experience
 - Status etc.

Data Source Abstraction

- We want to structure the system like this:



- Core code includes the monad, caching support etc.
- Core is *generic*: no data sources built-in

How do we arrange this?

- Three ways that a data source interacts with core:
 - issuing a data fetch request
 - persistent state
 - fetching the data
- Package this up in a type class

```
class DataSource req where  
  ...
```

- Let's look at requests first...

parameterised
over the type of
requests

Example Request type

```
data ExampleReq a where
  CountAardvarks :: String -> ExampleReq Int
  ListWombats    :: Id      -> ExampleReq [Id]
  deriving Typeable
```

it's a GADT, where the type parameter is the type of the result of this request

- Core has a single way to issue a request

```
dataFetch :: DataSource req => req a -> Haxl a
```

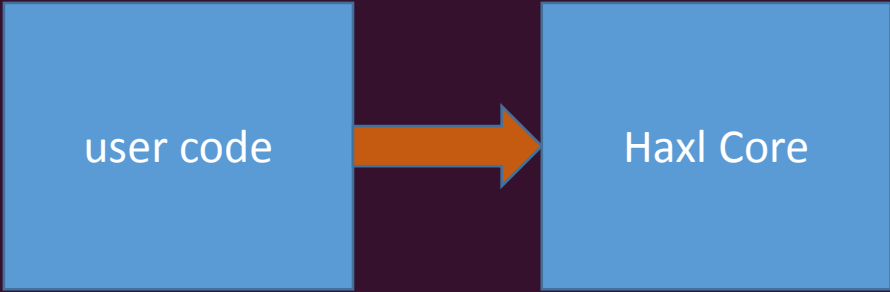
- Note how the result type matches up.

- Clean data source abstraction
- Means that we can plug in any set of data sources at *runtime*
 - e.g. mock data sources for testing and experimentation
 - core code can be built & tested independently

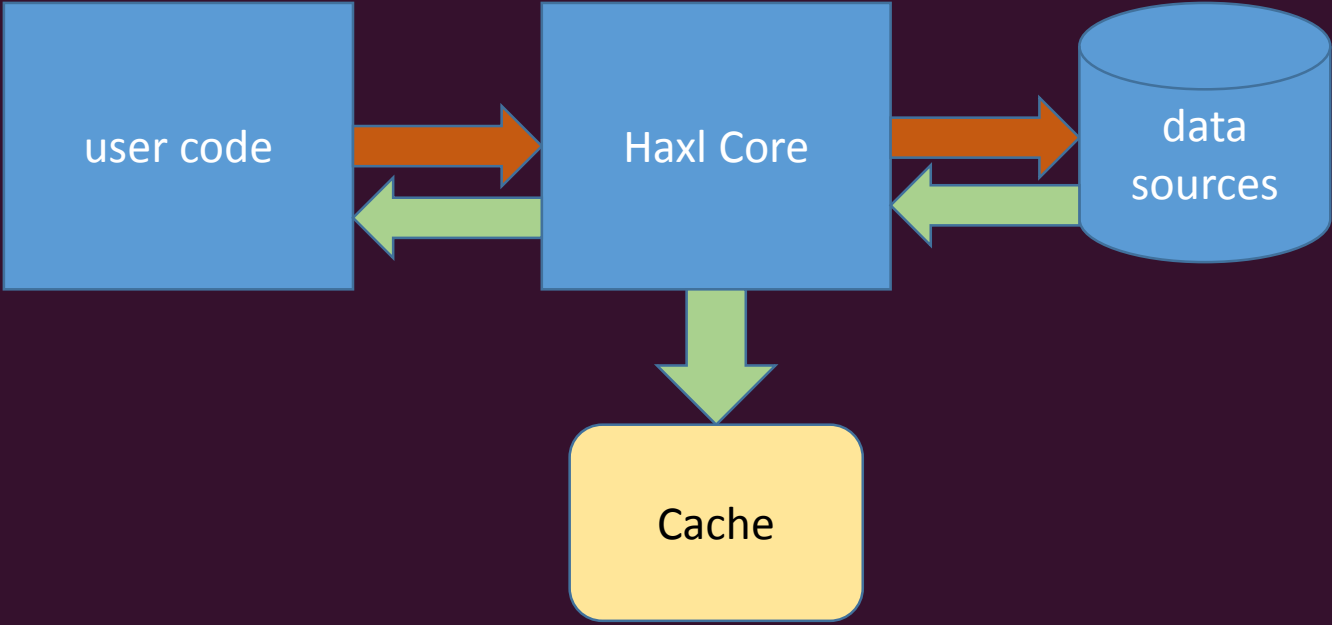
Replayability

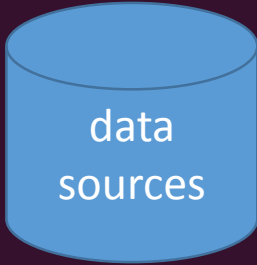
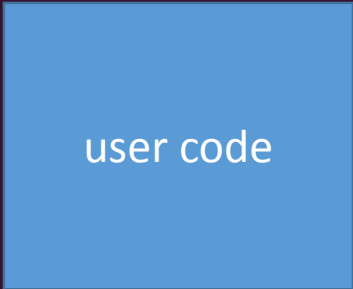
- The Haxl monad and the type system give us:
 - Guarantee of no side effects, except via `dataFetch`
 - Guarantee that everything is cached
 - The ability to replay requests...

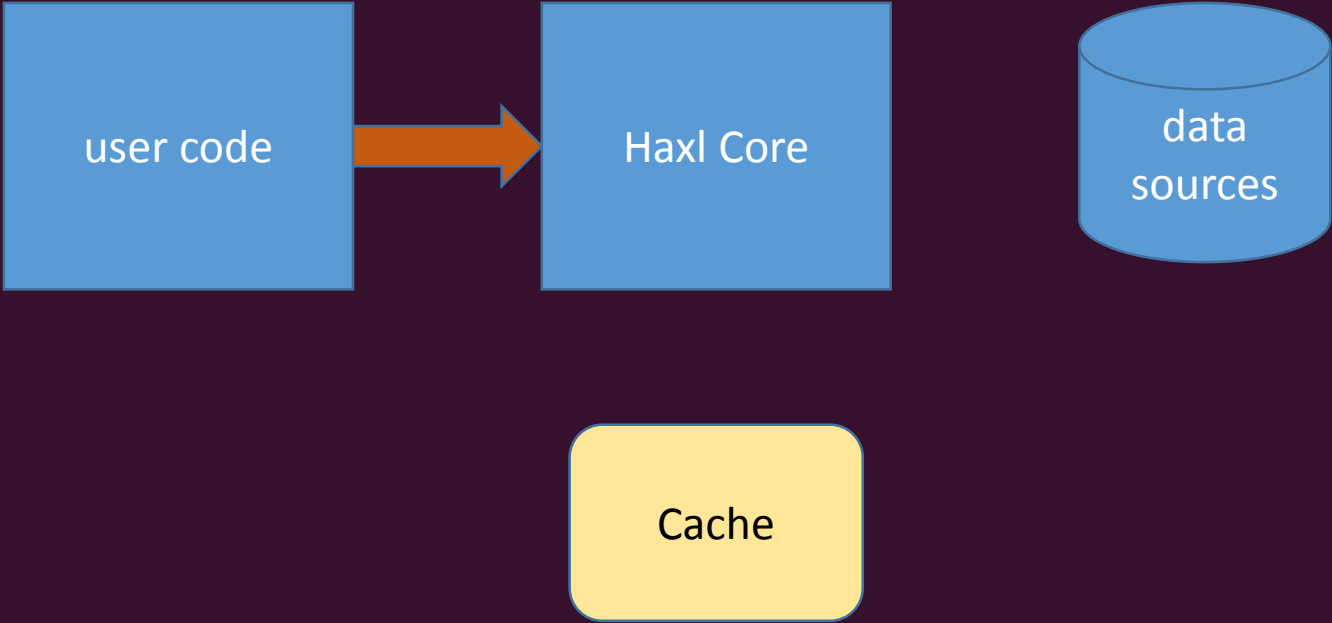
user code

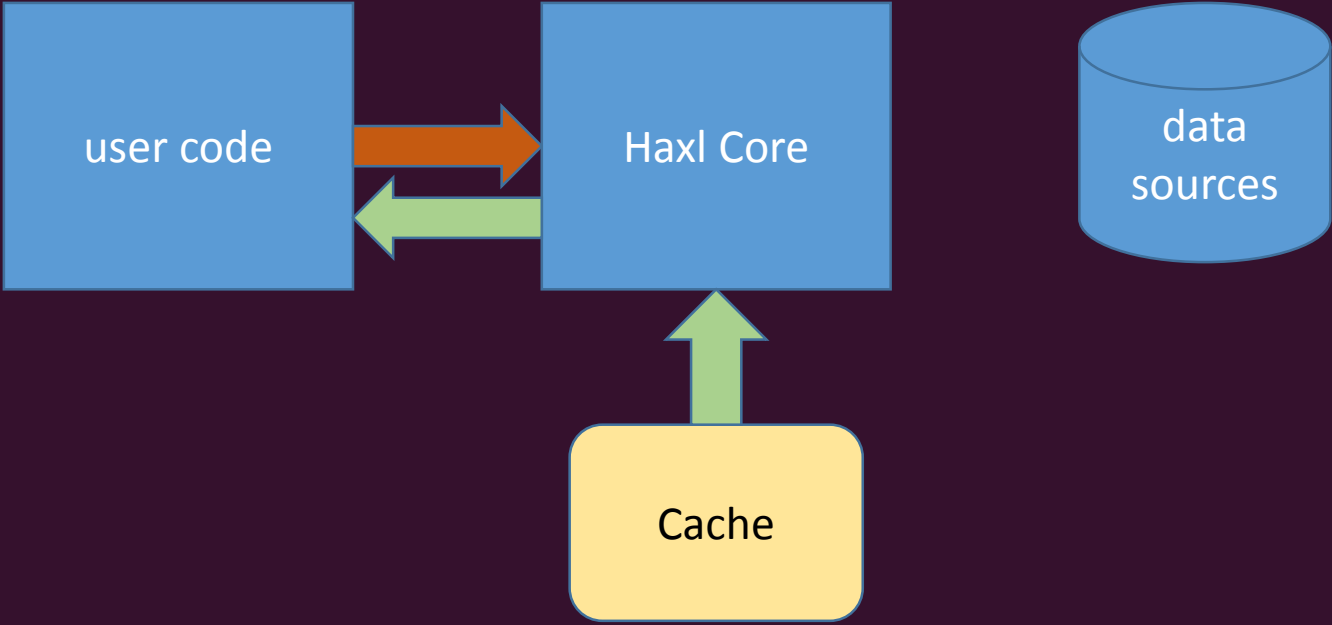


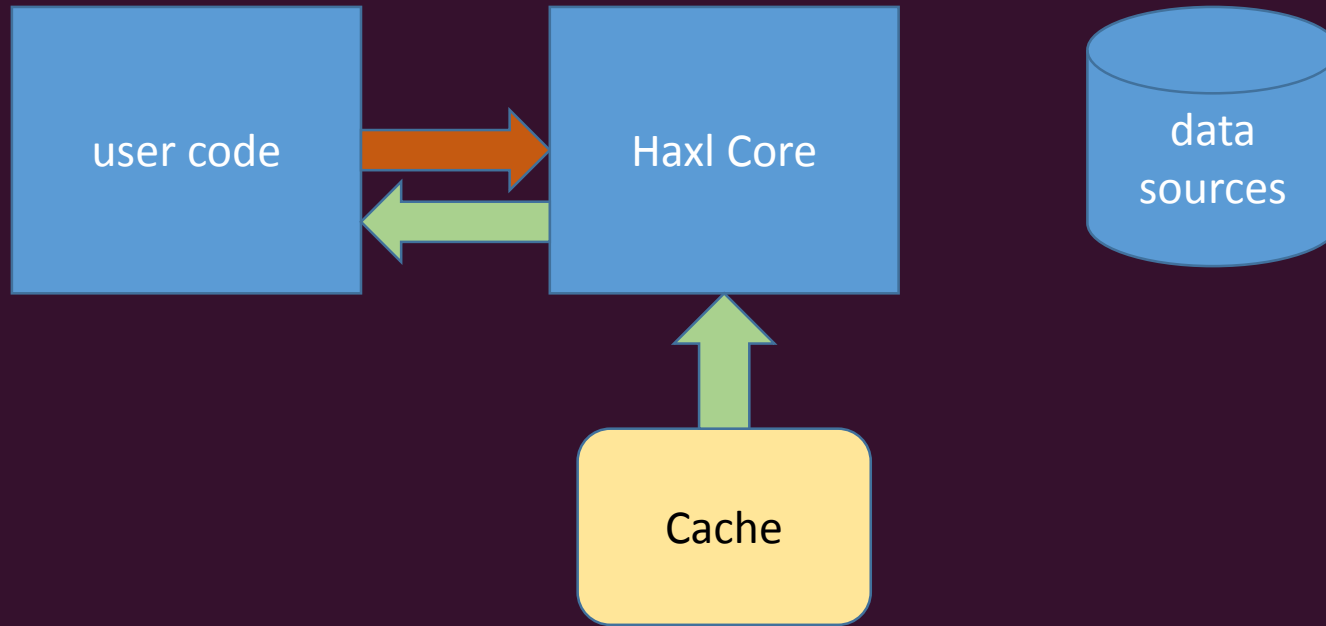












- The data sources change over time
- But if we *persist the cache*, we can re-run the user code and get the same results
- Great for
 - testing
 - fault diagnosis
 - profiling

Scaling

- Each server has lots of cores, pounded by requests from other boxes constantly.



Hot code swapping

- 1-2K machines, new code pushed many times per day
- Use GHC's built-in linker
 - Had to modify it to unload code
 - GC detects when it is safe to release old code
- We can swap in new code while requests are still running on the old code

Status

- Prototyped most features (including hot code swapping & scaling)
- Core is written
- We have a few data sources, more in the works
- Busy hooking it up to the infrastructure
- Can play around with the system in GHCi, including data sources

Questions?